



UC SANTA BARBARA

# Information *as* Control: The Role of Communication in Distributed Systems

Bryce L. Ferguson

For the ECE Department at UC Santa Cruz

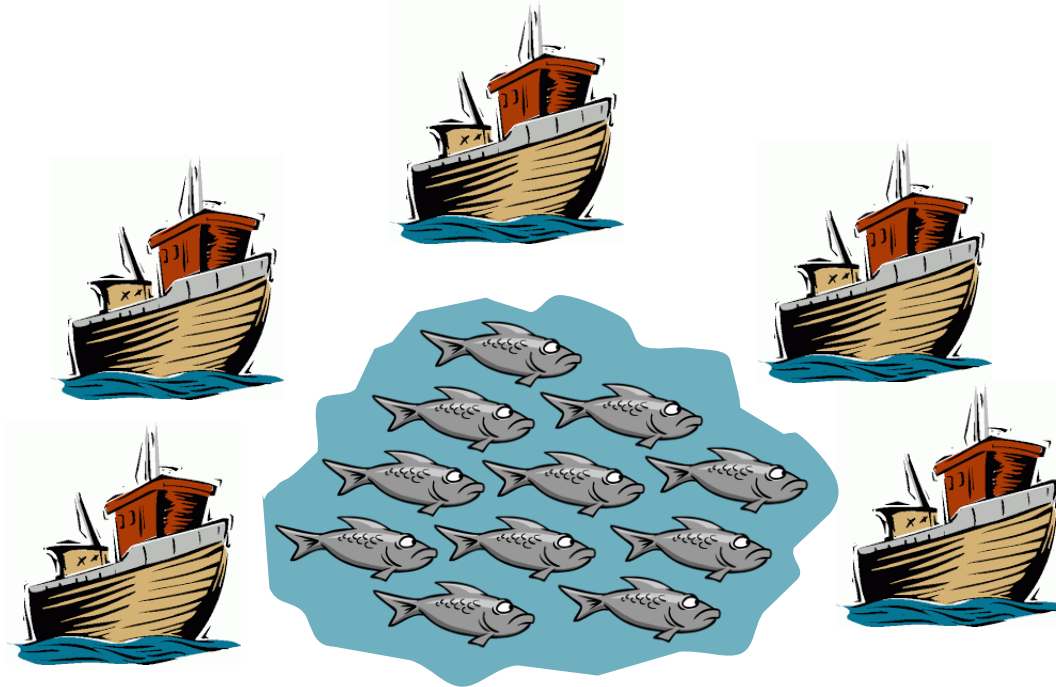
May 15<sup>th</sup>, 2023

Supported by NSF, ONR, AFOSR, and ARL

# Tragedy of the Commons & Distributed Control

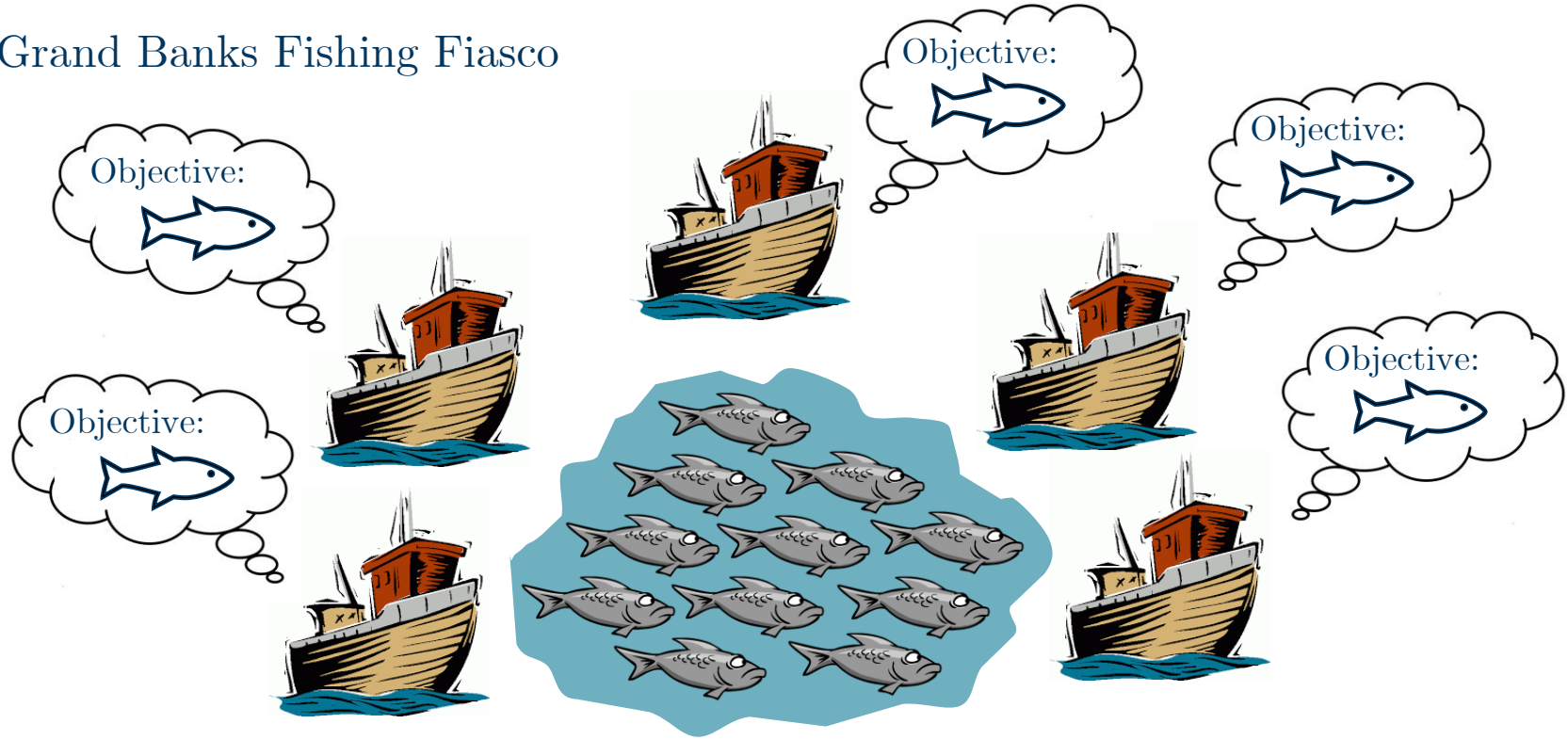
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## The Grand Banks Fishing Fiasco



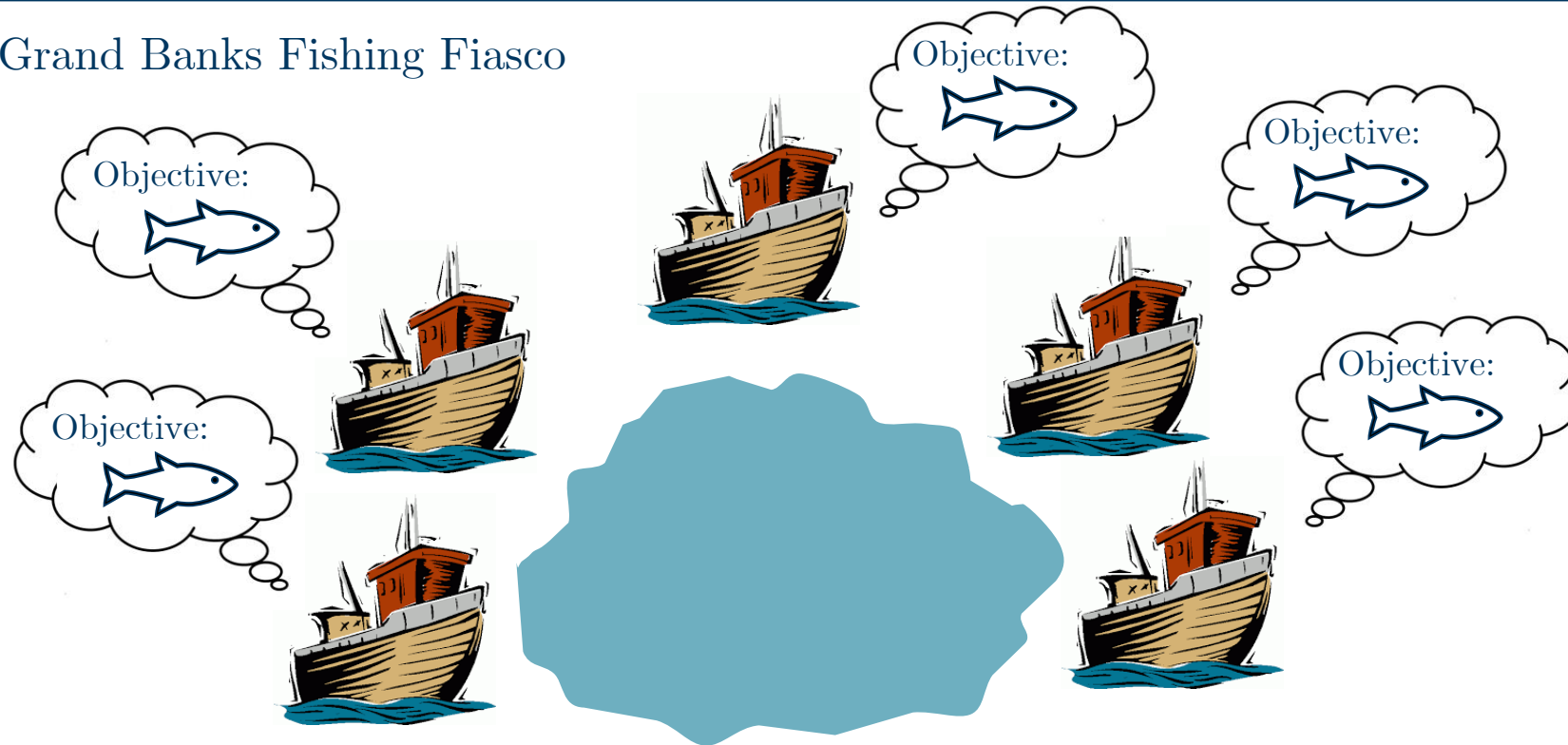
# Tragedy of the Commons & Distributed Control

The Grand Banks Fishing Fiasco



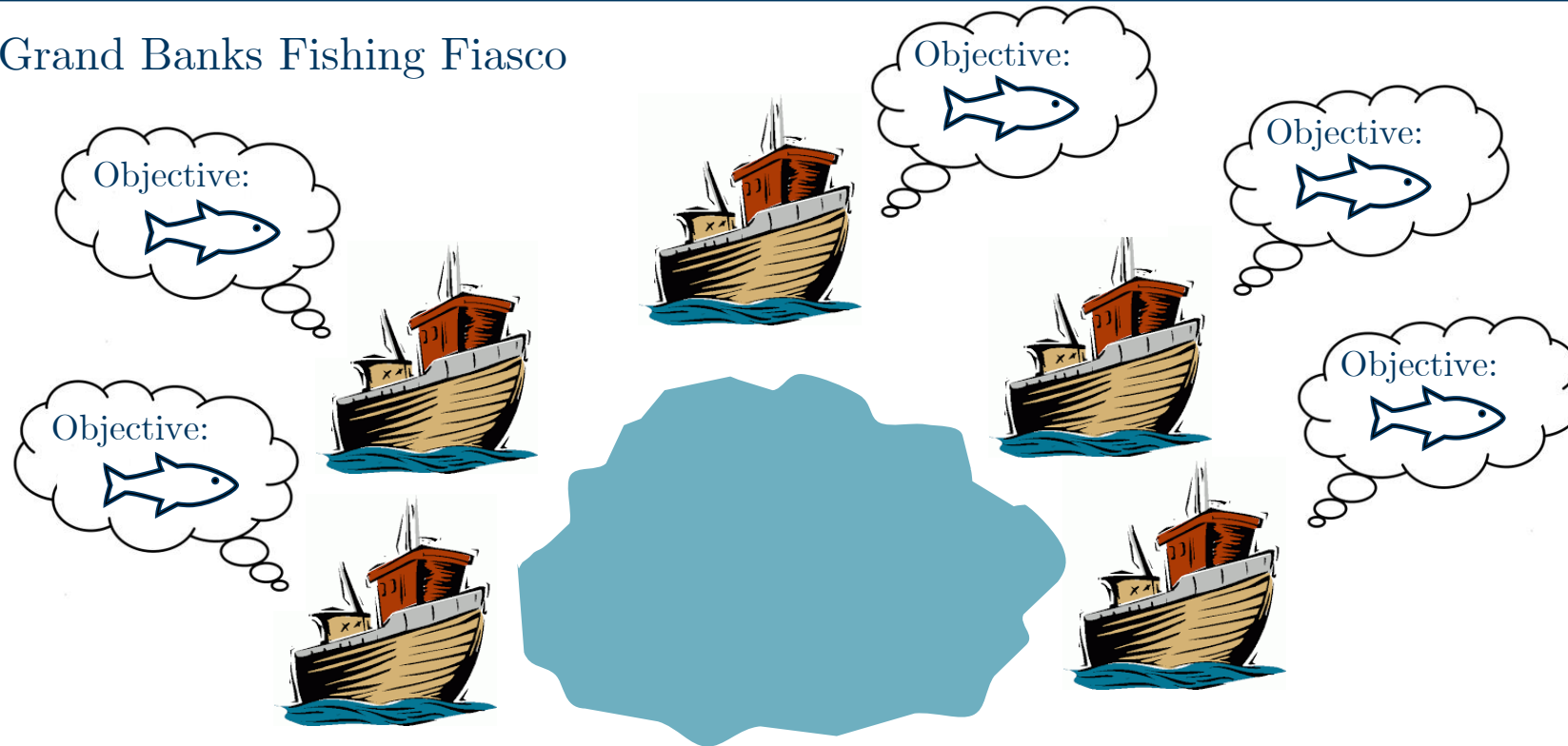
# Tragedy of the Commons & Distributed Control

## The Grand Banks Fishing Fiasco



# Tragedy of the Commons & Distributed Control

## The Grand Banks Fishing Fiasco



**Local Decision Making**  
(catch many fish)



**Sub-optimal Global Behavior**  
(fish population disappears)


# Tragedy of the Commons & Distributed Control

The Grand Banks Fishing Fiasco

## Environment

**THE BALTIMORE SUN**  
**Climate change and the "tragedy of the commons"**  
 By Dylan Selterman  
 Baltimore Sun • Jan 10, 2023

**CLIMATE**  
**Deforestation in Brazilian Amazon hits tragic record in 2022**  
 PUBLISHED WED, JUL 13 2022-9:36 AM EDT




Objective:

## Finance/Business

**Gartner: Corporate Sustainability Suffers From Tragedy of the Commons**  
 Emma Chervek | Reporter March 4, 2023 7:00 AM


**FORBES**  
**The Tragedy Of The Commons: How A Life Crisis Effects Family Wealth**  
 Matthew Erskine Contributor ©



## Traffic/Congestion

**The New York Times**  
**Data Driving New Approaches to Transportation**  
 Analyzing digital streams of information from electric scooters and motor-assisted bicycles is helping solve travel congestion issues.  
 By Norman Meyerson  
 Feb. 5, 2020  
 This article is part of a series on how businesses evolve.  
 Five seconds after a 1000 cars operated databank is noting the location. It

**FINANCIAL TIMES**  
**Uber exemplifies the Tragedy of the Commons**  
 From Tad Borek, San Francisco, CA, US — Wednesday's most read letter

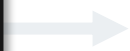


## Cloud Computing Services

**DATA ECONOMY**  
**ChatGPT is at capacity right now; yes, it is really annoying**  
 Just like any other service...

**VentureBeat**  
**3 ways data teams can avoid a tragedy of the cloud commons**  
 Clinton Ford, Unravel Data

Lately, the "ChatGPT is at capacity" message is an outstanding chatbot. "ChatGPT is at capacity" when the chatbot server is down, network trouble, etc. Know that the chatbot's servers are overwhelmed by the volume of requests they're receiving when you encounter any of these issues. However, a few simple steps can fix the widely used chatbot.



# Tragedy of the Commons & Distributed Control

Internet Connected Devices



**Local Decision Making**  
(device algorithms)



**Sub-optimal Global Behavior**  
(system errors/inefficiencies/privacy)

# Tragedy of the Commons & Distributed Control

Internet Connected Devices



**Local Decision Making**  
(device algorithms)



**Sub-optimal Global Behavior**  
(system errors/inefficiencies/privacy)

How do we elicit *coordination* among many *decision-makers*?

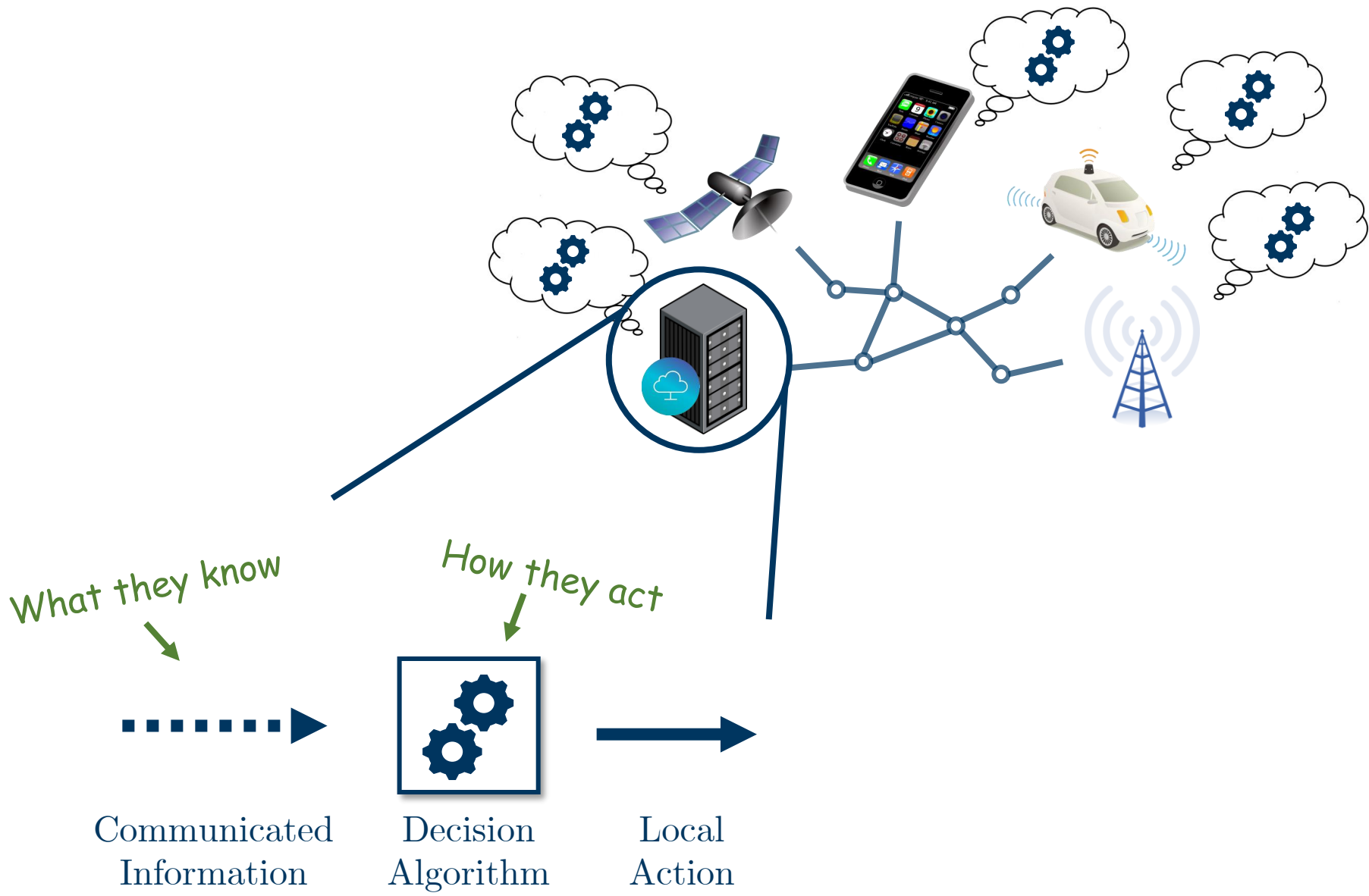


# Coordination in Distributed Systems

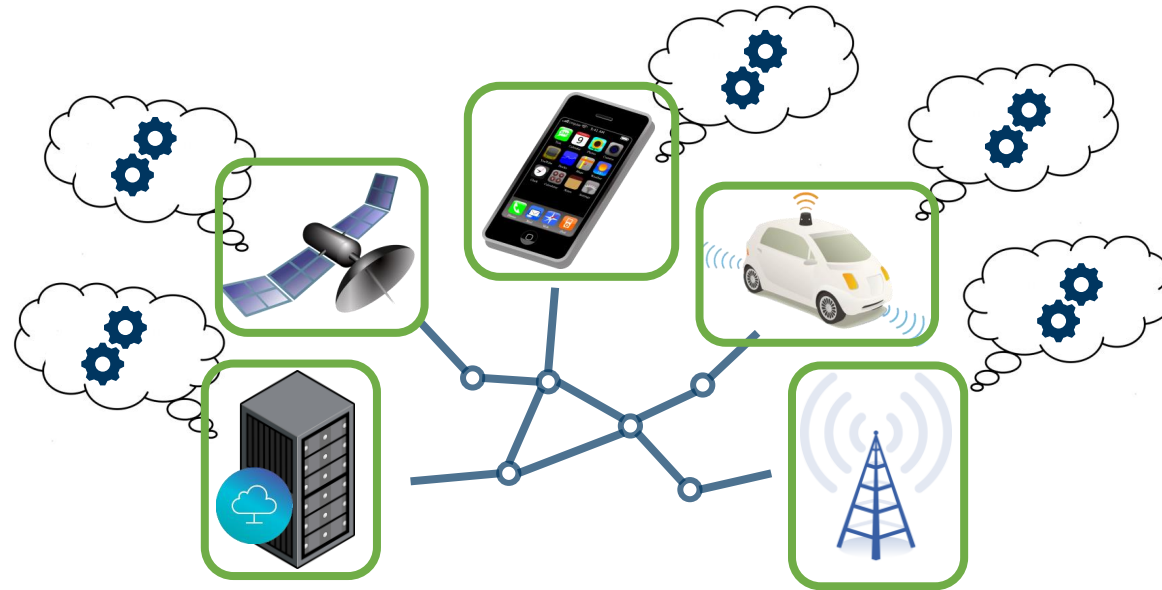
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# Coordination in Distributed Systems

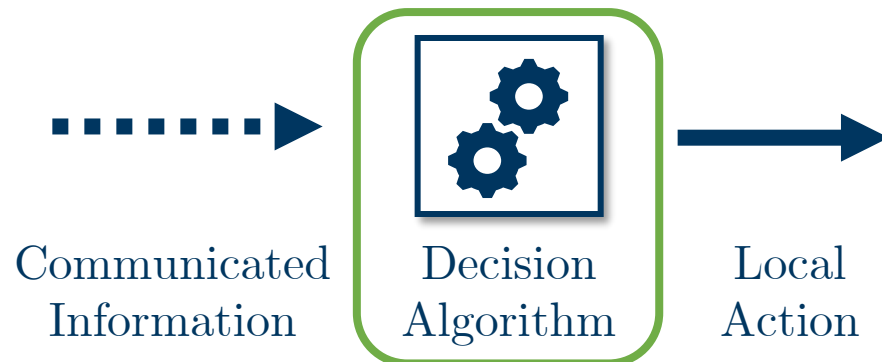


# Coordination in Distributed Systems



## Objective Design

- Design local decision algorithm
- Utilize limited information

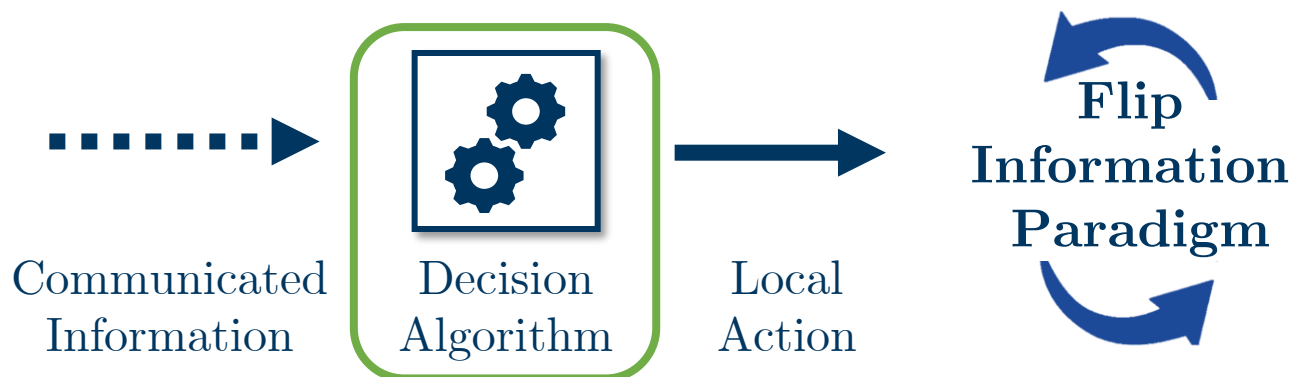


# Coordination in Distributed Systems



## Objective Design

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# Coordination in Distributed Systems



## Objective Design

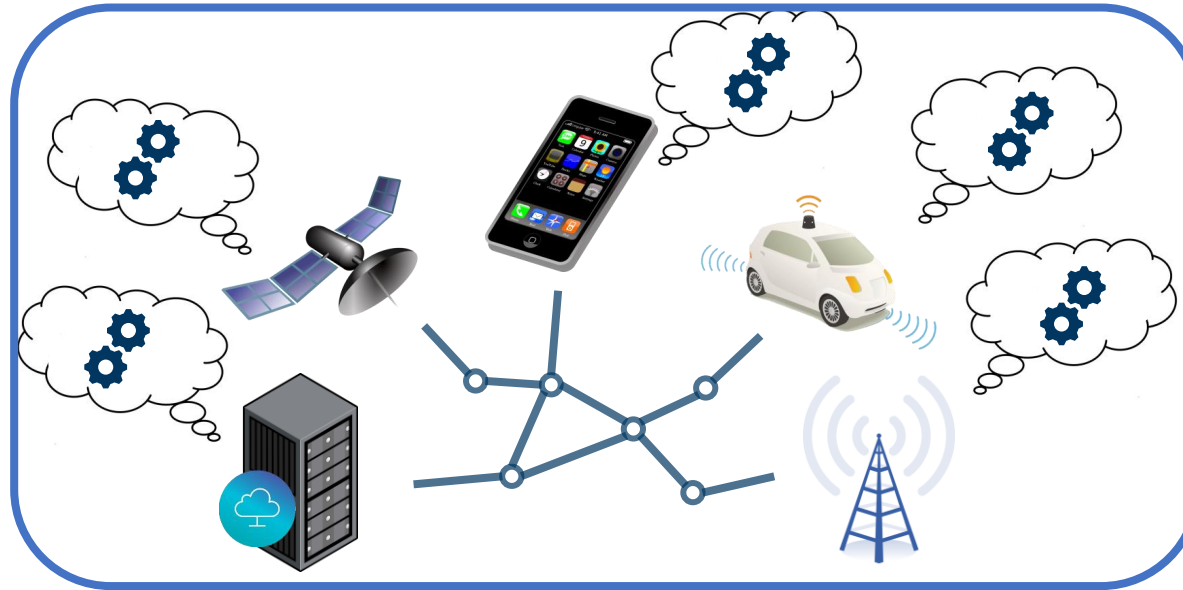
- Design local decision algorithm
- Utilize limited information

## Information Design

- Design communication structure
- Alter inputs to existing decision algorithms

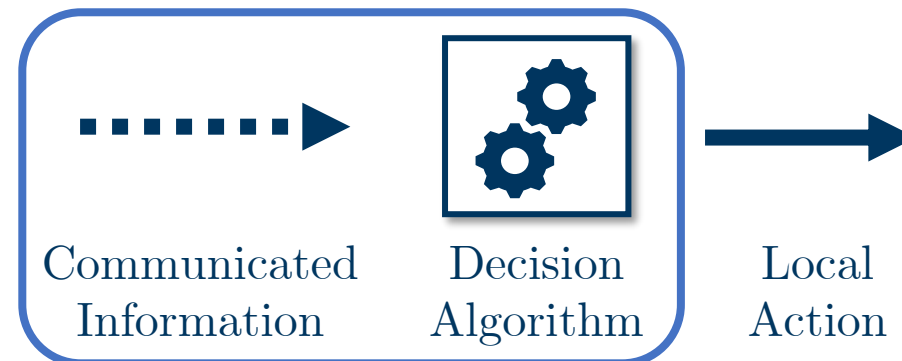


# Coordination in Distributed Systems



## Intelligent Information System Design

- Design communication and local algorithms together
- Improved performance at the cost of increased complexity



# Types of Information-Communication

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## Information Sharing

Collaboratively exchange information

## Information Provisioning

Strategically send out information

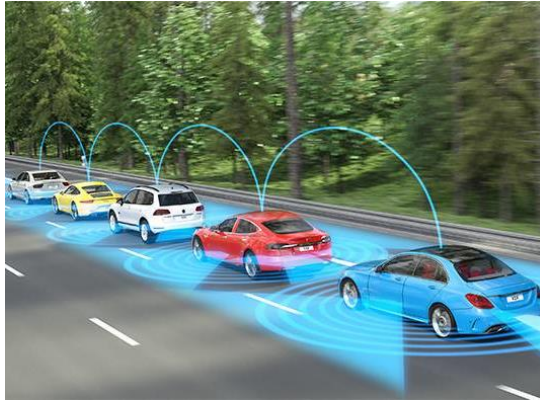
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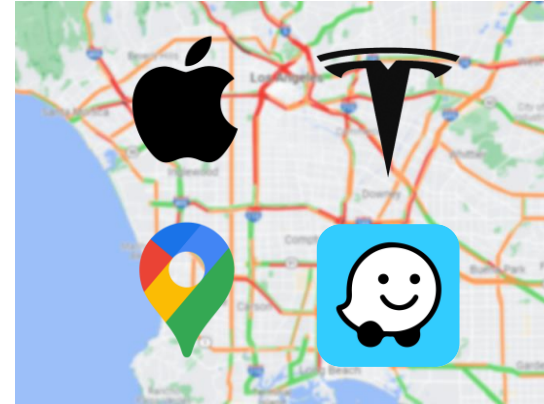
E.g.,  
Autonomous  
Driving



Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling

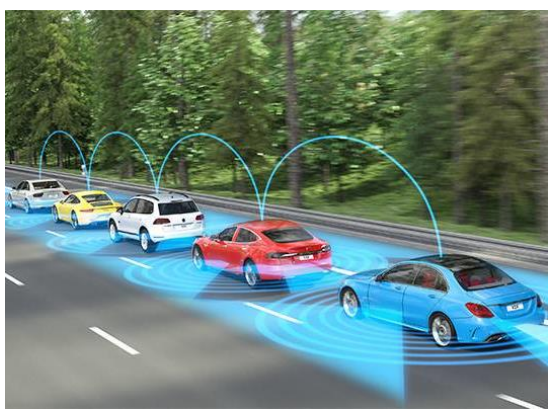


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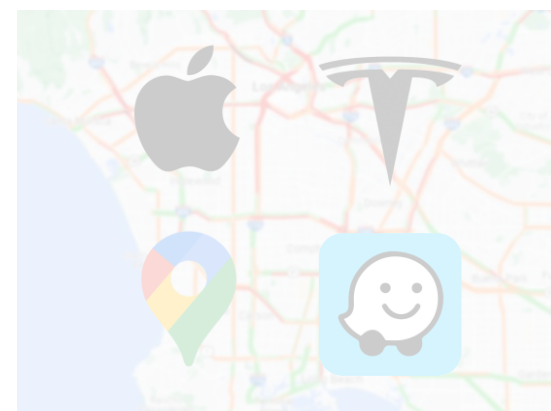
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Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



Local information  
Local decision



Local information  
Local decision



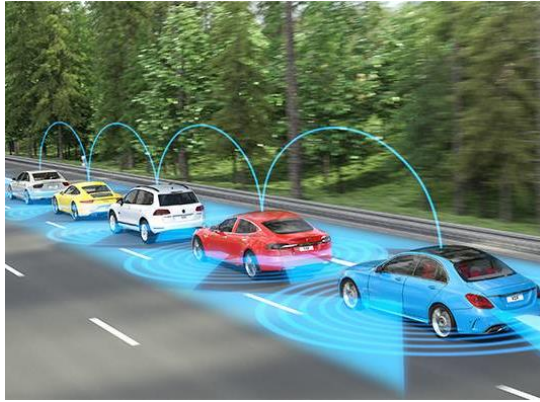
Local information  
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# Types of Information-Communication

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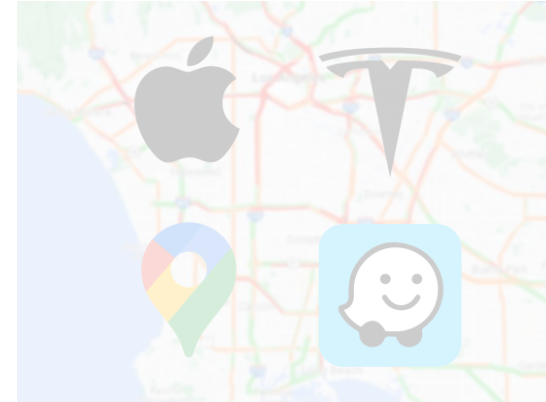
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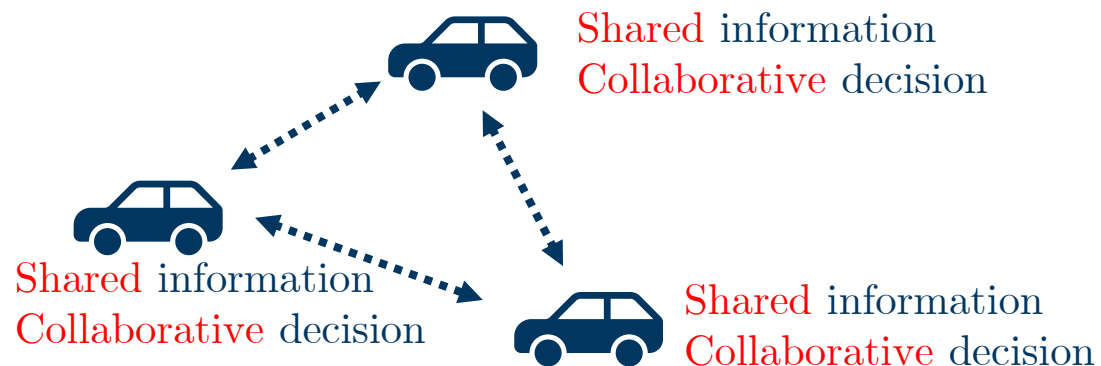
Vehicle Platooning

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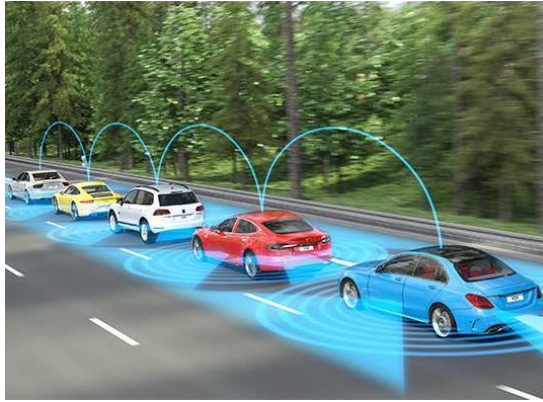


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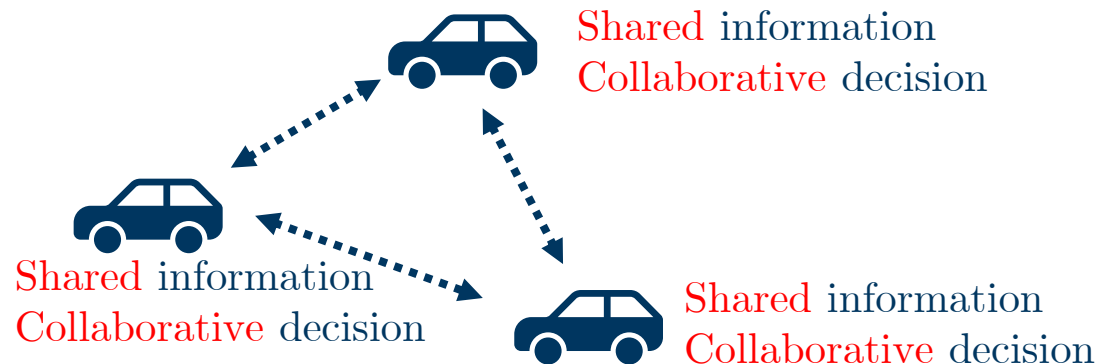
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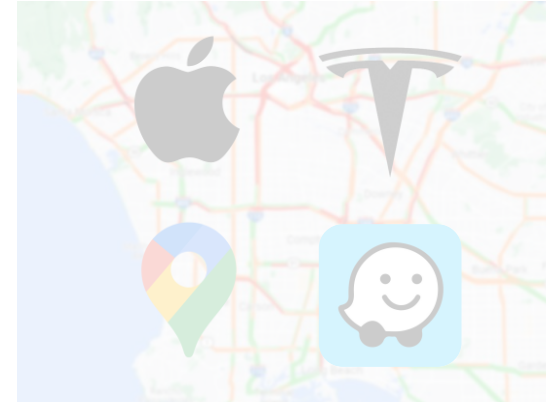


Vehicle Platooning



## Information Provisioning

Strategically send out information



Traffic Signaling

What information is shared?

How is information used?

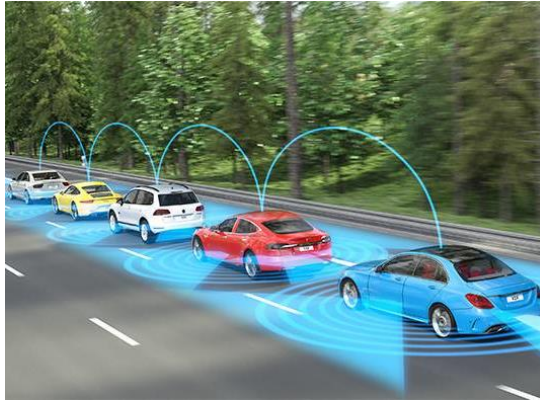
Network Coordination [Jackson & Watts '02]  
Collaborative Control [Fong et. al. '03]  
Coalition Formation [Ray & Vohra '15]

# Types of Information-Communication

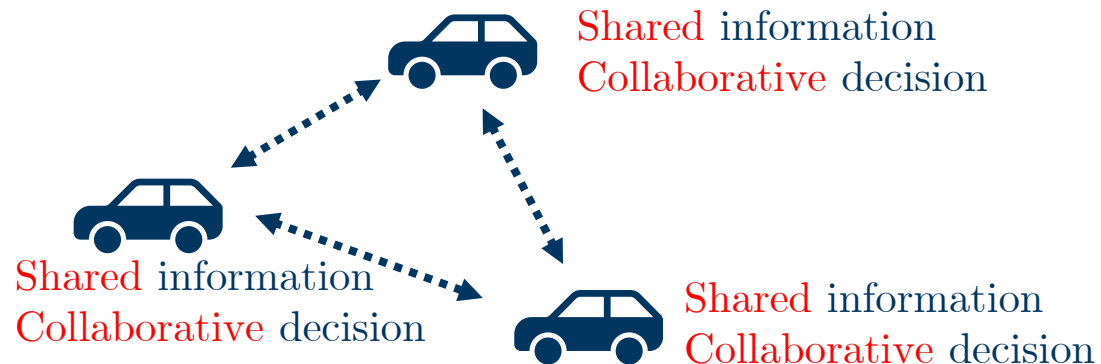
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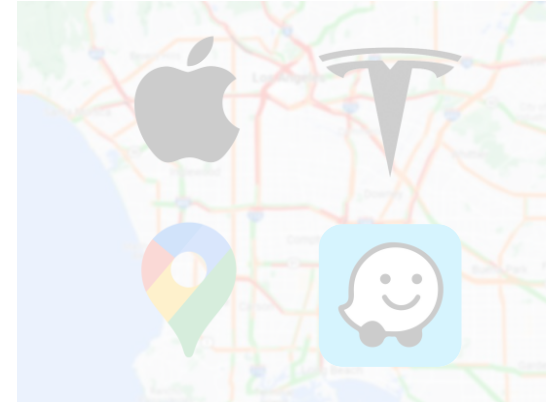


Vehicle Platooning



## Information Provisioning

Strategically send out information



Traffic Signaling

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Network Coordination [Jackson & Watts '02]  
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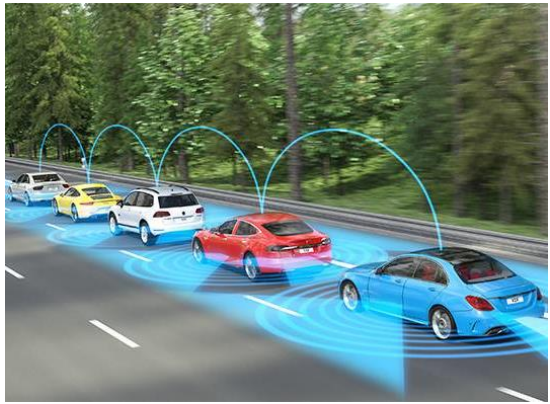
Focus on convergence to equilibrium behavior

# Types of Information-Communication

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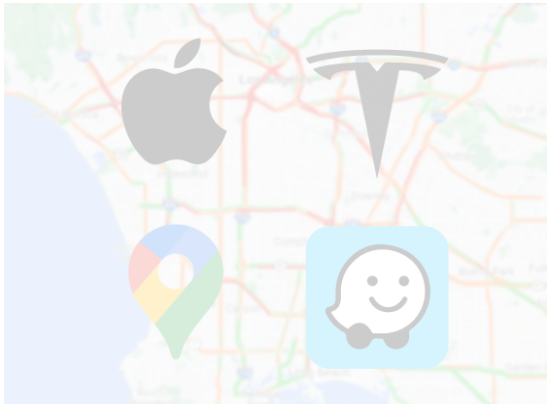
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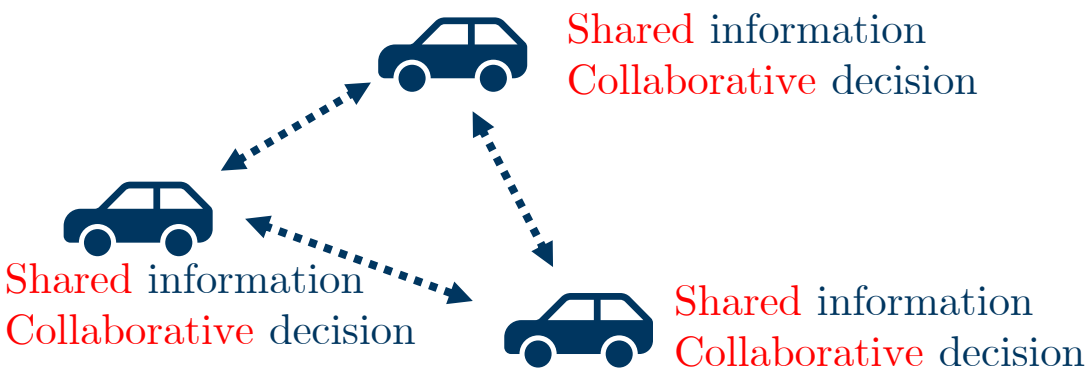
Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



What information is shared?

How is information used?

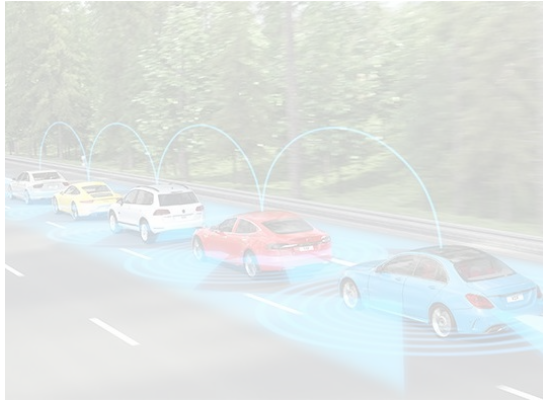
My work:

Quantify the *benefit* of collaborative information sharing to system performance

# Types of Information-Communication

## Information Sharing

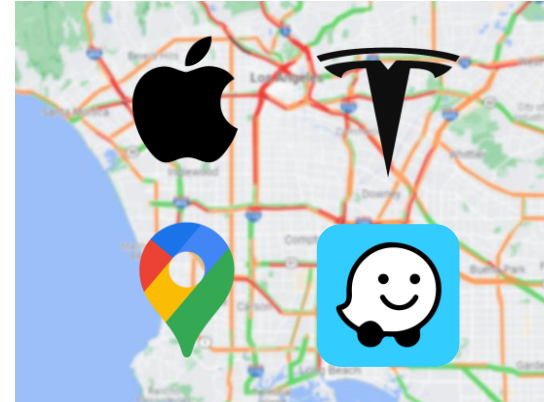
Collaboratively exchange information



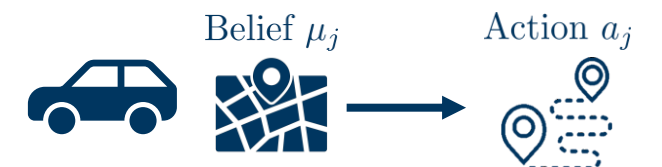
Vehicle Platooning

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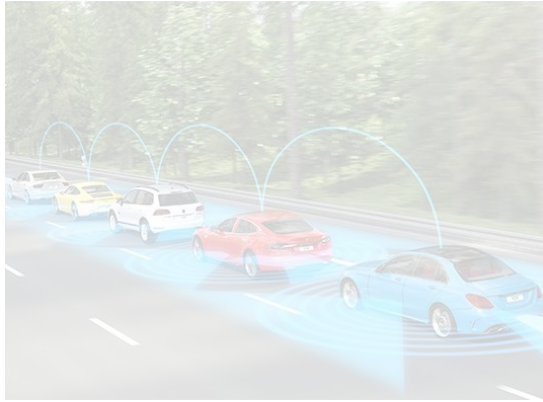
Traffic Signaling



# Types of Information-Communication

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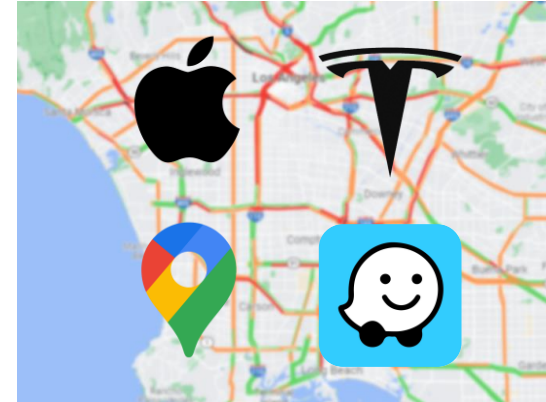
Collaboratively exchange information



Vehicle Platooning

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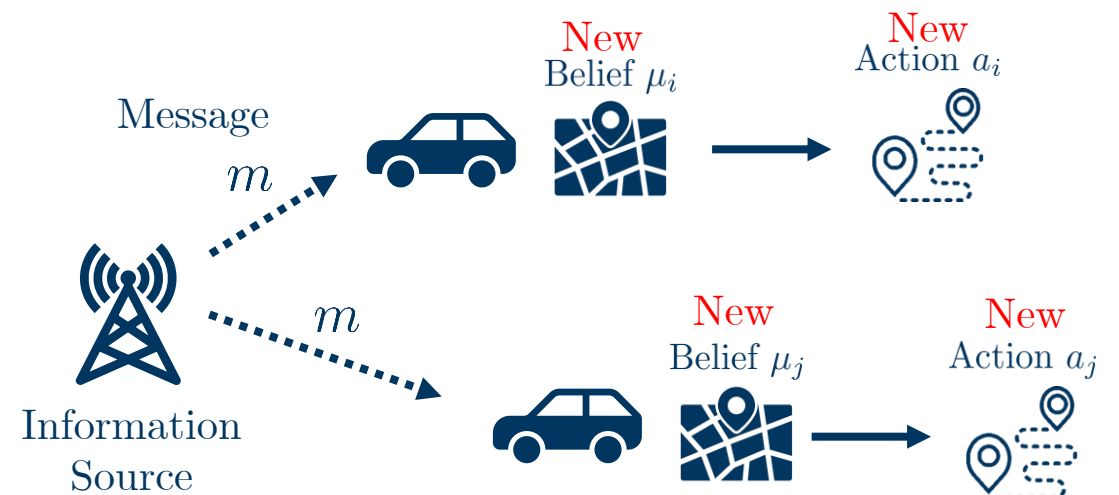
Traffic Signaling

Social Systems:

- Matching Markets [Ostrovsky et. al. '10]
- Social Media [Romero et. al. '11]
- Elections [Alonso et. al. '16]

Bayesian Persuasion: (more general framework)

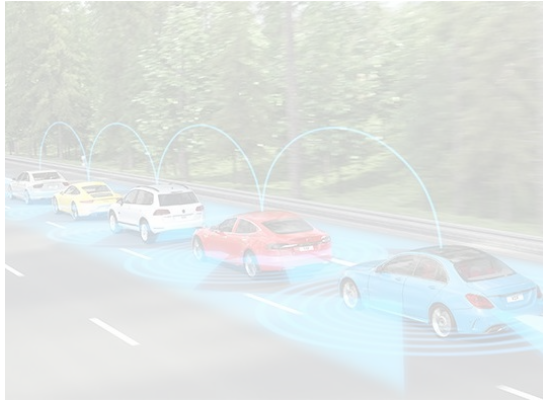
- Public [Kamenica & Gentzkow '11]
- Private [Arieli & Babichenko '19]



# Types of Information-Communication

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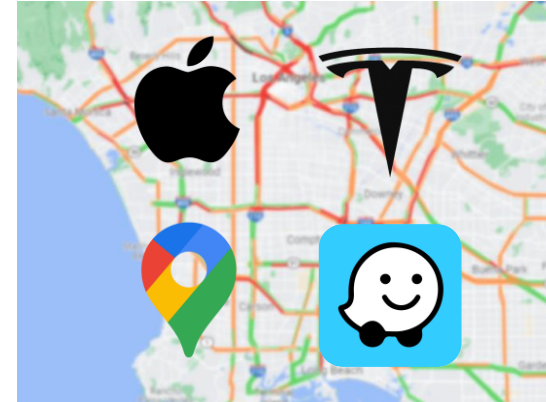
Collaboratively exchange information



Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling

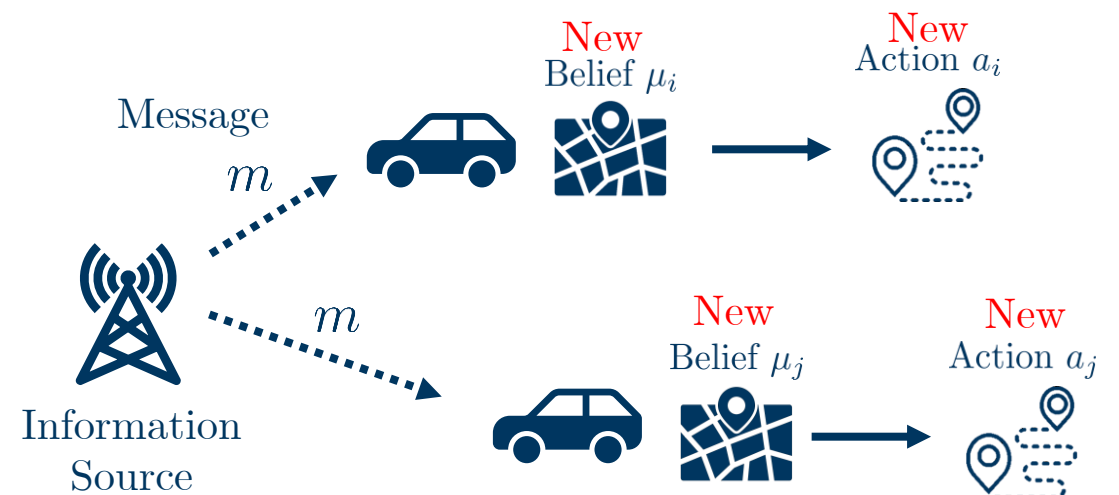
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Bayesian Persuasion: (more general framework)

- Public [Kamenica & Gentzkow '11]
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Identify scenarios where revealing info helps / hurts

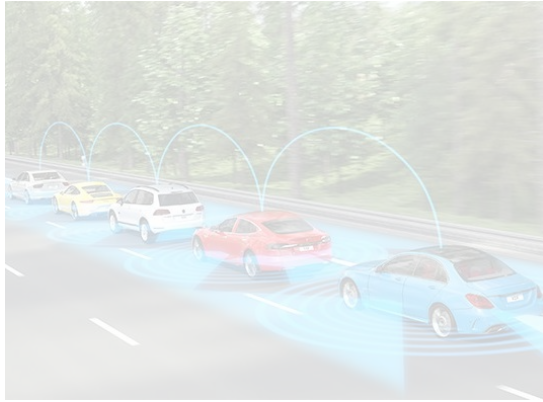




# Types of Information-Communication

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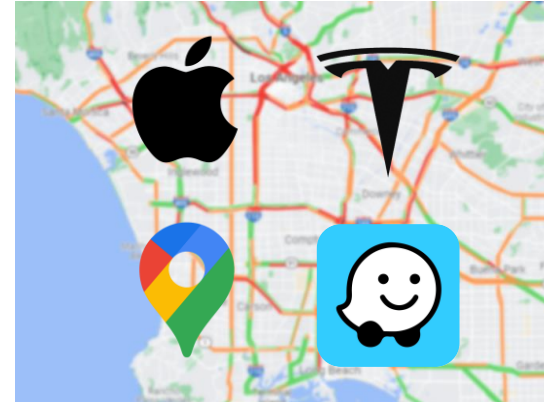
Collaboratively exchange information



Vehicle Platooning

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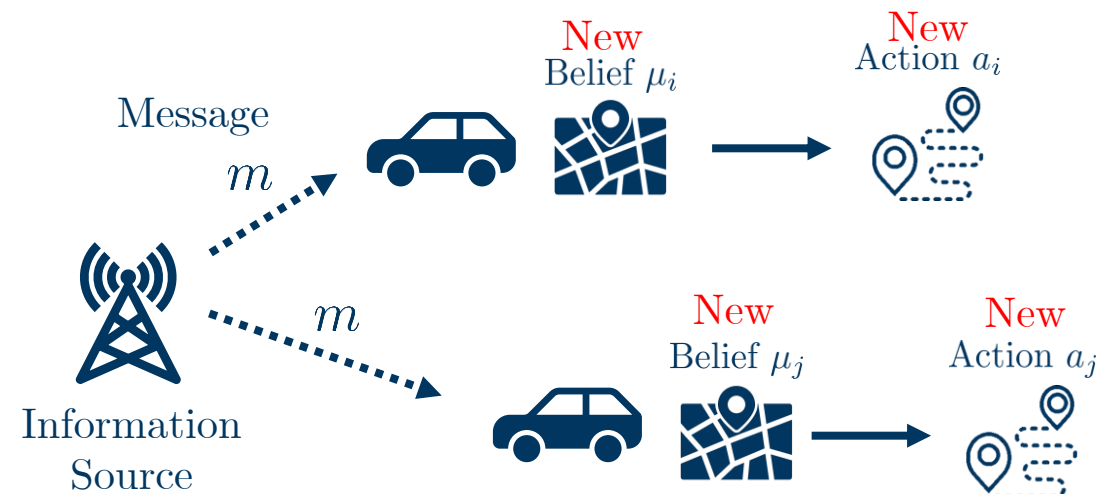
Strategically send out information



Traffic Signaling

My work:

Quantify how information provisioning affects *algorithmic guarantees* of distributed systems



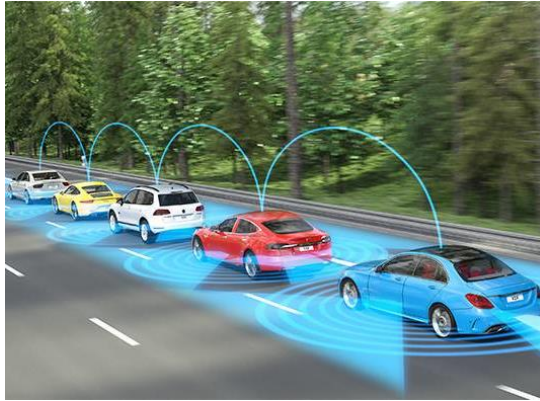
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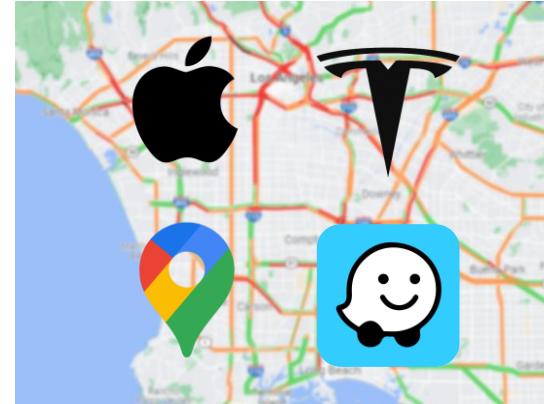
E.g.,  
Autonomous  
Driving



Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling

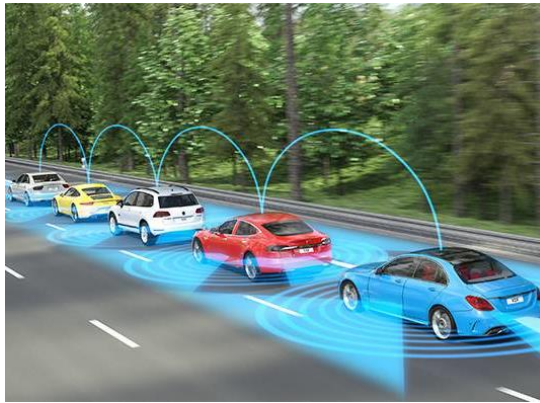
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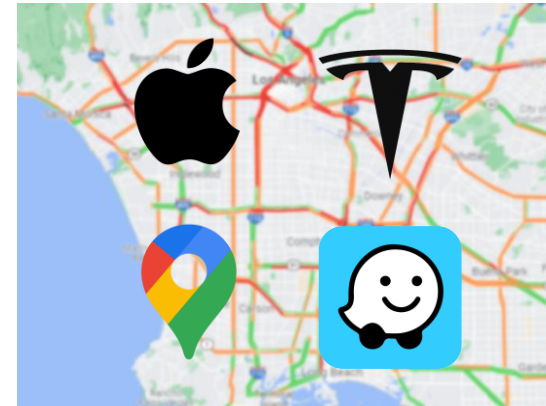
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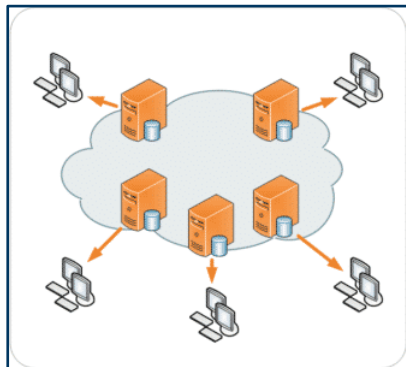
Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



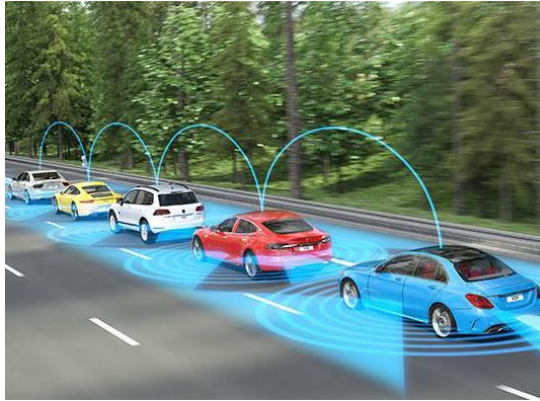
Cloud computing

# Types of Information-Communication

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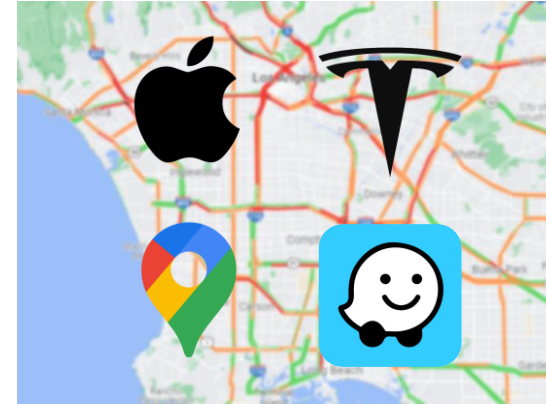
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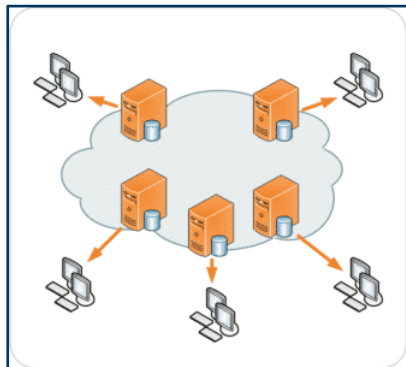
Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



Cloud computing



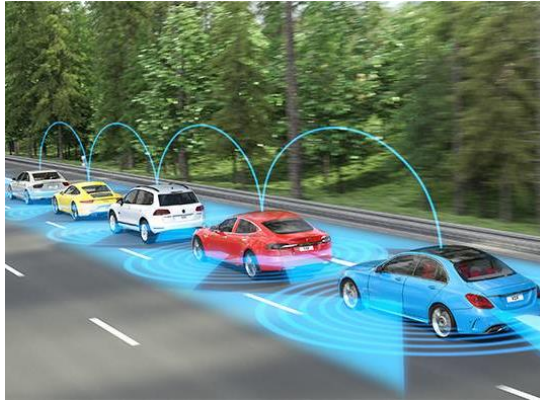
Supply-chain

# Types of Information-Communication

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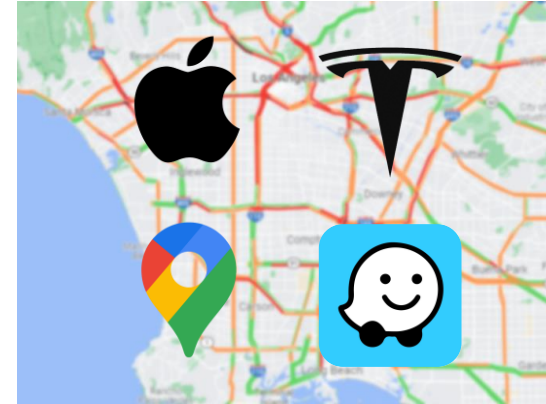
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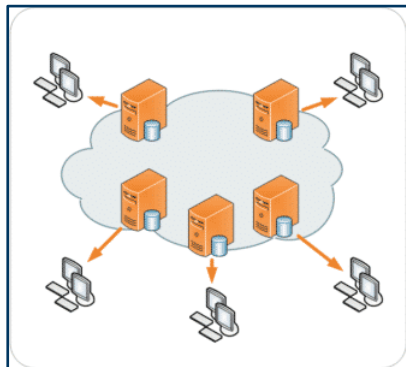
Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



Cloud computing



Supply-chain



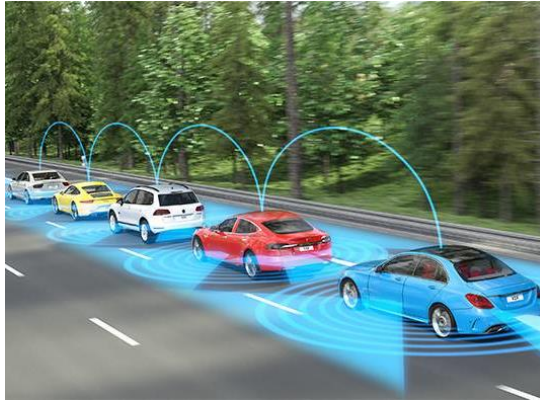
Power Grids

# Types of Information-Communication

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Collaboratively exchange information

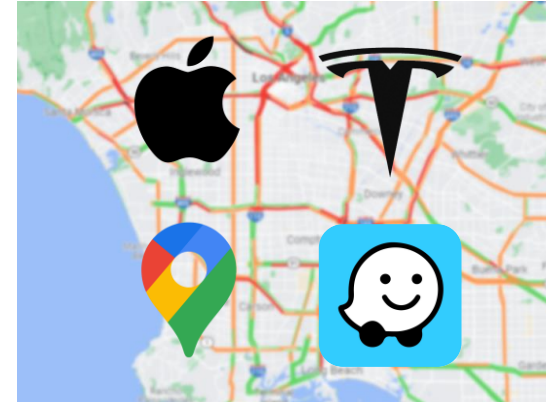
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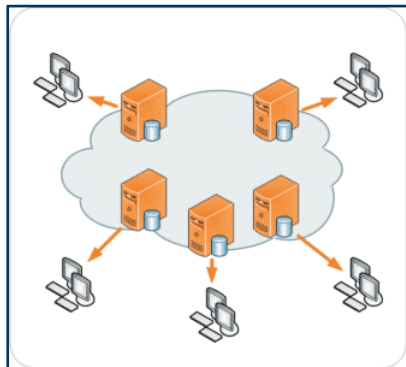
Vehicle Platooning

## Information Provisioning

Strategically send out information



Traffic Signaling



Cloud computing



Supply-chain



Power Grids



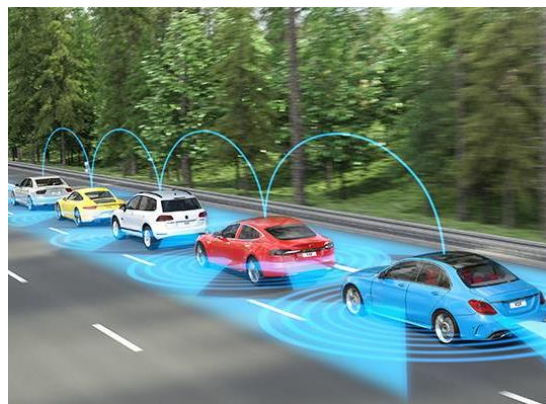
Fleet Robotics

# Types of Information-Communication

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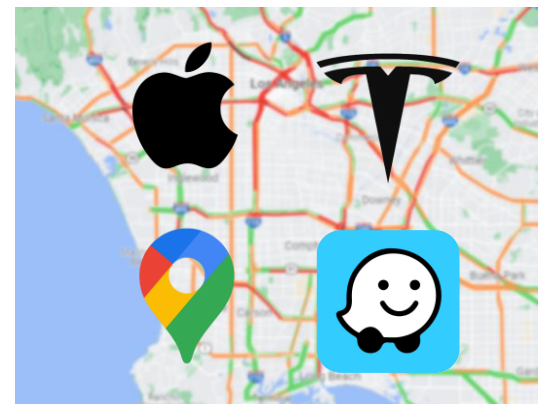
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Driving



Vehicle Platooning

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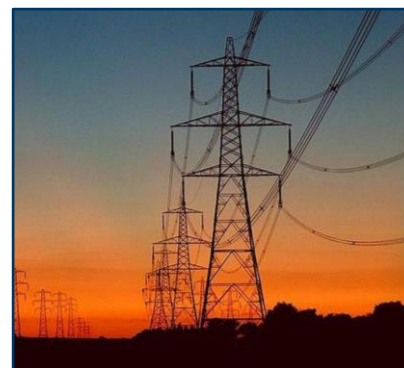
Traffic Signaling



Cloud computing



Supply-chain



Power Grids



Fleet Robotics



Military Ops.

# Types of Information-Communication

## Information Sharing

Collaboratively exchange information

## Information Provisioning

Strategically send out information

E.g.,  
Autonomous  
Driving

## The Effects of Information-Communication

Understand the *benefits* and *costs* of increasing the amount of communication in a distributed system

Benefit: improved system performance,  
better local decision making

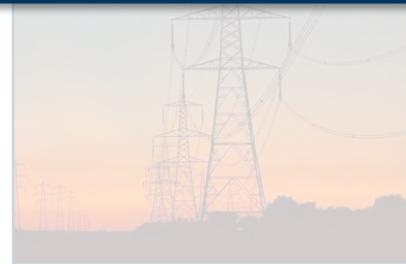
Cost: physical/computational burdens,  
unexpected behavior can emerge



Cloud computing



Supply-chain



Power Grids



Fleet Robotics



Military Ops.



# Unreliable Communicators

Unreliable

Information Sharing

Collaboratively exchange information

E.g.,  
Autonomous  
Driving



Unreliable

Information Provisioning

Strategically send out information



Simon Weckert

My work:

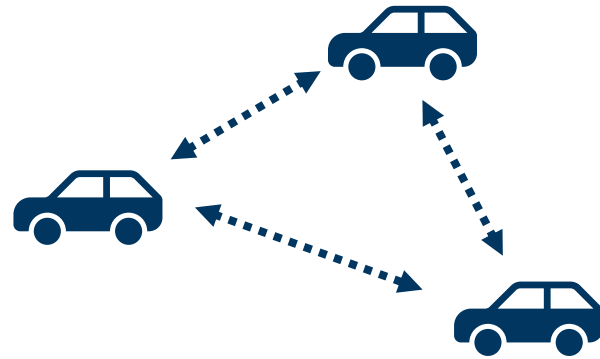
What are the effects of *unreliable communicators*  
and what can we do to mitigate them?

# Outline

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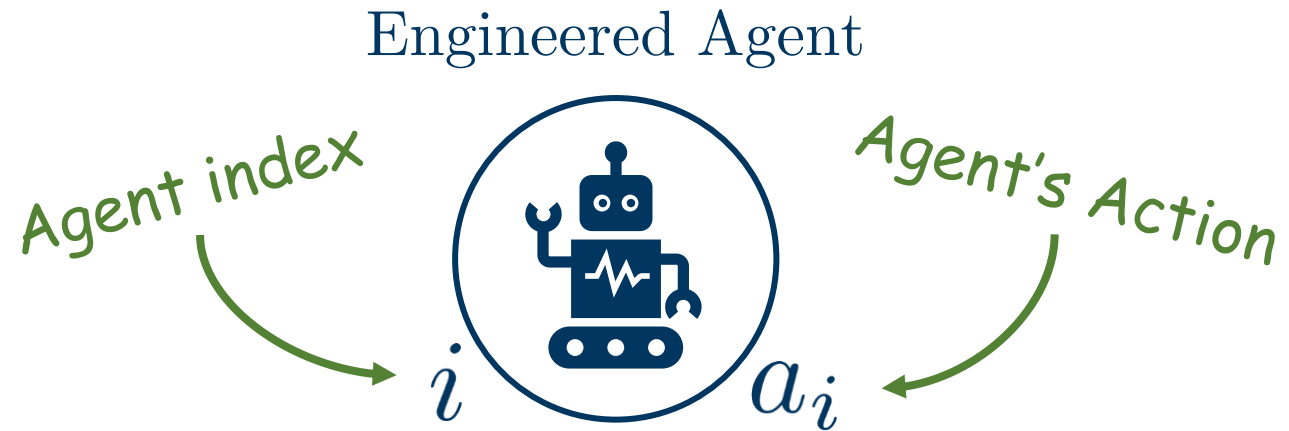
- I. Information Sharing
- II. Information Provisioning
- III. Unreliable Communicators
- IV. Conclusion and Directions

# I. Information Sharing



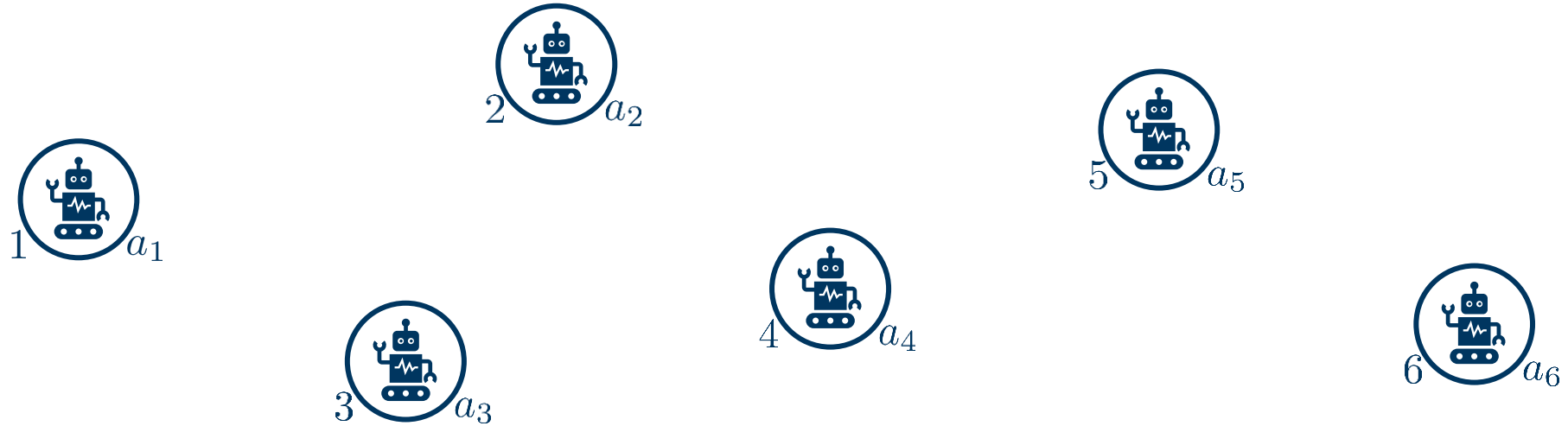
# Communication Design in Engineered System

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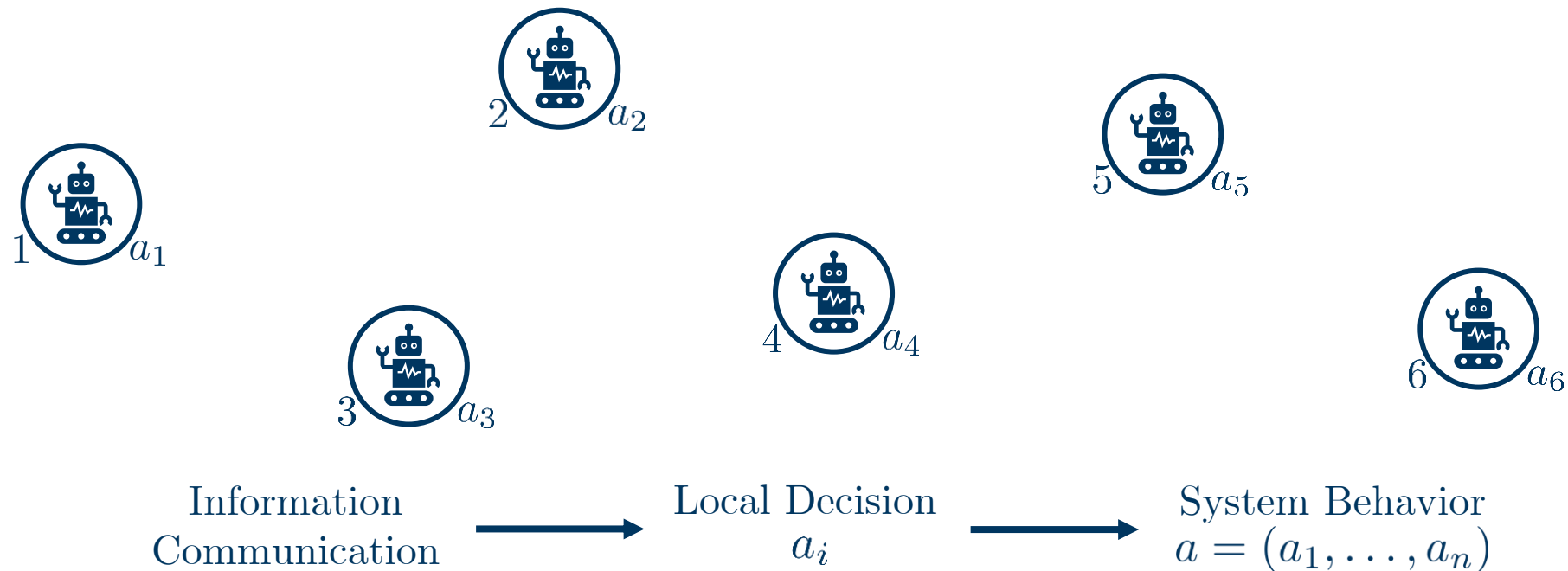
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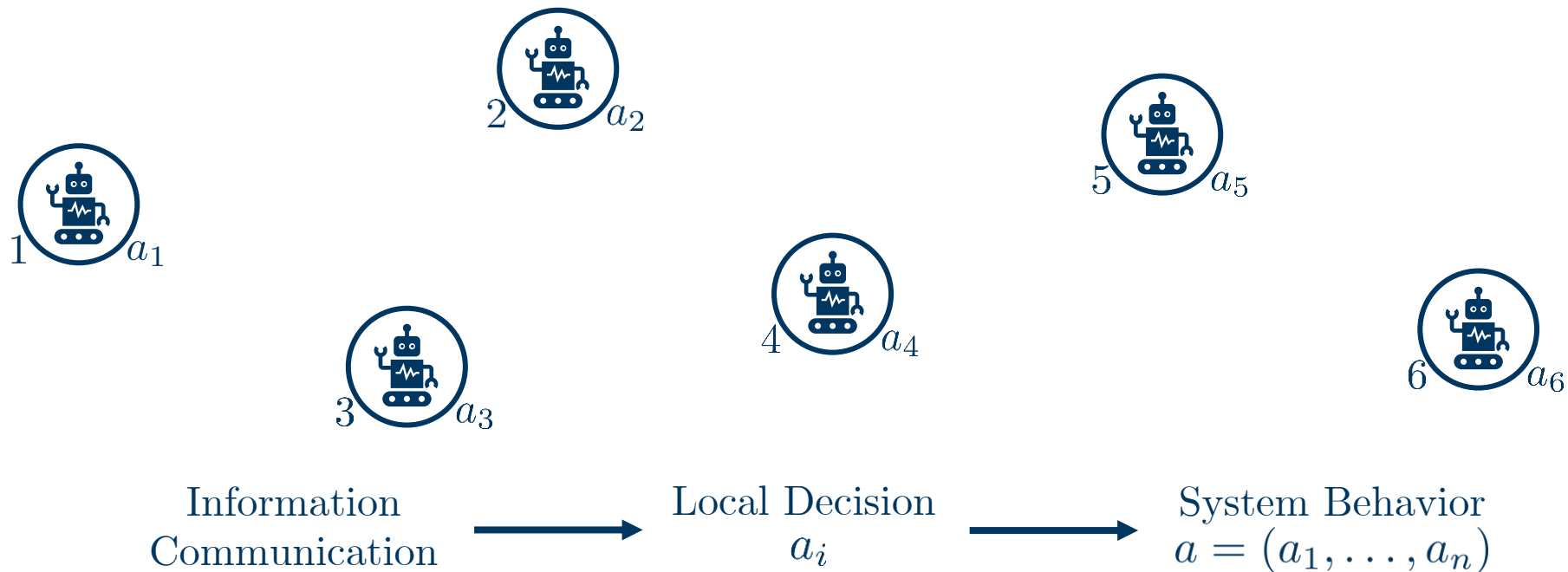
# Communication Design in Engineered System

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# Communication Design in Engineered System

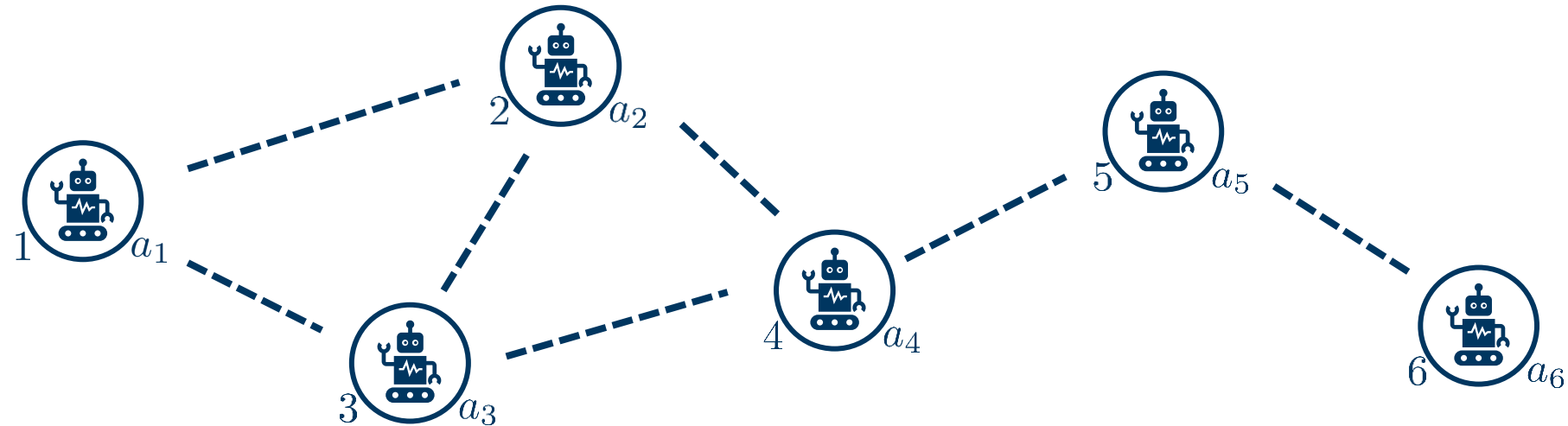
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How should agents communicate?

# Communication Design in Engineered System

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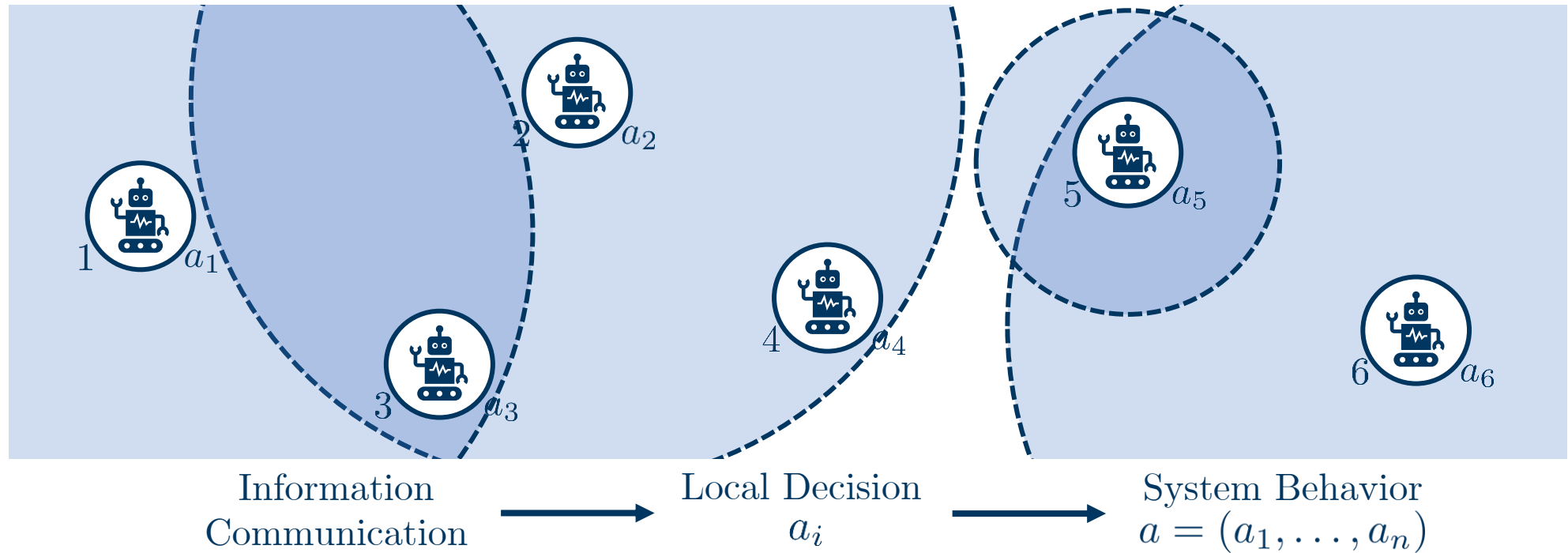


How should agents communicate?

- Networked communication?



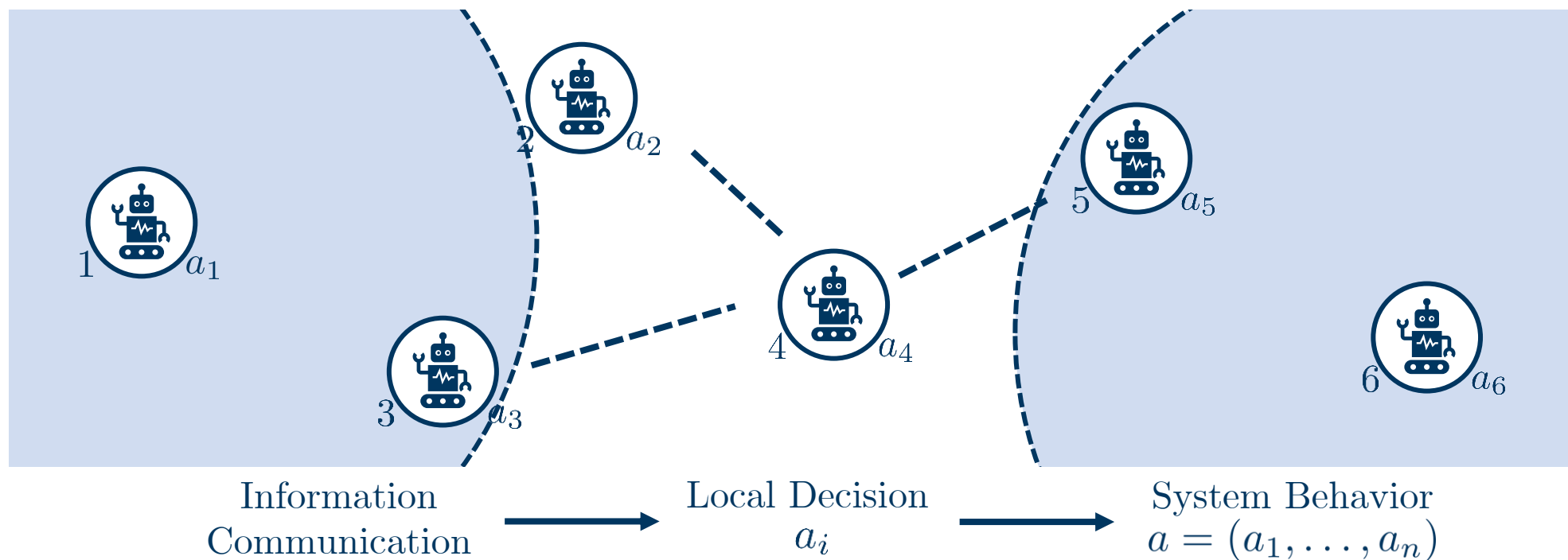
# Communication Design in Engineered System



How should agents communicate?

- Networked communication?
- Spatially Local communication?

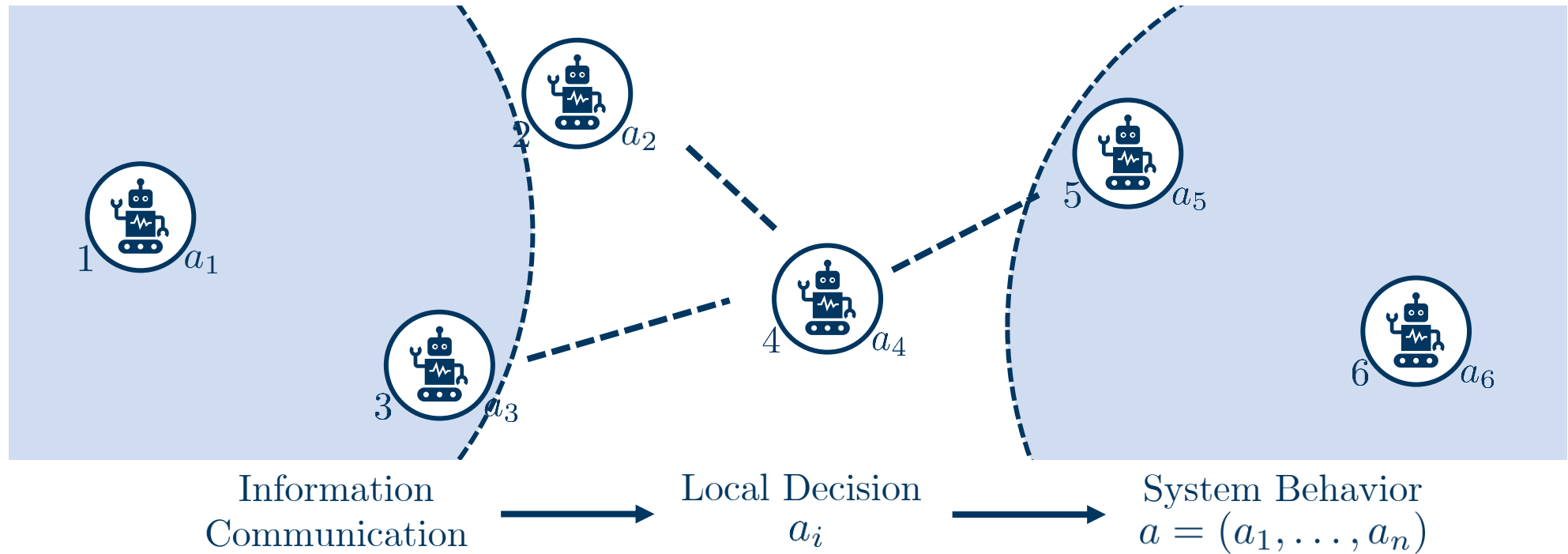
# Communication Design in Engineered System



How should agents communicate?

- Networked communication?
- Spatially Local communication?
- Pairwise communication?
- Etc.

# Communication Design in Engineered System



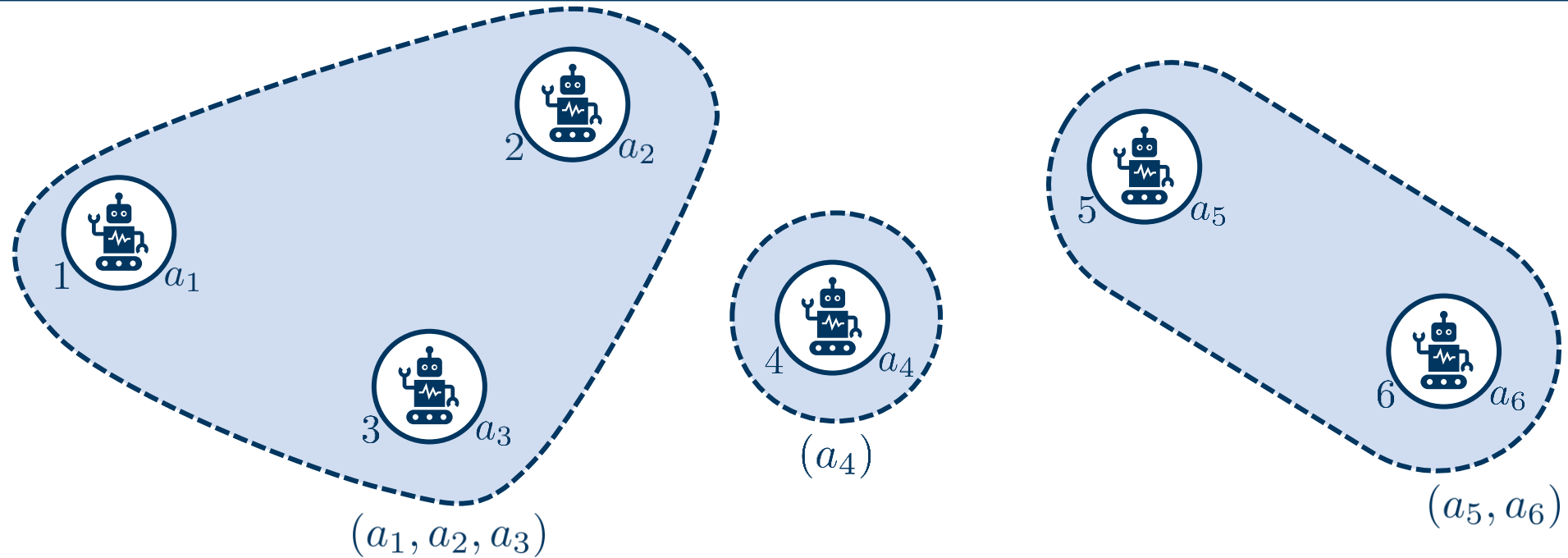
How should agents communicate?

- Networked communication?
- Spatially Local communication?
- Pairwise communication?
- Etc.

How does communication affect decision-making?

# Collaborative Communication & Decision-Making

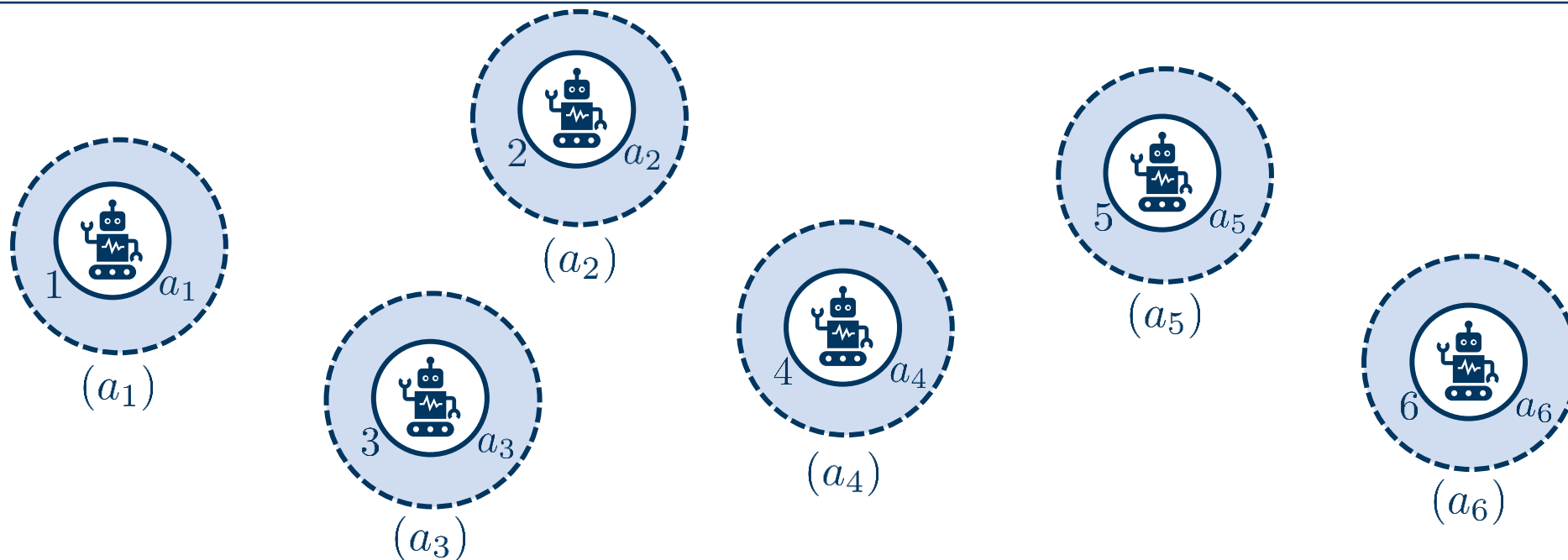
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Collaborative Multi-Agent System: Defined groups of agents can communicate and collaborate in making their decision.

# Collaborative Communication & Decision-Making

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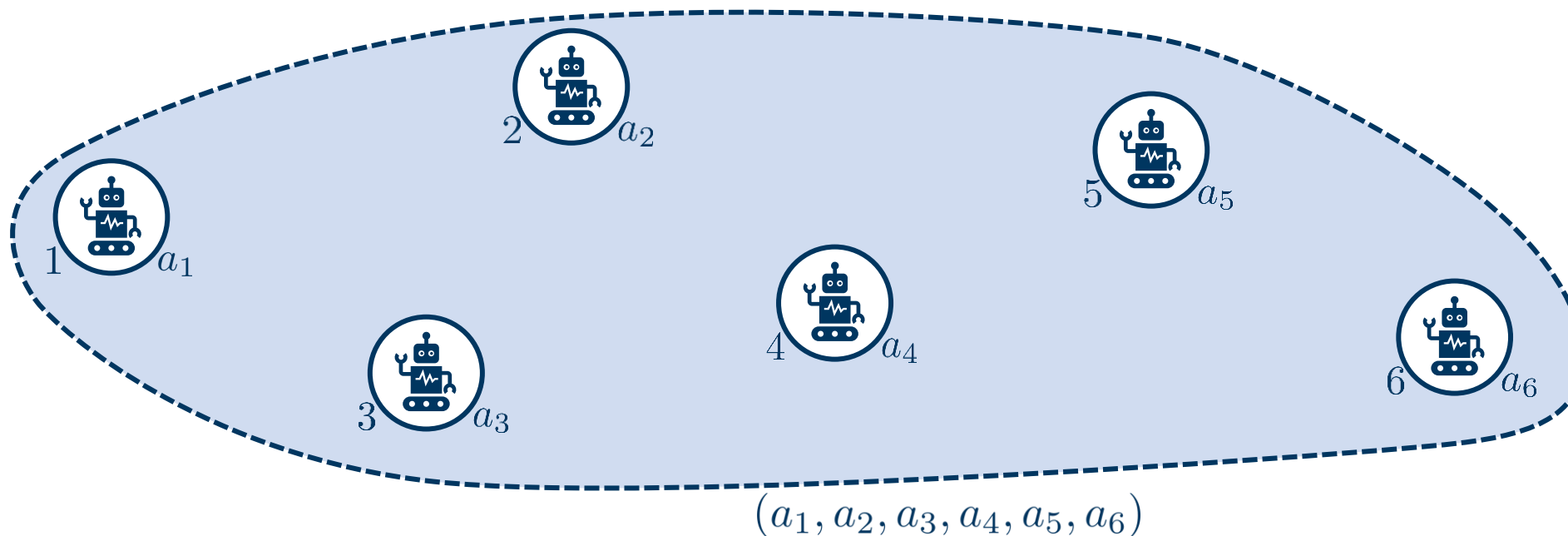


Collaborative Multi-Agent System: Defined groups of agents can communicate and collaborate in making their decision.

Distributed Decision-making  
(Bad performance / low complexity)

# Collaborative Communication & Decision-Making

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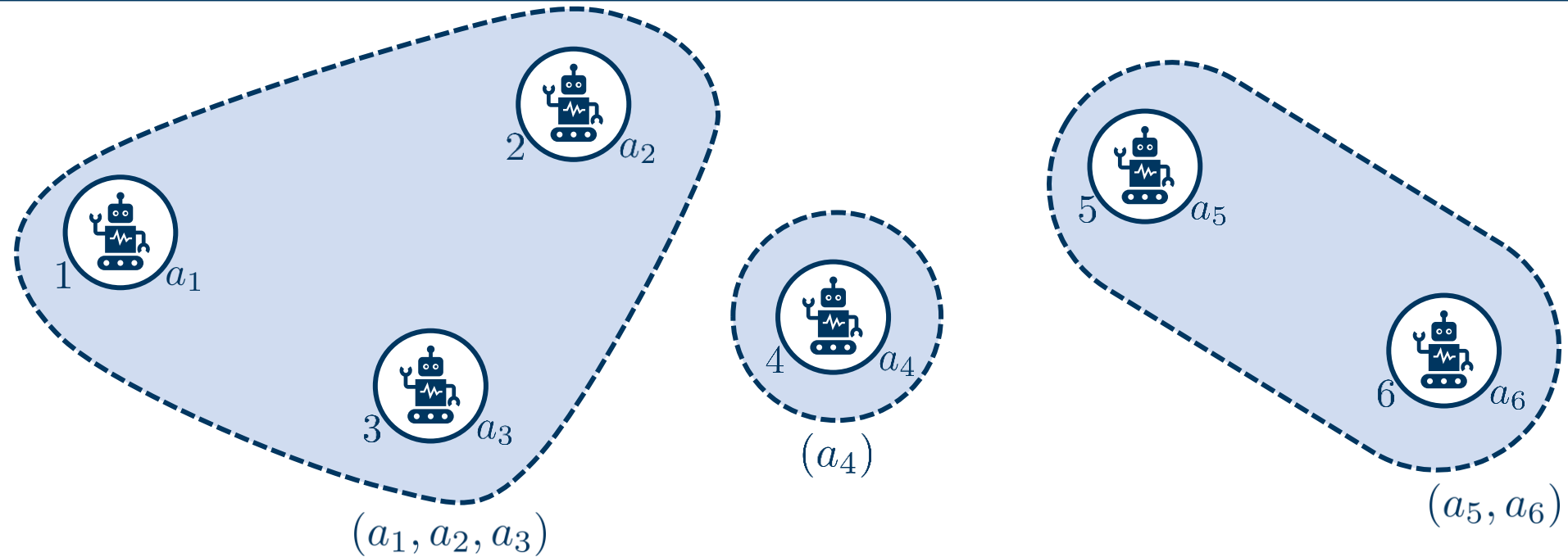


Collaborative Multi-Agent System: Defined groups of agents can communicate and collaborate in making their decision.

Distributed Decision-making  
(Bad performance / low complexity)

Centralized Decision-making  
(Best performance / high complexity)

# Collaborative Communication & Decision-Making



Collaborative Multi-Agent System: Defined groups of agents can communicate and collaborate in making their decision.

Distributed Decision-making  
(Bad performance / low complexity)

*Partially Collaborative*  
Decision-making

Centralized Decision-making  
(Best performance / high complexity)

Study the *benefits* and *costs* of collaborative multi-agent systems

# Distributed Decision-Making in Engineered Systems

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## Game Theoretic approach to Multi-Agent Control



# Distributed Decision-Making in Engineered Systems

---

## Game Theoretic approach to Multi-Agent Control

System Objective:

$$\max_{a_1, \dots, a_n} W(a_1, \dots, a_n)$$

# Distributed Decision-Making in Engineered Systems

---

## Game Theoretic approach to Multi-Agent Control

System Objective:

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Local

Decision-Making:

# Distributed Decision-Making in Engineered Systems

## Game Theoretic approach to Multi-Agent Control

System Objective:

$$\max_{a_1, \dots, a_n} W(a_1, \dots, a_n)$$



Agent 1



Agent 2



Agent 3

...



Agent  $n$

Local

Decision-Making:

$$\max_{a_1} W(a_1, a_{-1})$$

$$\max_{a_2} W(a_2, a_{-2})$$

$$\max_{a_3} W(a_3, a_{-3})$$

$$\max_{a_n} W(a_n, a_{-n})$$

# Distributed Decision-Making in Engineered Systems

## Game Theoretic approach to Multi-Agent Control

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Dynamics:

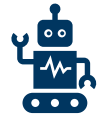
$$a(t+1) = \text{update}(a(t))$$

# Distributed Decision-Making in Engineered Systems

## Game Theoretic approach to Multi-Agent Control

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Dynamics:

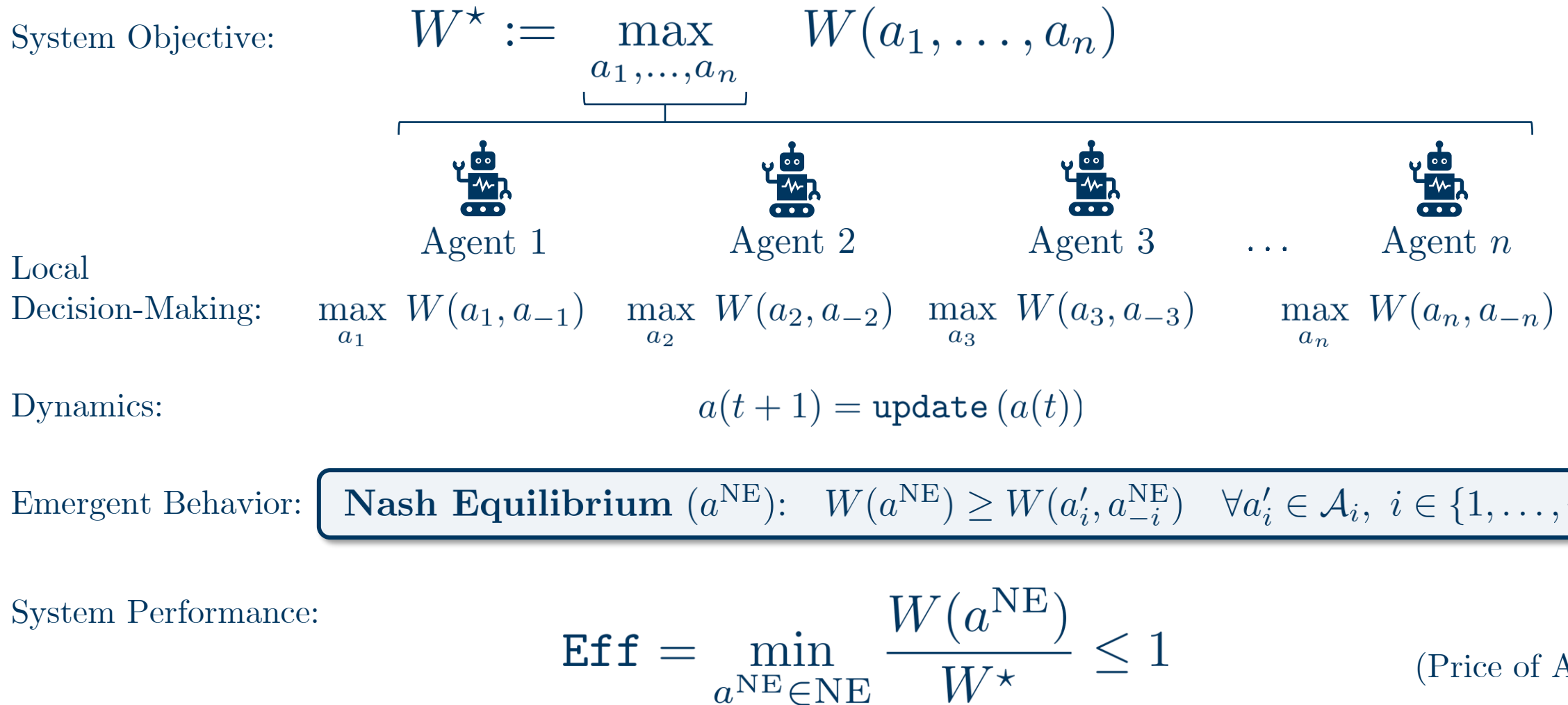
$$a(t+1) = \text{update}(a(t))$$

Emergent Behavior:

$$\text{Nash Equilibrium } (a^{\text{NE}}): \quad W(a^{\text{NE}}) \geq W(a'_i, a_{-i}^{\text{NE}}) \quad \forall a'_i \in \mathcal{A}_i, i \in \{1, \dots, n\}$$

# Distributed Decision-Making in Engineered Systems

## Game Theoretic approach to Multi-Agent Control



# Focus of Existing Research

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## Dynamics

### Reaching Nash Equilibria

#### Convergence Rate

- *Uncoupled dynamics do not lead to Nash equilibrium.* American Economic Review. Hart S, Mas-Colell A. 2003
- *The complexity of computing a Nash equilibrium.* SIAM Journal on Computing, Daskalakis C, Goldberg PW, Papadimitriou CH. 2009.

#### Asynchrony

- *Nash equilibrium seeking in noncooperative games.* IEEE Transactions on Automatic Control. Frihauf P, Krstic M, Basar T. 2011.
- *Learning efficient Nash equilibria in distributed systems.* Games and Economic behavior. Pradelski BS, Young HP. 2012.

#### Noise and Perturbations

- *Learning with Bandit Feedback in Potential Games.* Advances in Neural Information Processing Systems. Heliou A, Cohen J, Mertikopoulos P. 2017.
- *Learning in Games: Robustness of Fast Convergence.* Advances in Neural Information Processing Systems. Foster D, Li Z, Lykouris T, Sridharan K, Tardos E. 2016.

## Performance

### Efficiency of Nash Equilibria

#### Quantify the Price of Anarchy

- *Worst-case equilibria.* Computer science review. Koutsoupias E, Papadimitriou C. 2009.
- *Nash equilibria in competitive societies, with applications to facility location, traffic routing and auctions.* Symposium on Foundations of CS. Vetta A. 2002
- *The price of anarchy of finite congestion games.* In Proceedings of the thirty-seventh annual ACM symposium on Theory of computing. Christodoulou G, Koutsoupias E. 2005.
- *Intrinsic robustness of the price of anarchy.* Journal of the ACM (JACM). Roughgarden T. 2015.

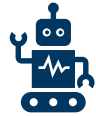
#### Optimize the Price of Anarchy

- *Covering games: Approximation through non-cooperation.* In International Workshop on Internet and Network Economics. Gairing M. 2009.
- *A unifying tool for bounding the quality of non-cooperative solutions in weighted congestion games.* Theory of Computing Systems. Bilò V. 2018.
- *Utility design for distributed resource allocation—part I: Characterizing and optimizing the exact price of anarchy.* IEEE Transactions on Automatic Control. Paccagnan D, Chandan R, Marden JR. 2019.

Challenge existing *informational* and *communication* assumptions  
of Nash equilibria and associated efficiency guarantees

# Collaborative Communication

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Agent 1

$$\max_{a_1} W(a_1, a_{-1})$$



Agent 2

$$\max_{a_2} W(a_2, a_{-2})$$



Agent 3

$$\max_{a_3} W(a_3, a_{-3})$$

...



Agent  $n$

$$\max_{a_n} W(a_n, a_{-n})$$

**No communication:** unilateral deviations

**Nash Equilibrium** ( $a^{\text{NE}}$ ):

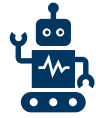
$$W(a^{\text{NE}}) \geq W(a'_i, a_{-i}^{\text{NE}})$$

$$\forall a'_i \in \mathcal{A}_i,$$

$$i \in N = \{1, \dots, n\}$$



# Collaborative Communication



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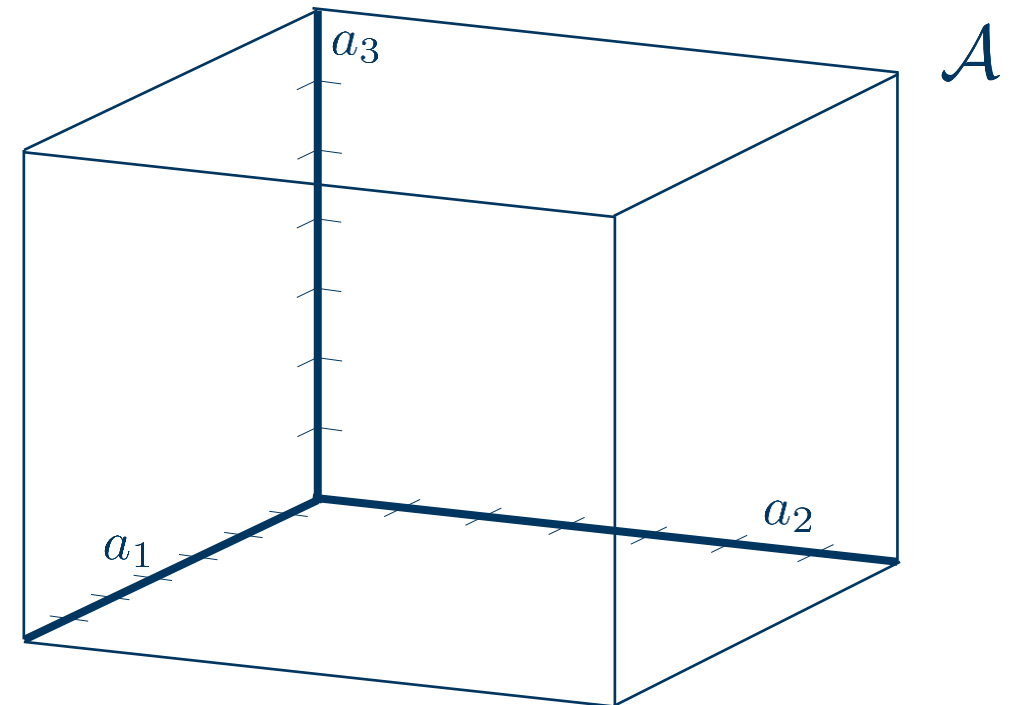
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**No communication:** unilateral deviations

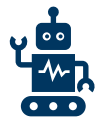
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$$\begin{aligned} W(a^{\text{NE}}) &\geq W(a'_i, a_{-i}^{\text{NE}}) \\ &\forall a'_i \in \mathcal{A}_i, \\ &i \in N = \{1, \dots, n\} \end{aligned}$$

E.g.,



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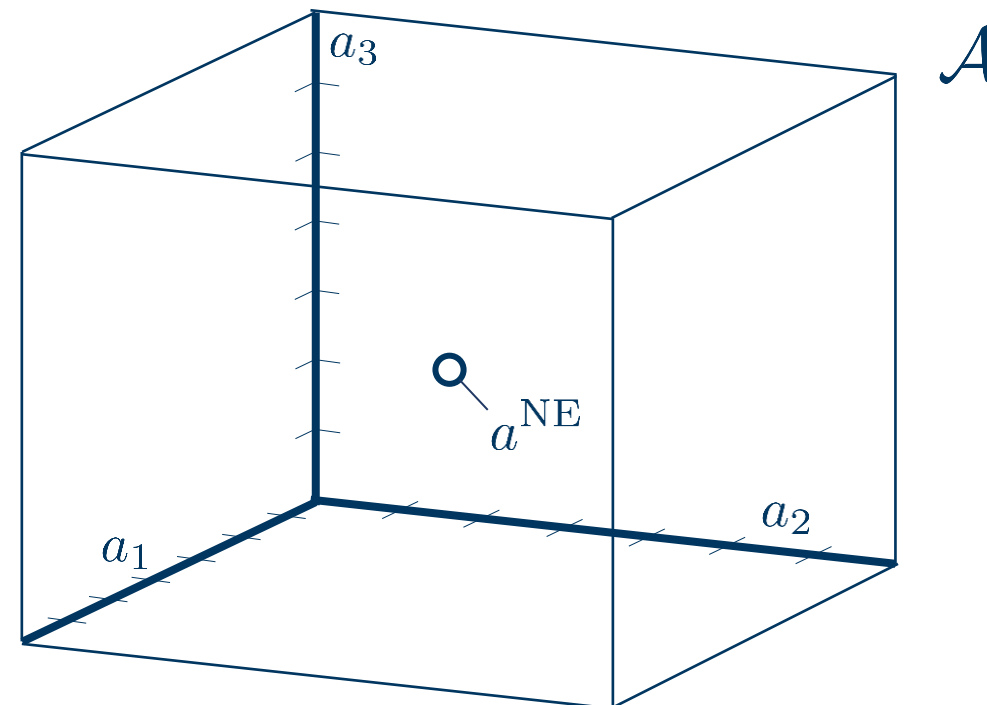
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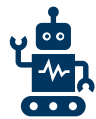
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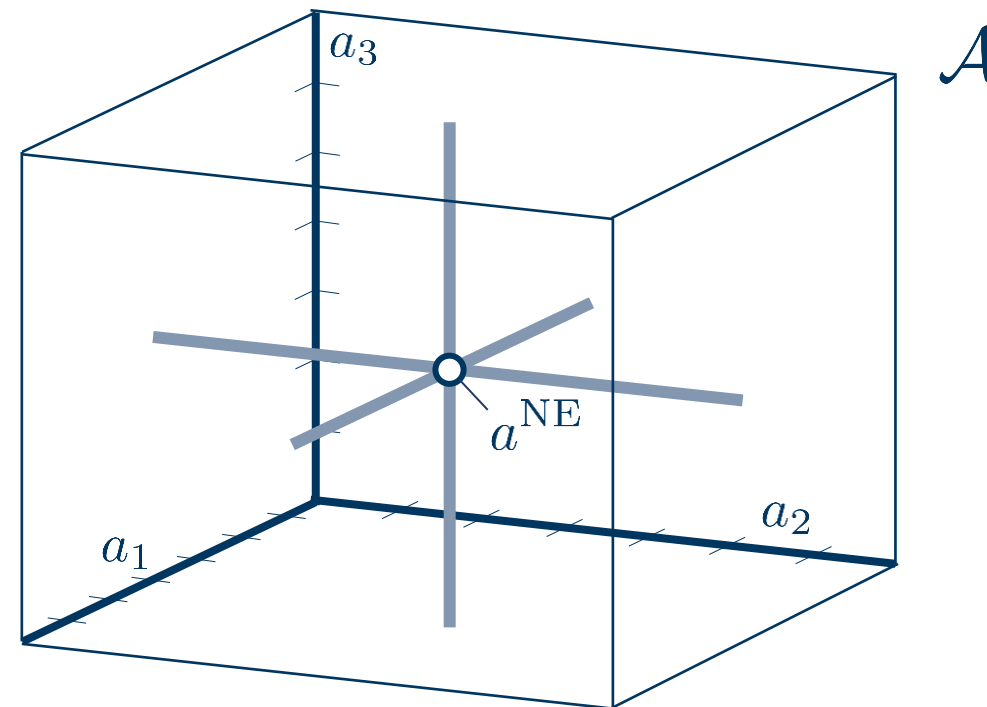
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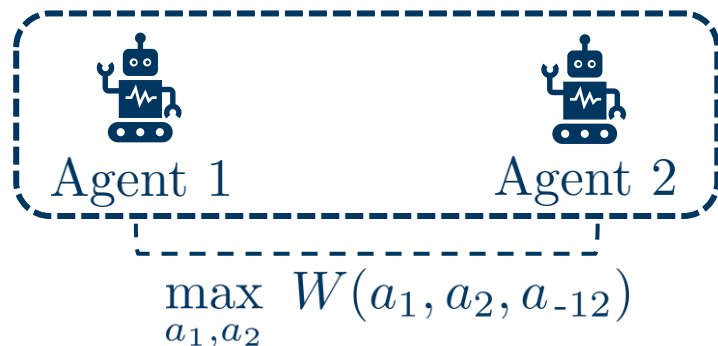
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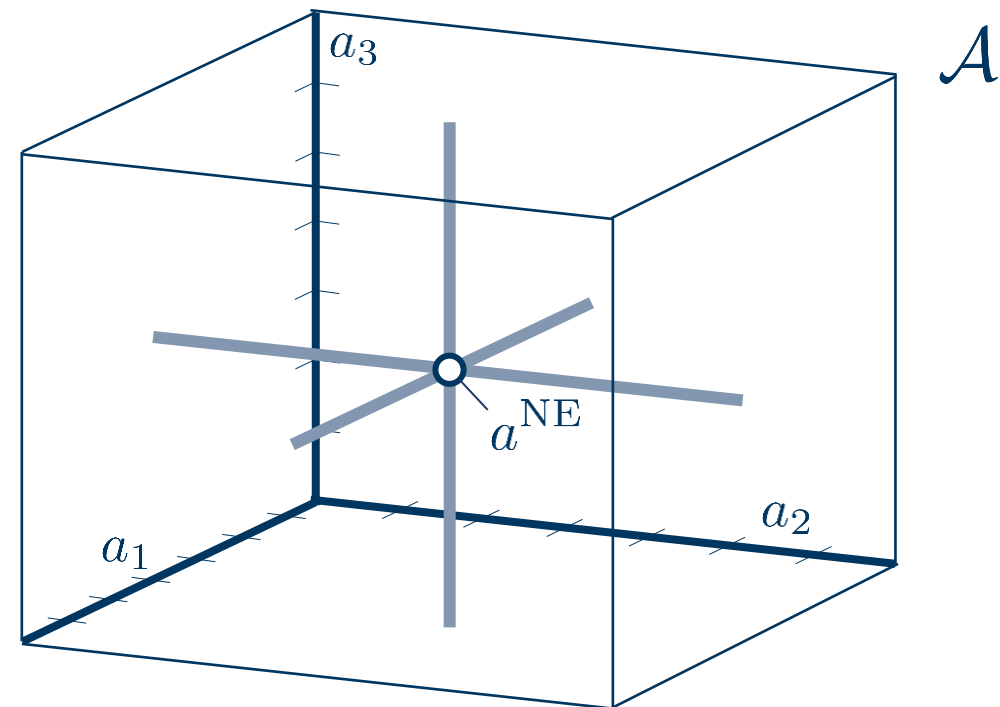
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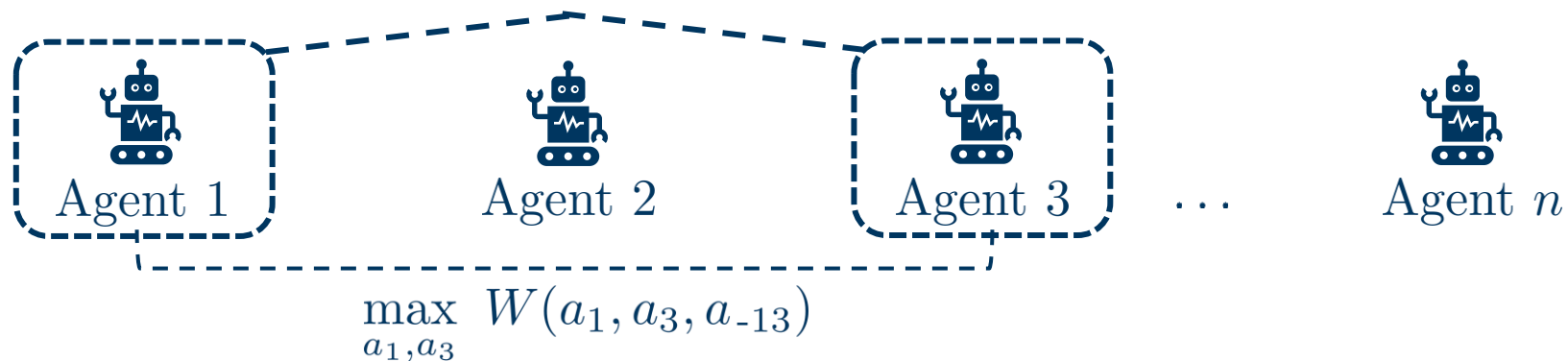
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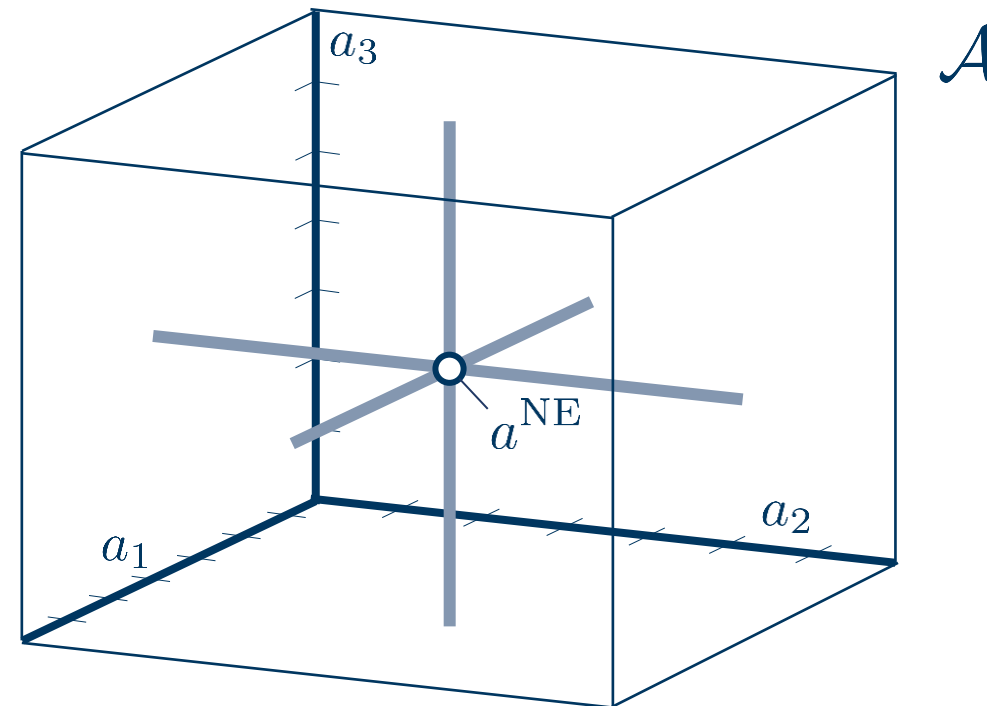
# Collaborative Communication



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# Collaborative Communication

  
Agent 1

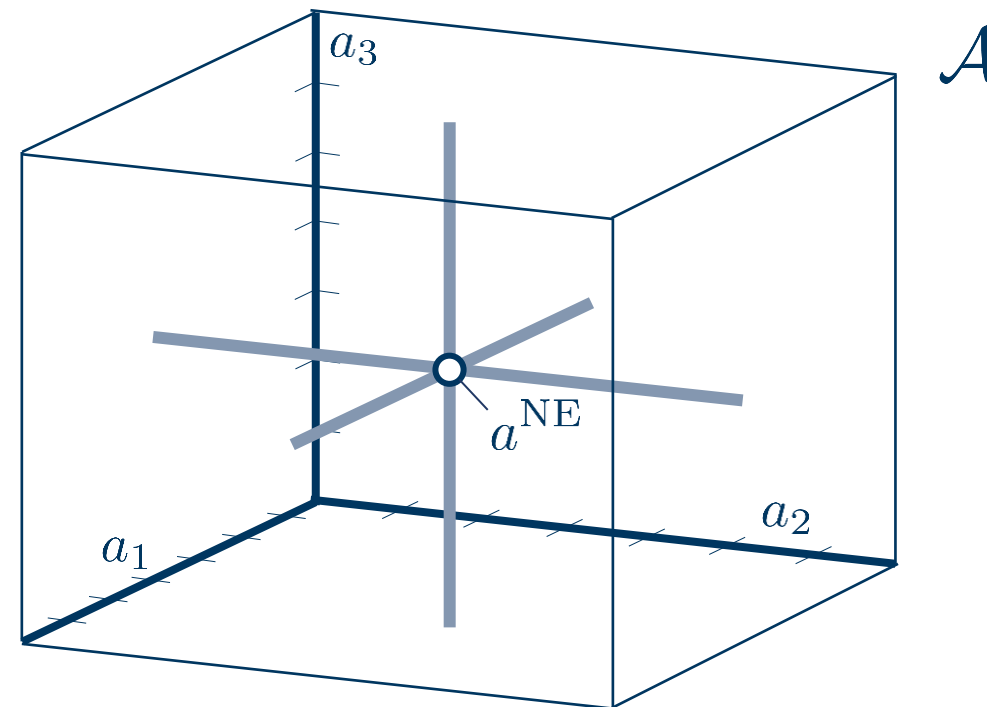


$$\max_{a_2, a_n} W(a_2, a_n, a_{-2n})$$

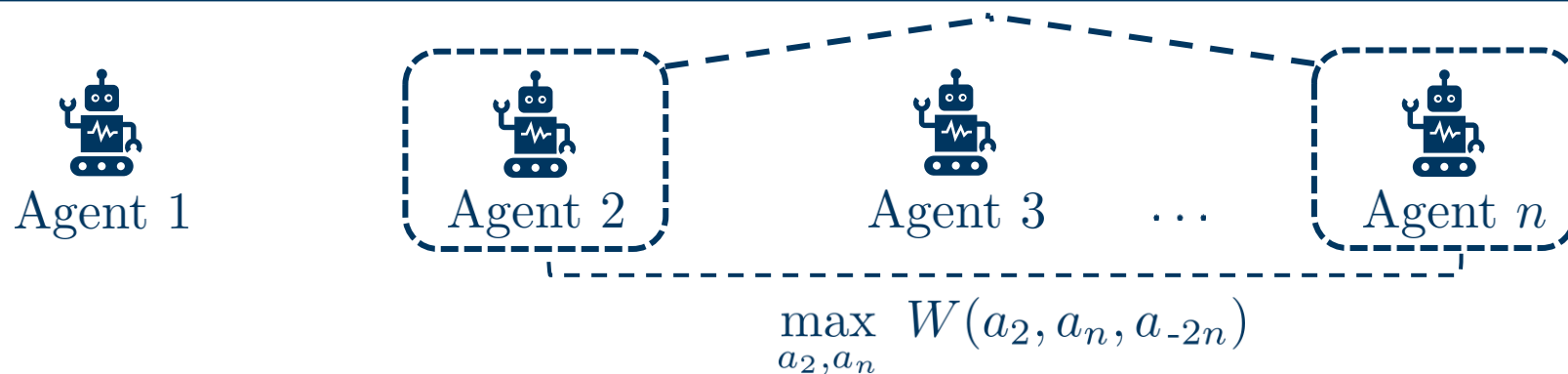
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E.g.,



# Collaborative Communication



Pair-wise communication: bilateral deviations

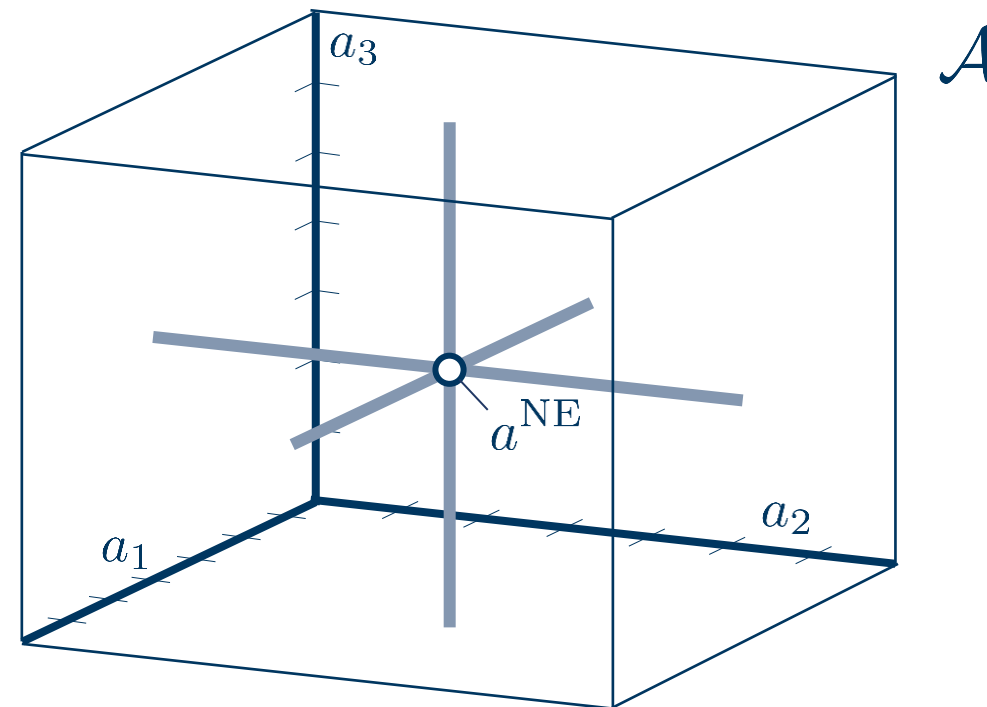
**2-Strong Nash Equilibrium** ( $a^{2\text{SNE}}$ ):

$$W(a^{2\text{SNE}}) \geq W(a'_i, a'_j, a_{-ij}^{2\text{SNE}})$$

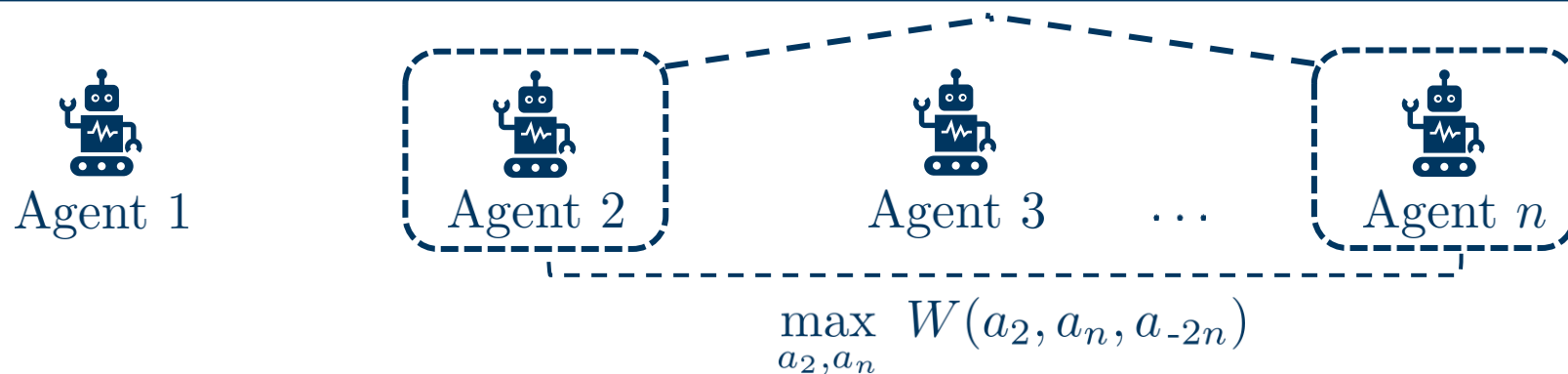
$$\forall a'_i \in \mathcal{A}_i, a'_j \in \mathcal{A}_j,$$

$$i, j \in N$$

E.g.,



# Collaborative Communication



Pair-wise communication: bilateral deviations

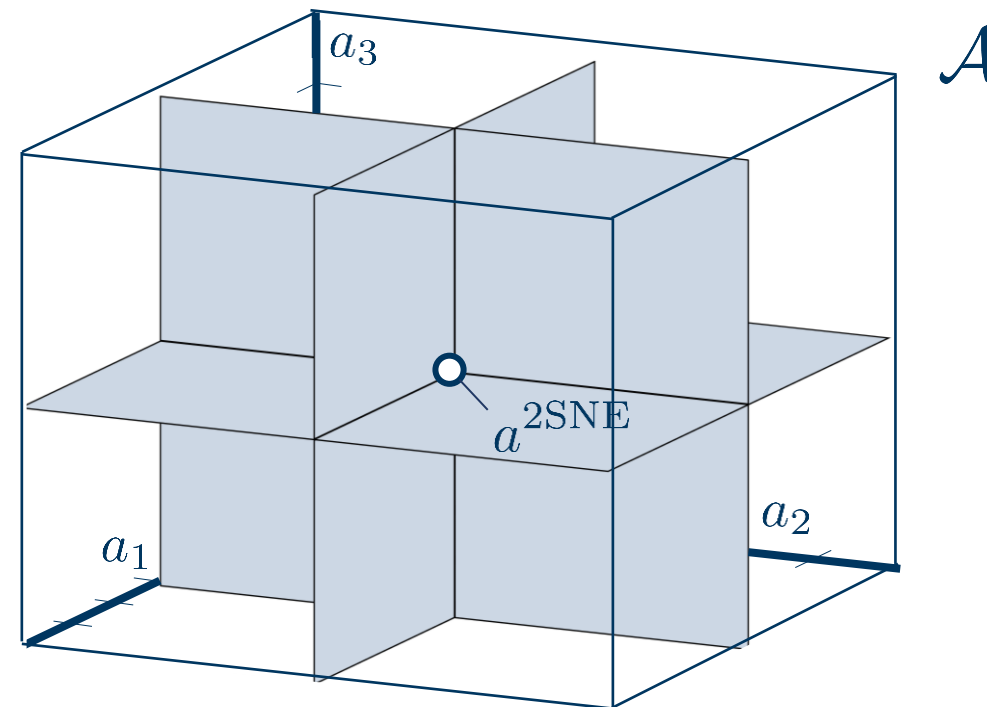
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E.g.,





# Collaborative Communication



$$\max_{a_{\Gamma}} W(a_{\Gamma}, a_{-\Gamma}), \quad \text{where } \Gamma \subseteq N$$

**Multi-agent communication:** group deviations

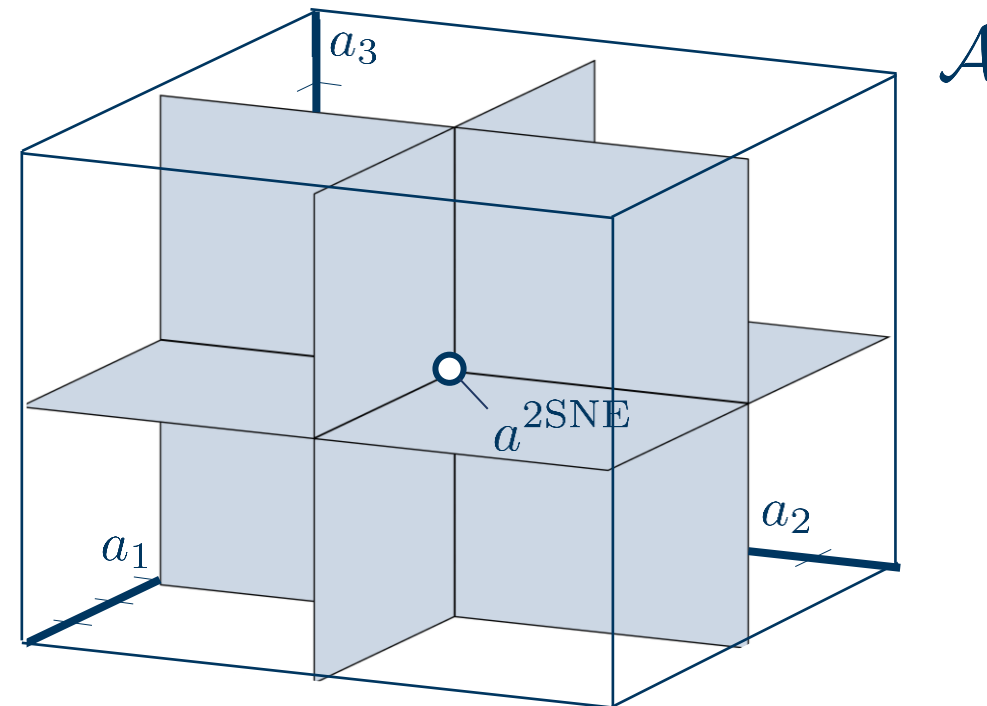
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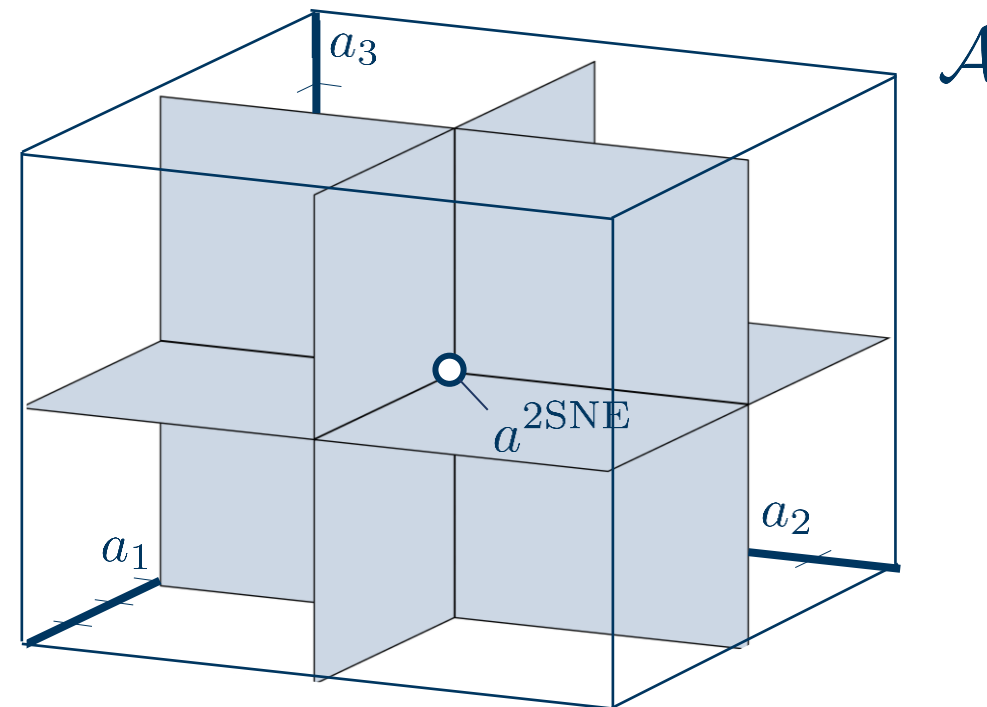
**$k$ -Strong Nash Equilibrium ( $a^{k\text{SNE}}$ ):**

$$W(a^{k\text{SNE}}) \geq W(a'_\Gamma, a_{-\Gamma}^{k\text{SNE}})$$

$$\forall a'_\Gamma \in \mathcal{A}_\Gamma = \prod_{i \in \Gamma} \mathcal{A}_i,$$

$$\Gamma \in \{Z \subseteq N : |Z| \leq k\}$$

E.g.,



# Collaborative Communication



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**Multi-agent communication:** group deviations

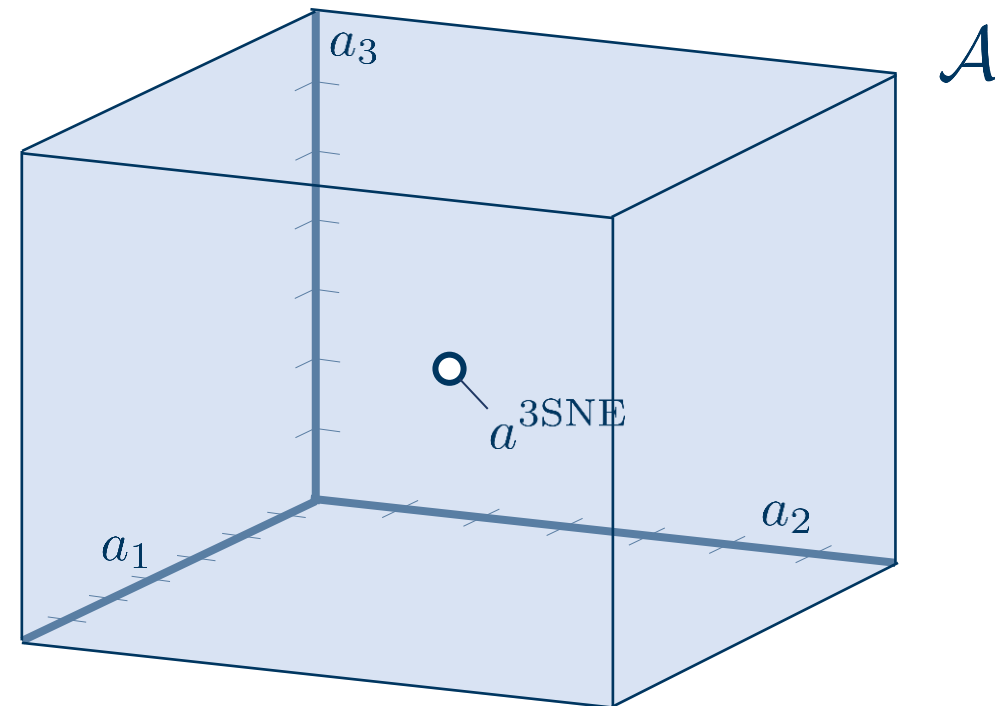
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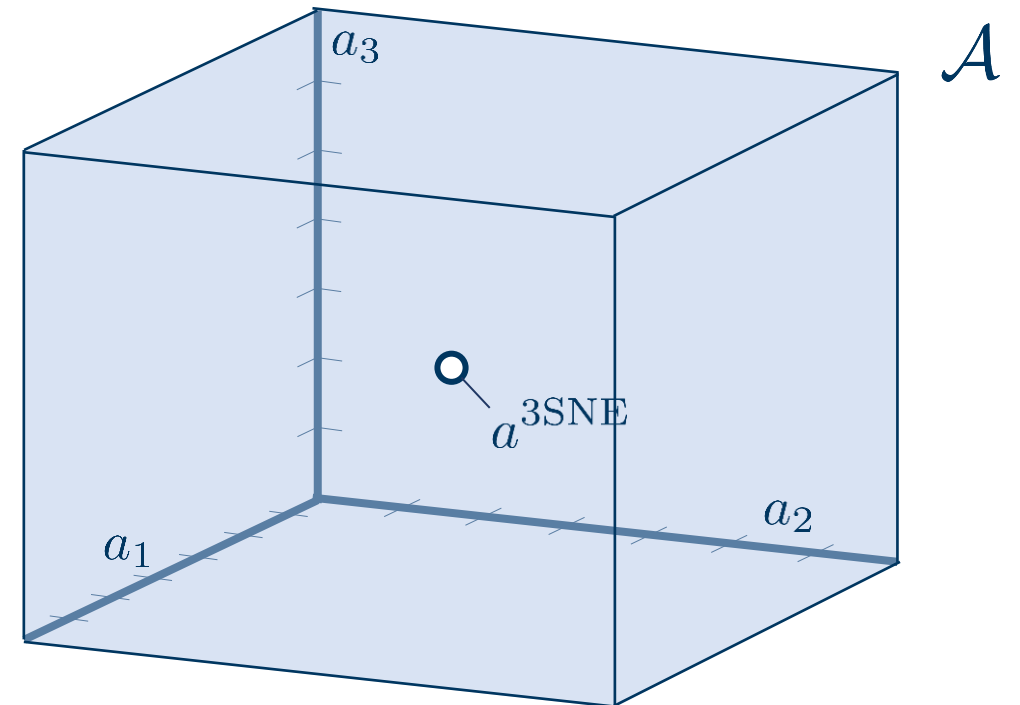
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E.g.,



Use communication to *improve efficiency* and bridge gap between *centralized* and *decentralized* performance

# $k$ – Strong Nash Equilibria

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## In General:

- Typically discussed in cooperative/cost-sharing games
- *Need not exist* (in general games)
- No guarantee of efficiency improvement (when such an equilibrium exists)

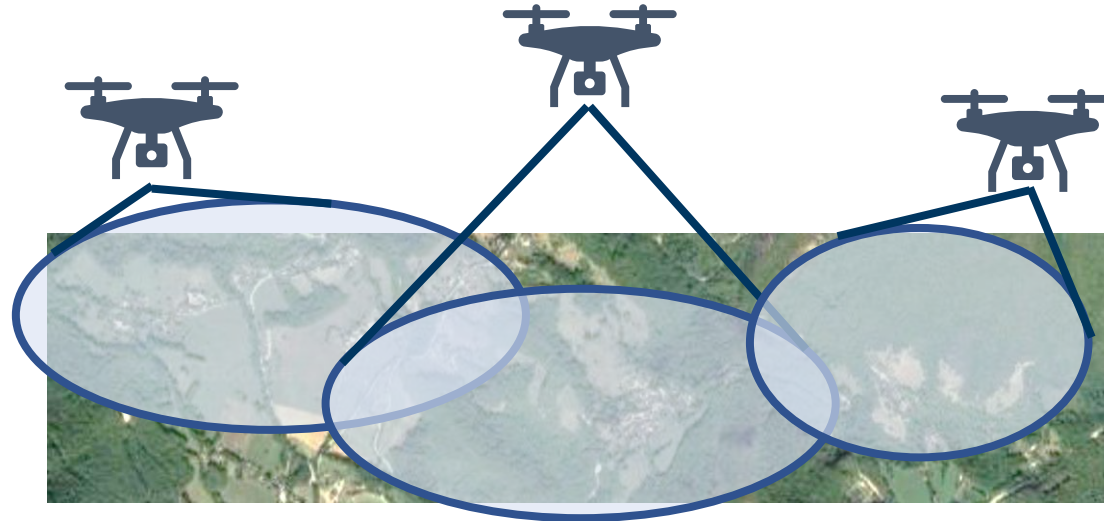
## In Our Work:

- Group of players collaborate to improve global (**common interest**) objective
- Existence guaranteed
- Optimal solution is a  $k$ -SNE
- Finite convergence time

Focus: 1. How does *efficiency* improve with communication ( $k$ )?  
2. What additional *complexity* is incurred?

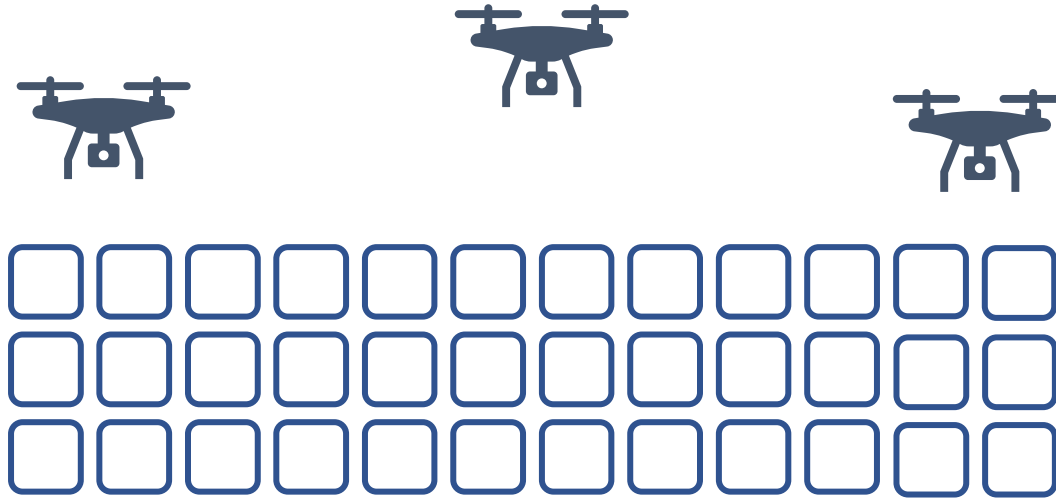
# Resource Allocation Problems

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# Resource Allocation Problems

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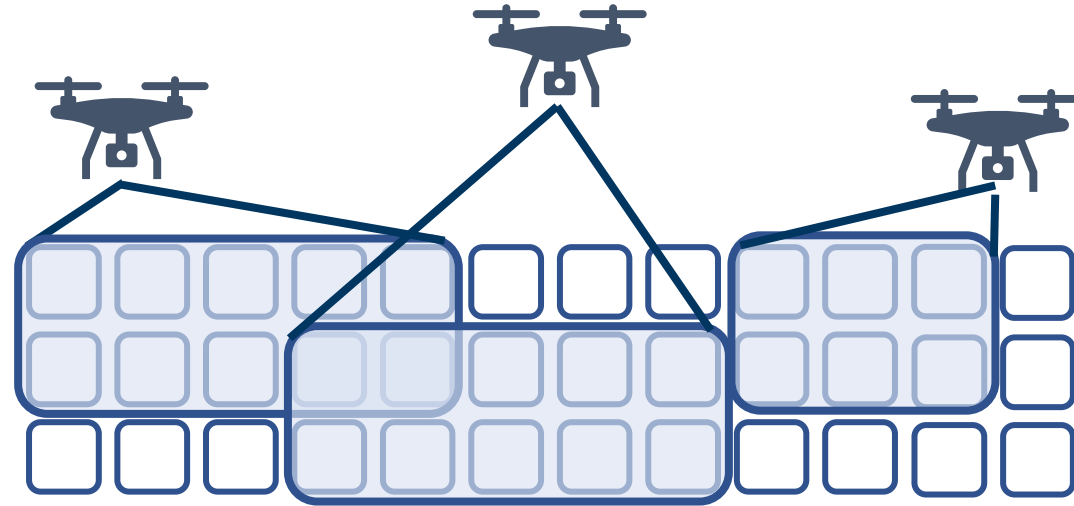


Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

Agents:  $i \in \mathcal{N} = \{1, \dots, n\}$

# Resource Allocation Problems

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Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

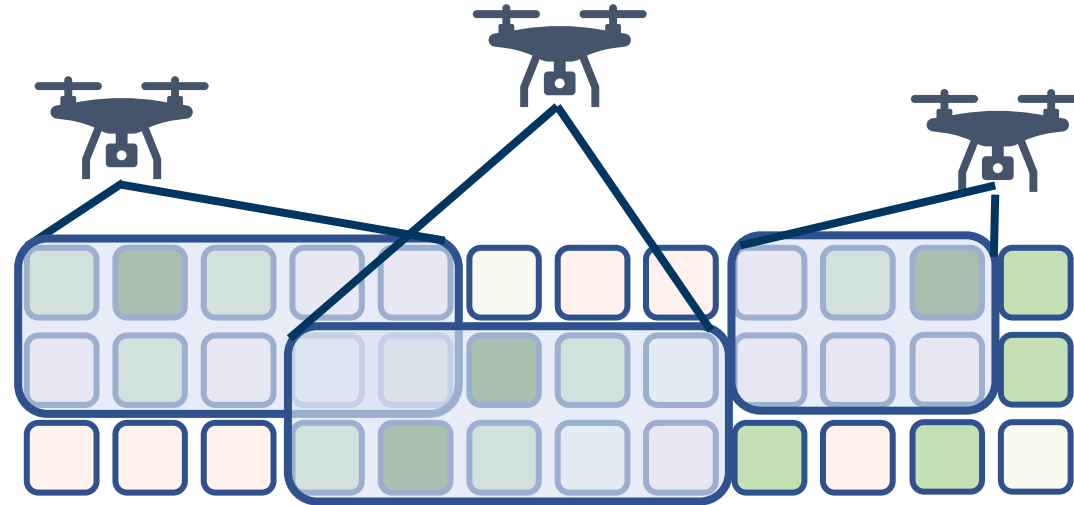
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Actions:  $a_i \subset \mathcal{R}, \mathcal{A}_i \subseteq 2^{\mathcal{R}}$



# Resource Allocation Problems

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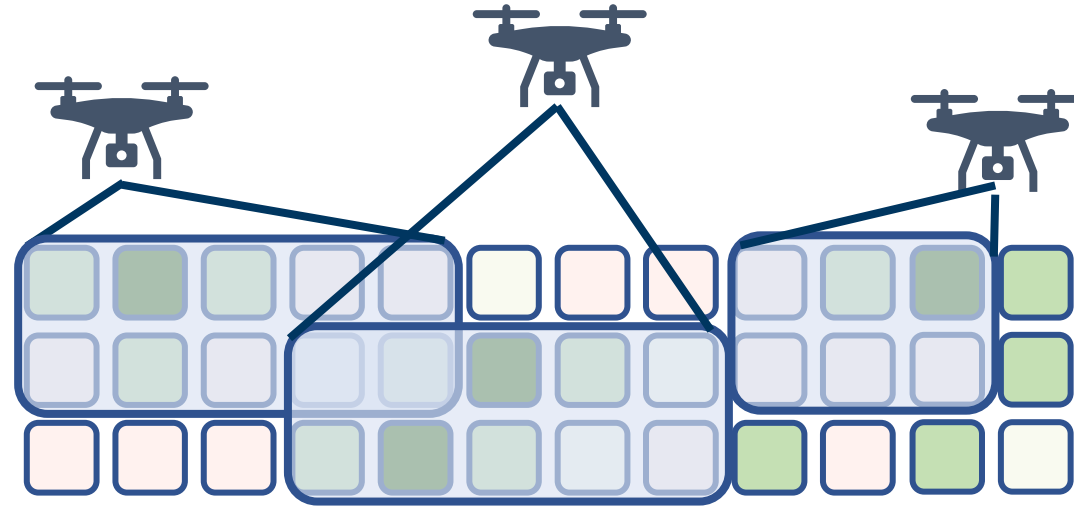
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Actions:  $a_i \subset \mathcal{R}, \mathcal{A}_i \subseteq 2^{\mathcal{R}}$

System Welfare  $W(a) = \sum_{r \in \mathcal{R}} v_r w(|a|_r)$

# Resource Allocation Problems



Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

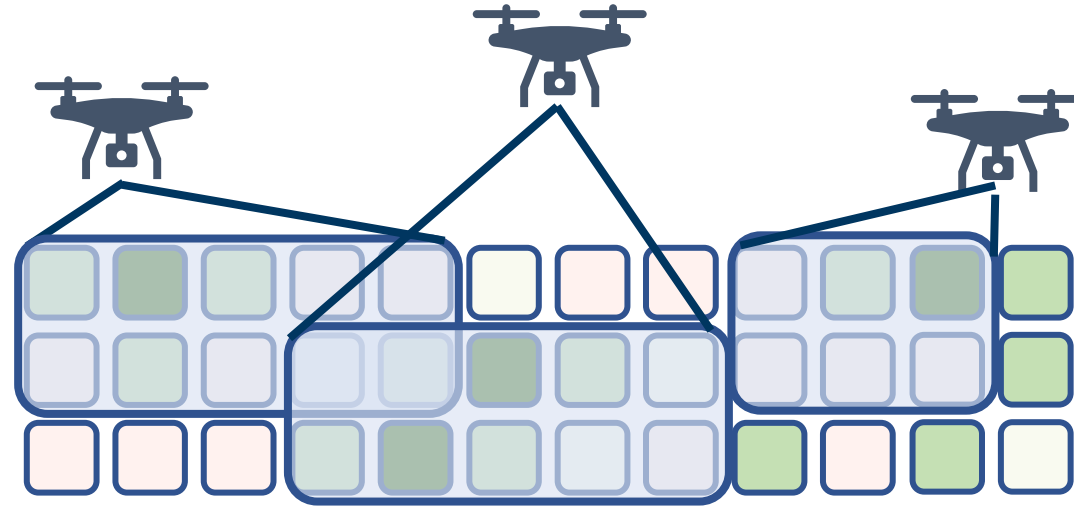
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# Resource Allocation Problems



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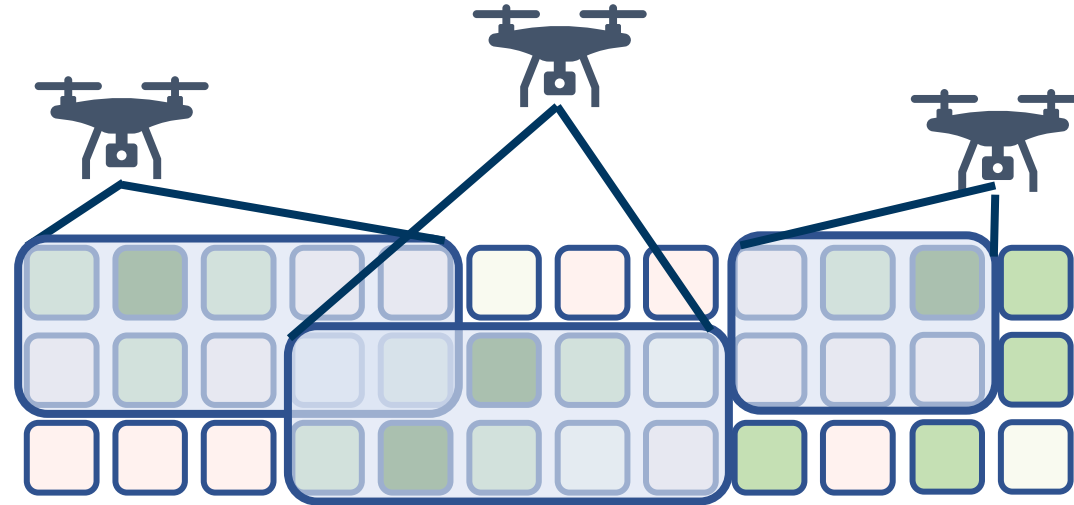
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Optimal Solution:  $a^{\text{opt}} \in \arg \max_{a \in \mathcal{A}_1 \times \dots \times \mathcal{A}_n} W(a)$



# Resource Allocation Problems



Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

Agents:  $i \in N = \{1, \dots, n\}$

Actions:  $a_i \subset \mathcal{R}, \mathcal{A}_i \subseteq 2^{\mathcal{R}}$

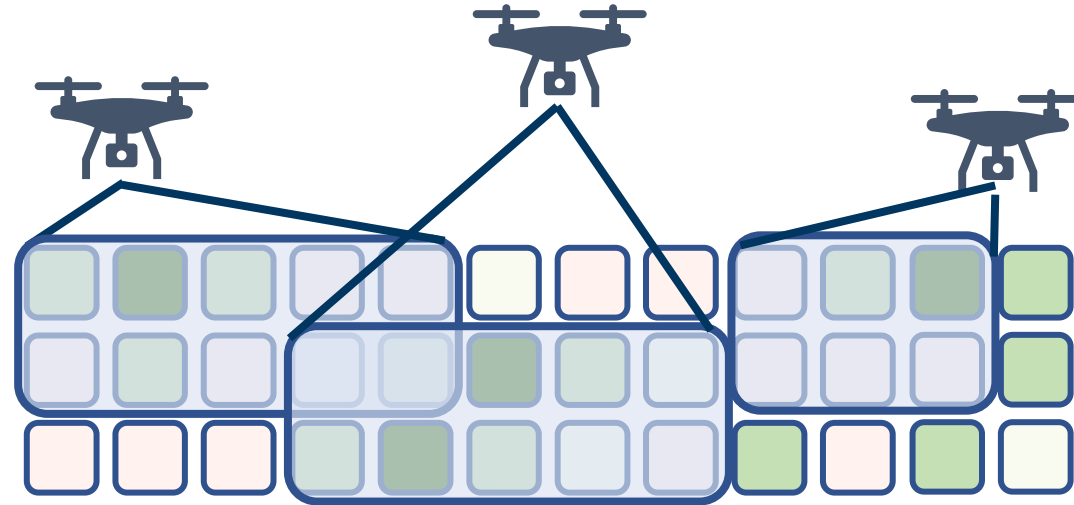
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**Combinatorial and NP-hard!**



# Resource Allocation Problems



Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

Utility:  $U_i(a_i, a_{-i}) = W(a_i, a_{-i})$

Agents:  $i \in N = \{1, \dots, n\}$

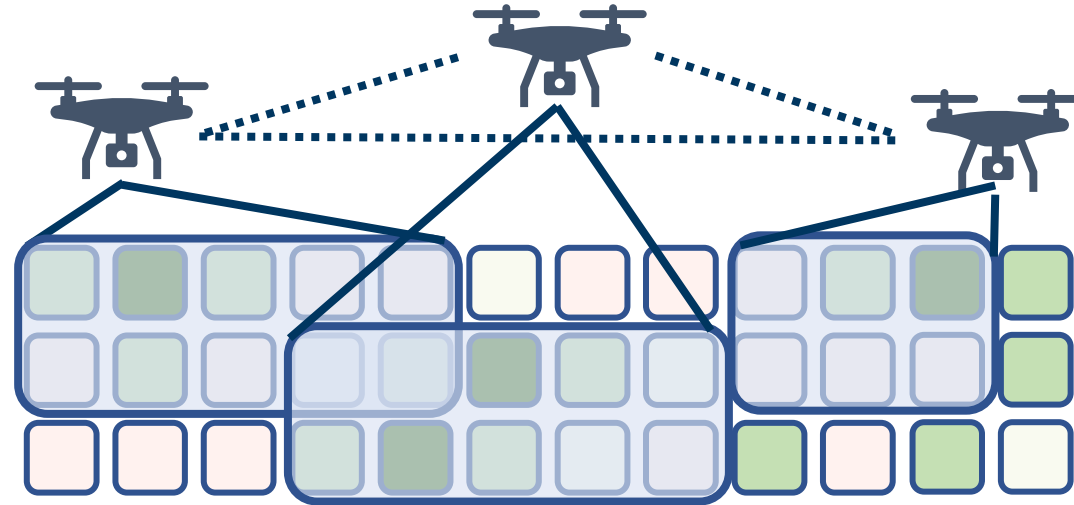
Actions:  $a_i \subset \mathcal{R}, \mathcal{A}_i \subseteq 2^{\mathcal{R}}$

System Welfare  $W(a) = \sum_{r \in \mathcal{R}} \underbrace{v_r w(|a|_r)}_{\text{Value of having } |a|_r \text{ agents share resource } r}$

Optimal Solution:  $a^{\text{opt}} \in \arg \max_{a \in \mathcal{A}_1 \times \dots \times \mathcal{A}_n} W(a)$

**Combinatorial and NP-hard!**

# Resource Allocation Problems



Resources:  $r \in \mathcal{R} = \{1, \dots, R\}$

Agents:  $i \in N = \{1, \dots, n\}$

Actions:  $a_i \subset \mathcal{R}, \mathcal{A}_i \subseteq 2^{\mathcal{R}}$

System Welfare  $W(a) = \sum_{r \in \mathcal{R}} \underbrace{v_r w(|a|_r)}_{\text{Value of having } |a|_r \text{ agents share resource } r}$

Optimal Solution:  $a^{\text{opt}} \in \arg \max_{a \in \mathcal{A}_1 \times \dots \times \mathcal{A}_n} W(a)$

**Combinatorial and NP-hard!**

Utility:  $U_{\Gamma}(a_{\Gamma}, a_{-\Gamma}) = W(a_{\Gamma}, a_{-\Gamma}), \Gamma \subseteq N$

Emergent System Behavior:

**$k$  – Strong  
Nash Equilibrium**

Efficiency:

$$\text{Eff}(k) = \min_{a^{k\text{SNE}} \in k\text{SNE}} \frac{W(a^{k\text{SNE}})}{W(a^{\text{opt}})}$$

Efficiency as a function of the level of communication between agents

# Efficiency of $k$ – Strong Nash Equilibria

How much does inter-agent communication *improve efficiency*?

Theorem 1.1:

[BLF, Paccagnan, Pradeslki, Marden CDC23\*]

For a resource allocation problem  $(\mathcal{R}, N, \mathcal{A}, \{v_r\}_{r \in \mathcal{R}}, w)$ , a  $k$ SNE approximates the optimal solution with

$$\text{Eff}(k) \geq P^*(n, w, k),$$

where  $P^*(n, w, k)$  is the solution to a linear program with  $k+1$  decision variables and  $\mathcal{O}(kn^3)$  constraints. Further, this bound is tight.

$$\begin{aligned}
 P^*(n, w, k) = & \max_{\theta \in \mathbb{R}_{\geq 0}^{|Z|}} \sum_{e, x, o} w(o + x) \theta(e, x, o) \\
 \text{s.t.} & \sum_{e, x, o} \left( \frac{n!}{(n - \zeta)!} w(e + x) - \sum_{\substack{0 \leq \alpha \leq \zeta \\ 0 \leq \beta \leq \alpha}} \binom{\zeta}{\alpha} \binom{\zeta - \alpha}{\beta} e^{\alpha} o^{\beta} (n - e - o)^{\zeta - \alpha - \beta} w(e + x + \beta - \alpha) \right) \theta(e, x, o) \geq 0 \\
 & \forall \zeta \in \{1, \dots, k\} \\
 & \sum_{e, x, o} w(e + x) \theta(e, x, o) = 1
 \end{aligned}$$

# Efficiency of $k$ – Strong Nash Equilibria

How much does inter-agent communication *improve efficiency*?

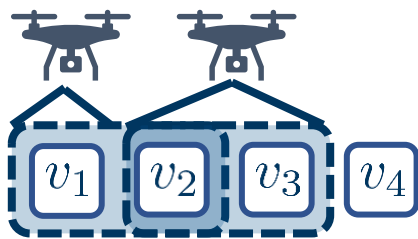
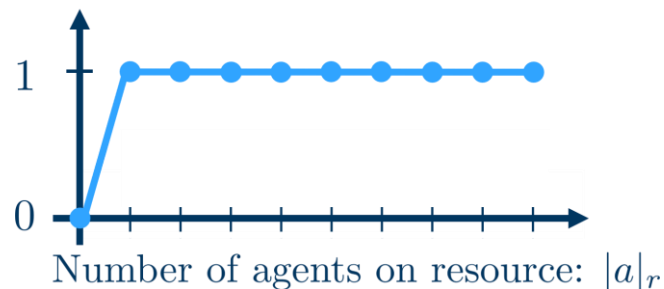
Theorem 1.1:

[BLF, Paccagnan, Pradeski, Marden CDC23\*]

$$\text{Eff}(k) \geq P^*(n, w, k)$$

Example: Covering Problems

$$w(|a|_r) = \mathbb{1} [ |a|_r \geq 1 ]$$



$$W(a) = v_1 + v_2 + v_3$$



# Efficiency of $k$ – Strong Nash Equilibria

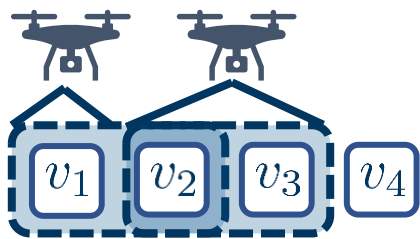
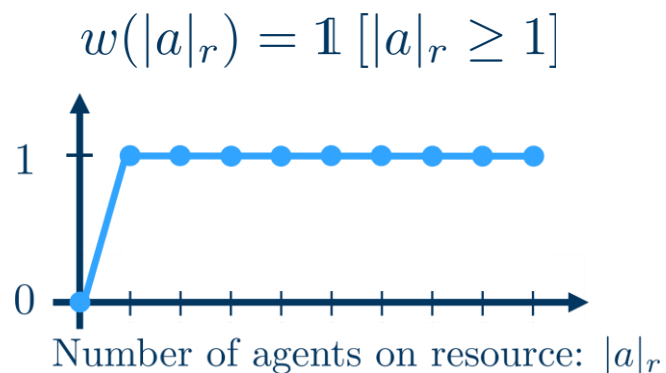
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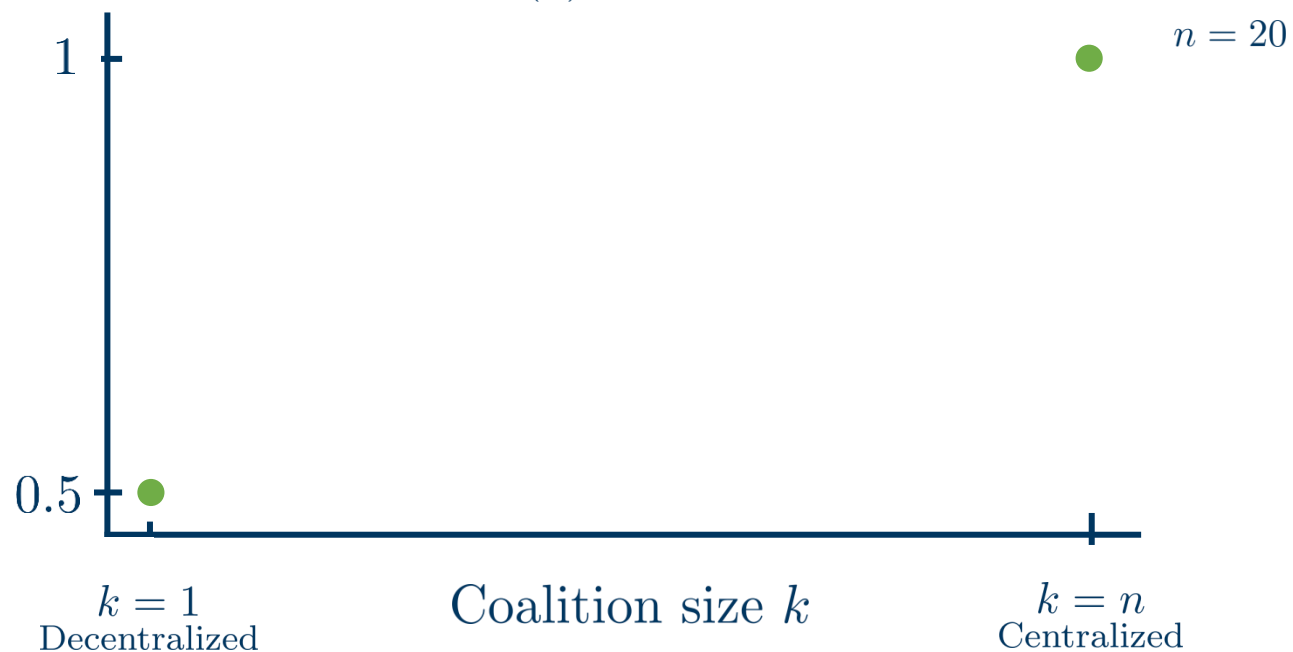
$$\text{Eff}(k) \geq P^*(n, w, k)$$

Example: Covering Problems



$$W(a) = v_1 + v_2 + v_3$$

$\text{Eff}(k)$  Lower-bound



# Efficiency of $k$ – Strong Nash Equilibria

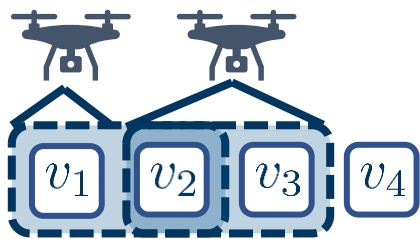
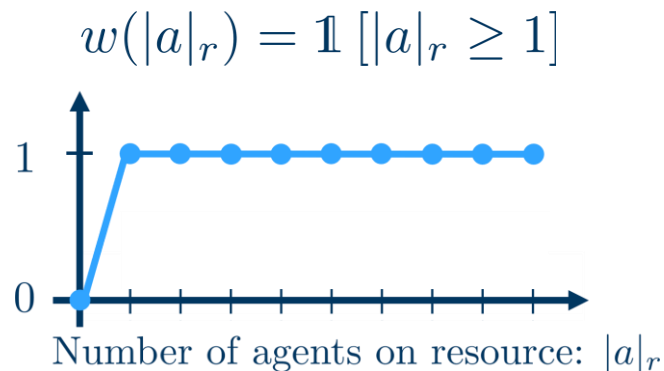
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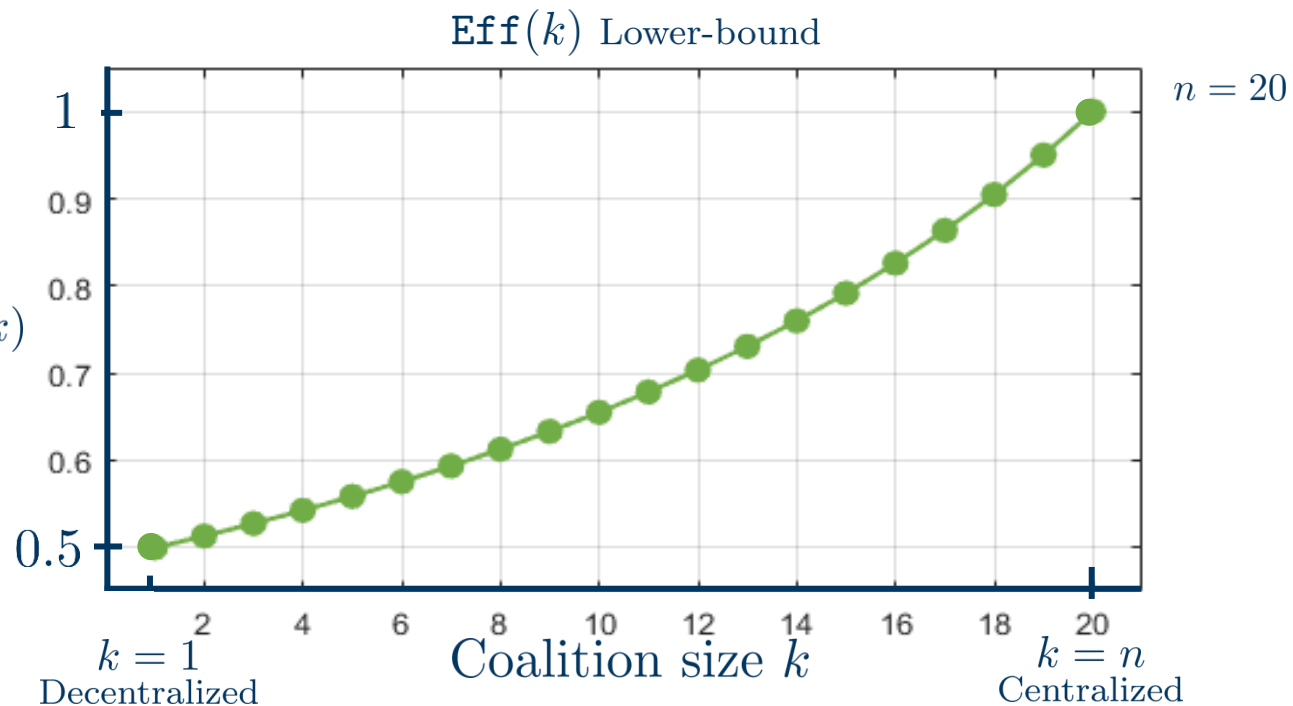
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$P^*(n, w, k)$



# Efficiency of $k$ – Strong Nash Equilibria

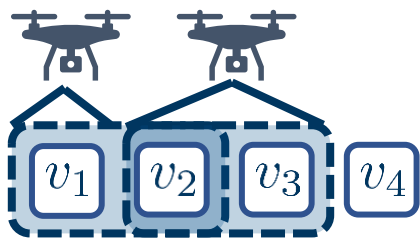
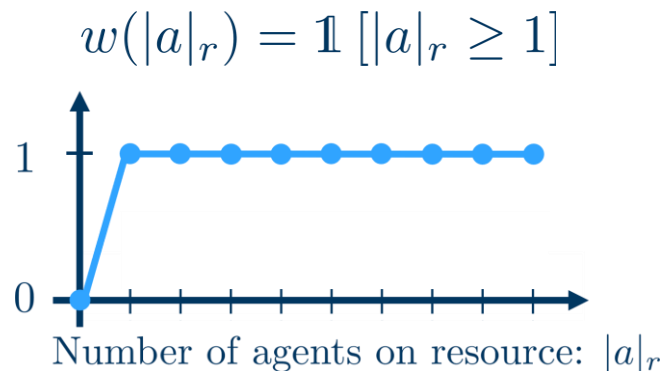
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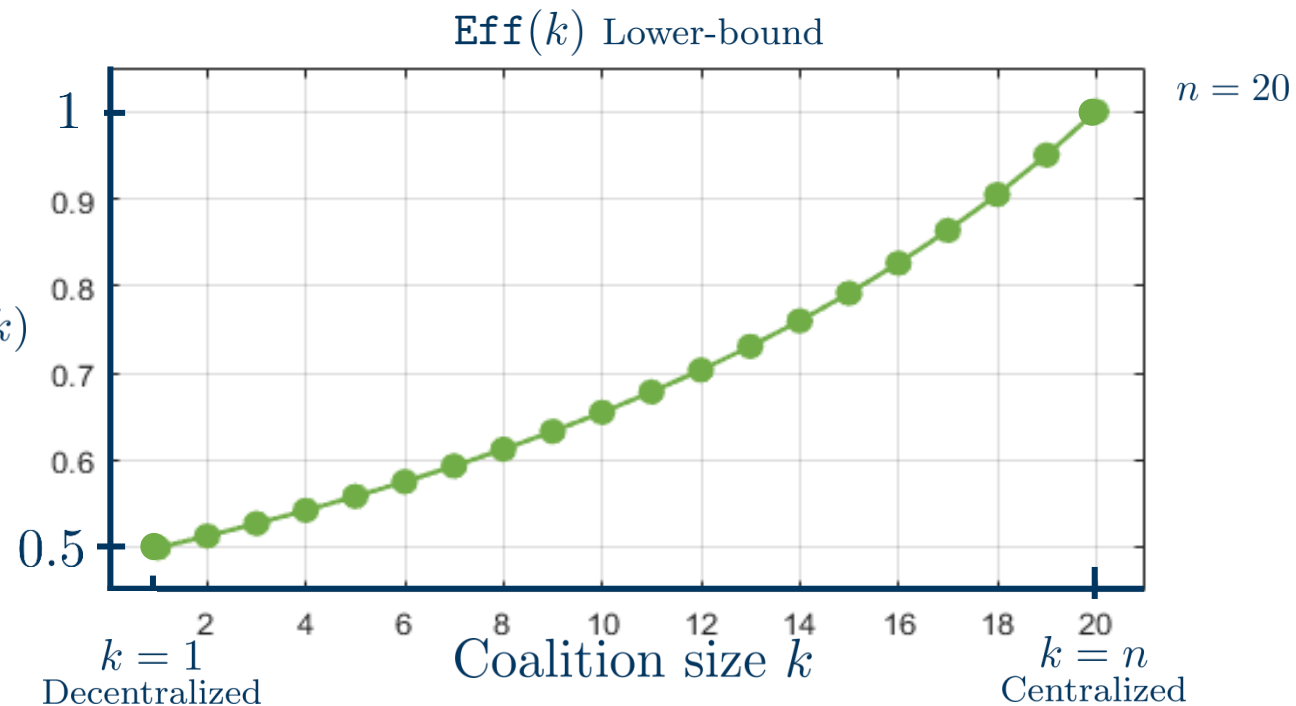
$$\text{Eff}(k) \geq P^*(n, w, k)$$

Example: Covering Problems



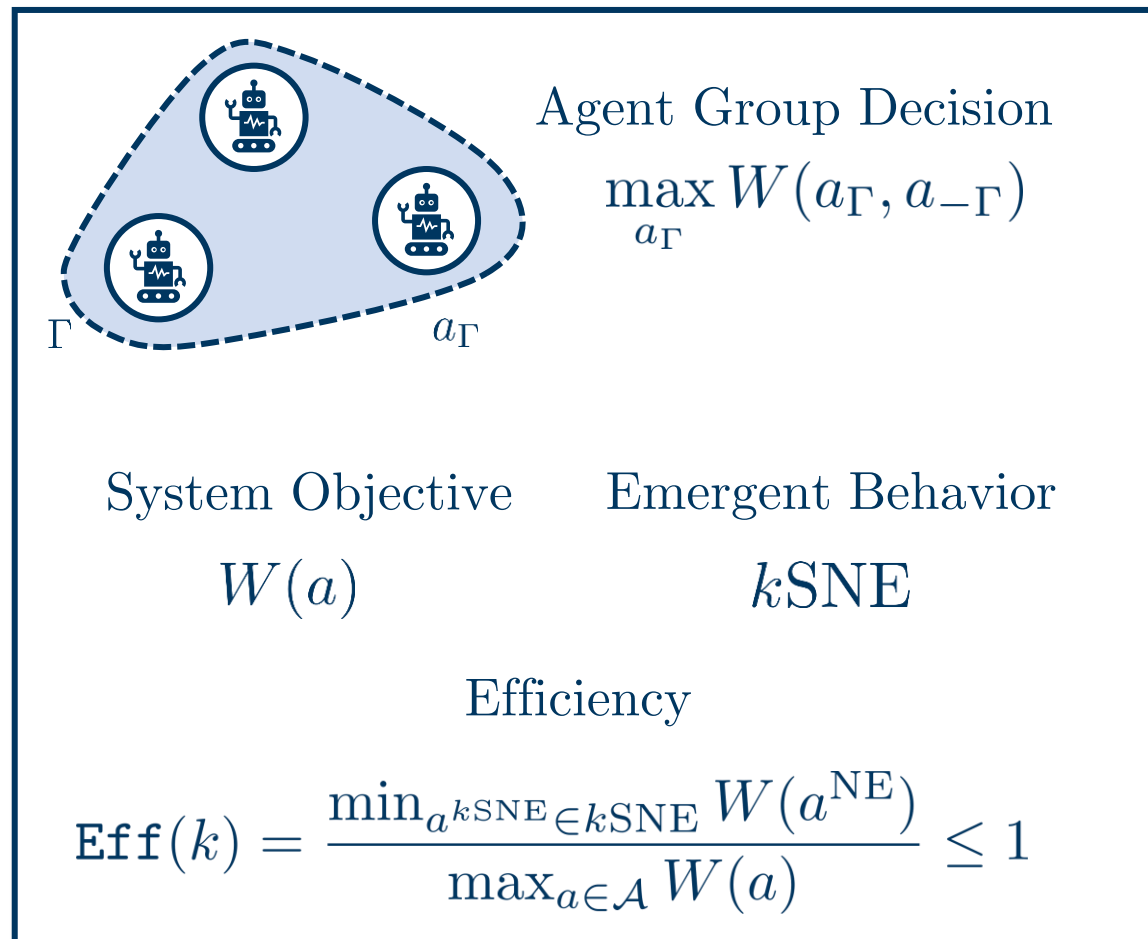
$$W(a) = v_1 + v_2 + v_3$$

$P^*(n, w, k)$



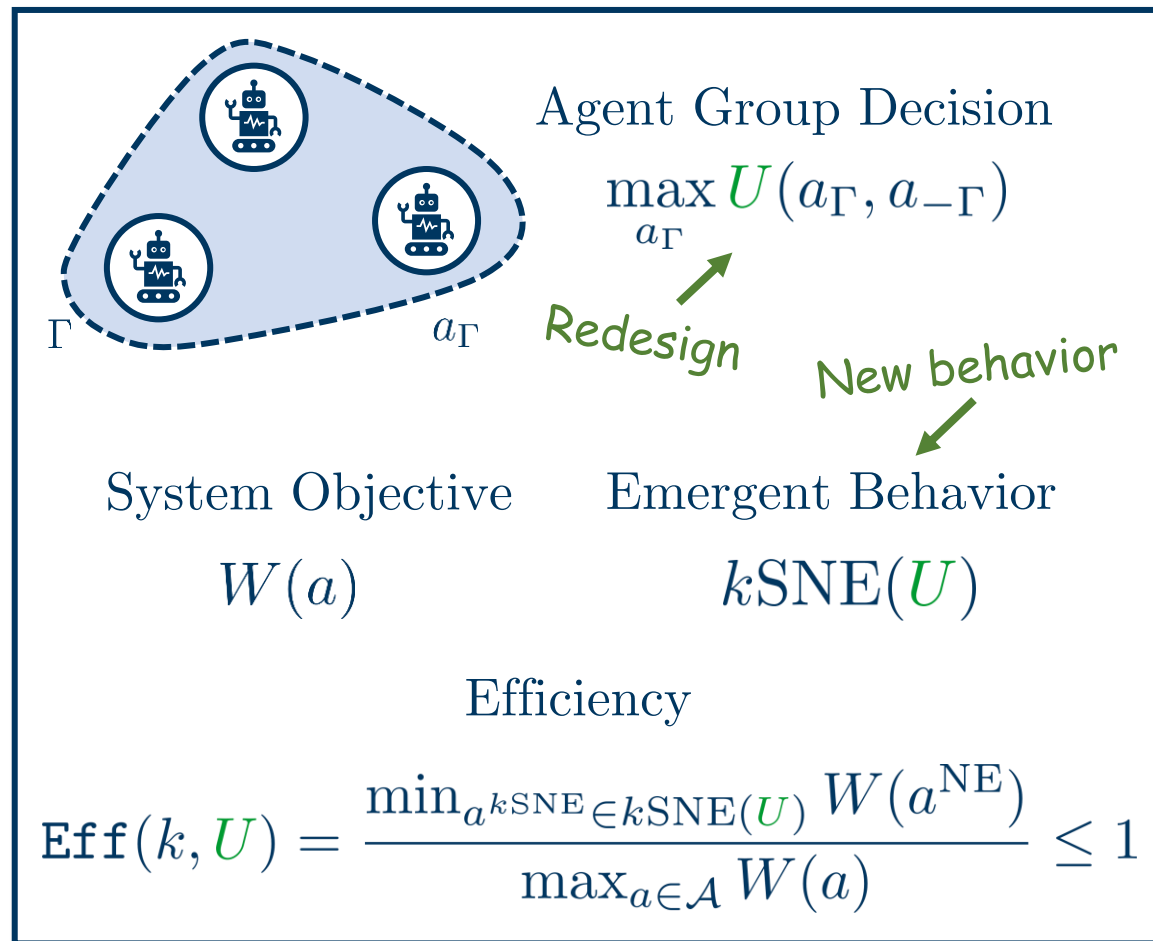
Thm1.1: Bridges gap between centralized and decentralized performance!

# Coalitional Utility Design



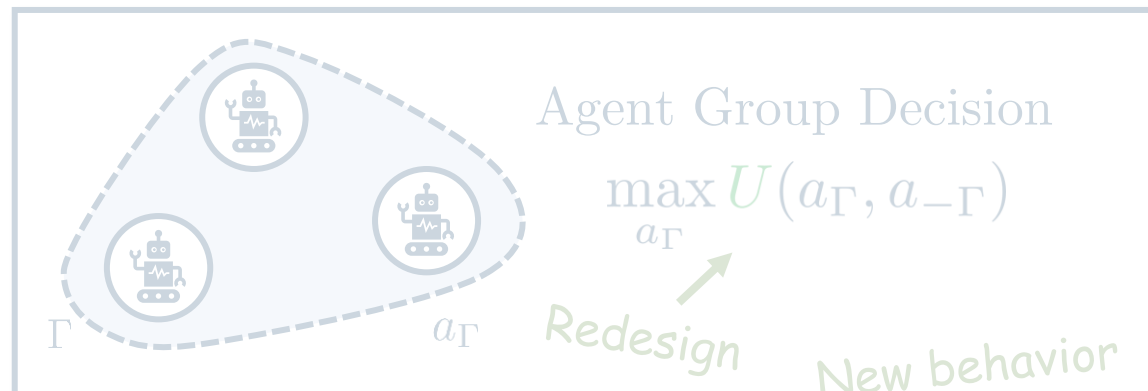
Design *group* objective to alter *global* behavior

# Coalitional Utility Design



Design *group* objective to alter *global* behavior

# Coalitional Utility Design



How much does *designing utility* help?

Efficiency

$$\text{Eff}(k, U) = \frac{\min_{a^{k\text{SNE}} \in k\text{SNE}(U)} W(a^{\text{NE}})}{\max_{a \in \mathcal{A}} W(a)} \leq 1$$

Design *group* objective to alter *global* behavior

# Coalition Utility Design

---

Proposition 1.2:

[BLF, Paccagnan, Pradeslki, Marden CDC23\*]

For a resource allocation problem  $(\mathcal{R}, N, \mathcal{A}, \{v_r\}_{r \in \mathcal{R}}, w)$ , under the optimal utility design, a  $k$ SNE approximates the optimal solution with

$$\text{Eff}(k, U^*) \geq Q^*(n, w, k),$$

where  $Q^*(n, w, k)$  is the solution to a linear program with  $nk + 1$  decision variables and  $\mathcal{O}(kn^3)$  constraints.

# Coalition Utility Design

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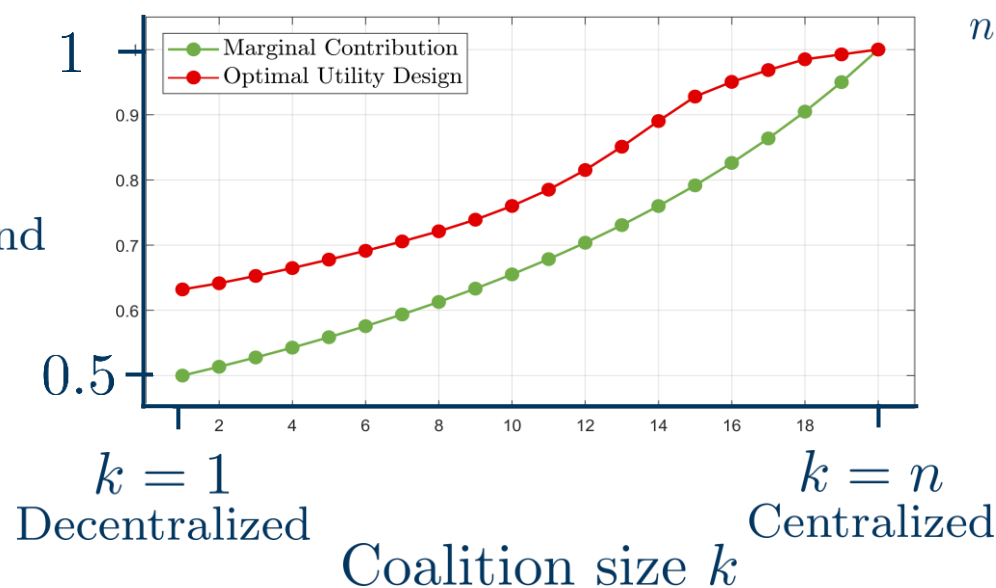
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where  $Q^*(n, w, k)$  is the solution to a linear program with  $nk + 1$  decision variables and  $\mathcal{O}(kn^3)$  constraints.

Example: Covering Problems

$$w(|a|_r) = \mathbb{1} [|a|_r \geq 1]$$

$\text{Eff}(k)$   
Lower-bound



$n = 20$

Utility Design can  
improve efficiency  
with coalitions



# A Performance/Complexity Trade-off

---

Distributed Decision-making  
(Bad performance / low complexity)



*Partially Collaborative*  
Decision-making

Centralized Decision-making  
(Best performance / high complexity)

# A Performance/Complexity Trade-off

---

Distributed Decision-making  
(Bad performance / low complexity)



*Partially Collaborative*  
Decision-making

Centralized Decision-making  
(Best performance / high complexity)

Increasing Communication

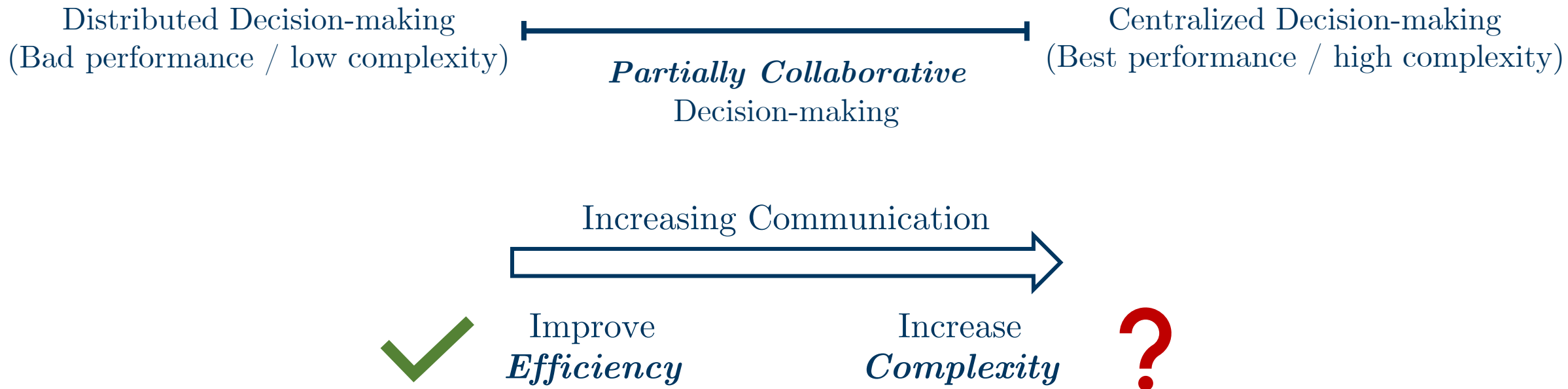


Improve  
*Efficiency*

Increase  
*Complexity*

# A Performance/Complexity Trade-off

---



We now have some understanding of performance,  
What happens to complexity?

# Complexity of $k$ SNE

---

How does communication affect *convergence rate*?

$n :=$  number of players

$m :=$  number of agent actions

$k :=$  size of coalitions

# Complexity of $k$ SNE

---

How does communication affect *convergence rate*?

Nash Equilibrium

$k$ -Strong Nash Equilibrium

Evaluating an equilibrium

$$\mathcal{O}(nm)$$

$$\mathcal{O}\left(\frac{n!}{(n-k)!k!}m^k\right)$$

$n :=$  number of players

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# Complexity of $k$ SNE

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Finding an equilibrium

$$\mathcal{O}(m^n)$$

?

As long  
or longer!

$n :=$  number of players

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# Complexity of $k$ SNE

---

How does communication affect *convergence rate*?

Nash Equilibrium

$k$ -Strong Nash Equilibrium

Evaluating an equilibrium

$$\mathcal{O}(nm)$$

$$\mathcal{O}\left(\frac{n!}{(n-k)!k!}m^k\right)$$

Finding an equilibrium

$$\mathcal{O}(m^n)$$

?

As long  
or longer!

The relative value of faster convergence is context dependent!  
 What techniques can we employ to *reduce complexity* with the  
*least sacrifice* to performance?

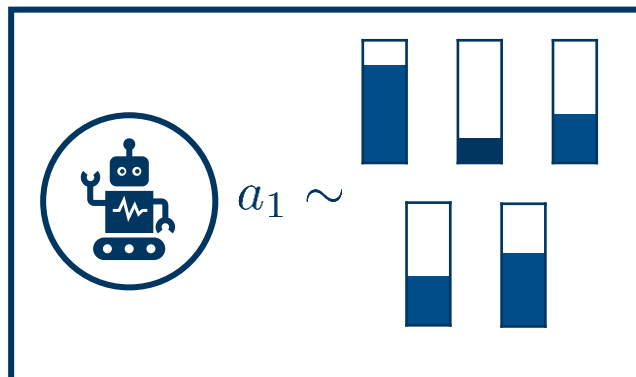
$n :=$  number of players

$m :=$  number of agent actions

$k :=$  size of coalitions

# Exploring the Trade-off

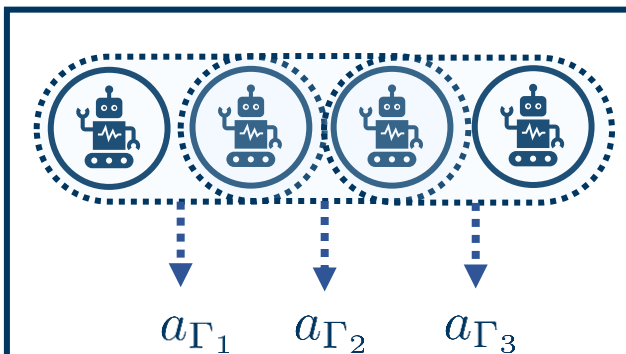
## Non-Deterministic Algorithms



Group regret-based decision-making

- Stochastic decision making to non-pure equilibrium concepts
- Coarse-correlated equilibria and smoothness

## Non-Equilibrium Algorithms

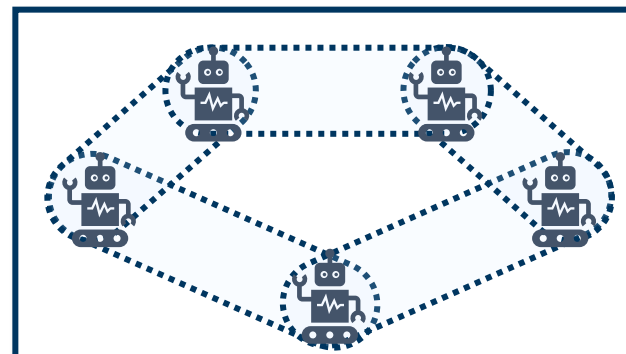


One-round walk & finite run time

- Each group revises their action finite times

$$\mathcal{O}\left(\frac{n!}{(n-k)!k!}m^k\right)$$

## Weaker Communication Structures

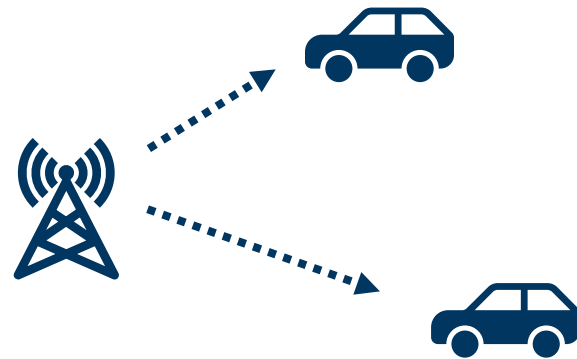


Not every subset communicates

- Lesser communication structure means worse performance guarantees
- Ideally, the reduction in complexity outweighs the loss in performance

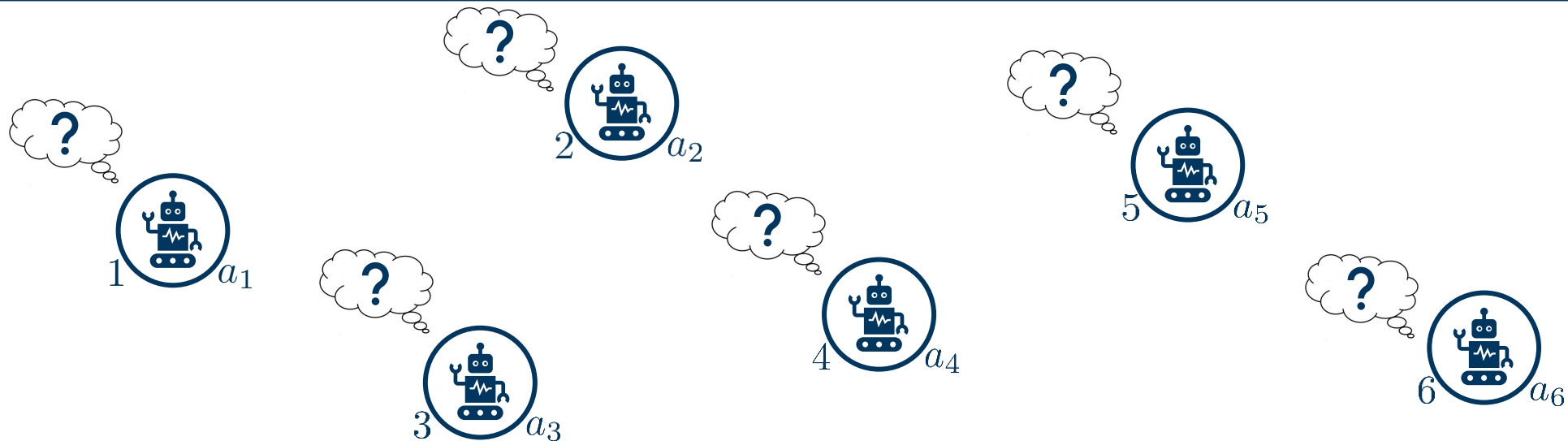


## II. Information Provisioning



# Information Provisioning

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Missing  
Information

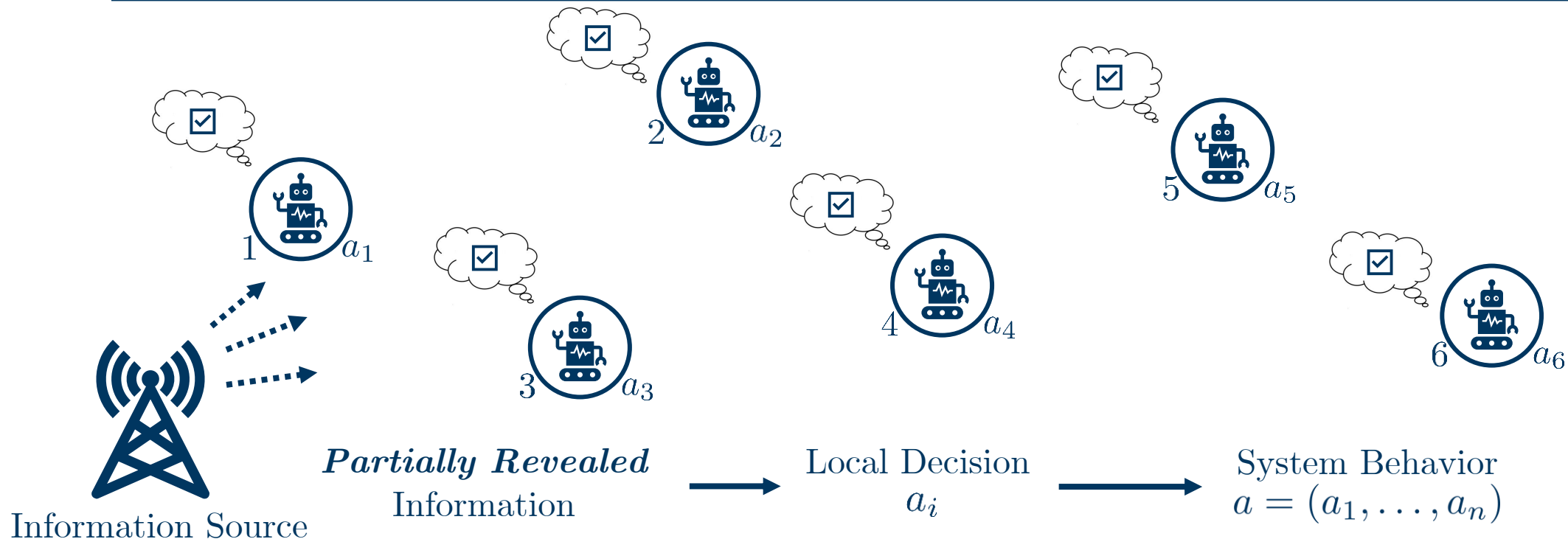


Local Decision  
 $a_i$



System Behavior  
 $a = (a_1, \dots, a_n)$

# Information Provisioning



How does *revealing information* to local decision makers affect system performance?

# Informing Decision-Makers

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Goal: surveil number of deer with a distributed drone team

# Informing Decision-Makers

---

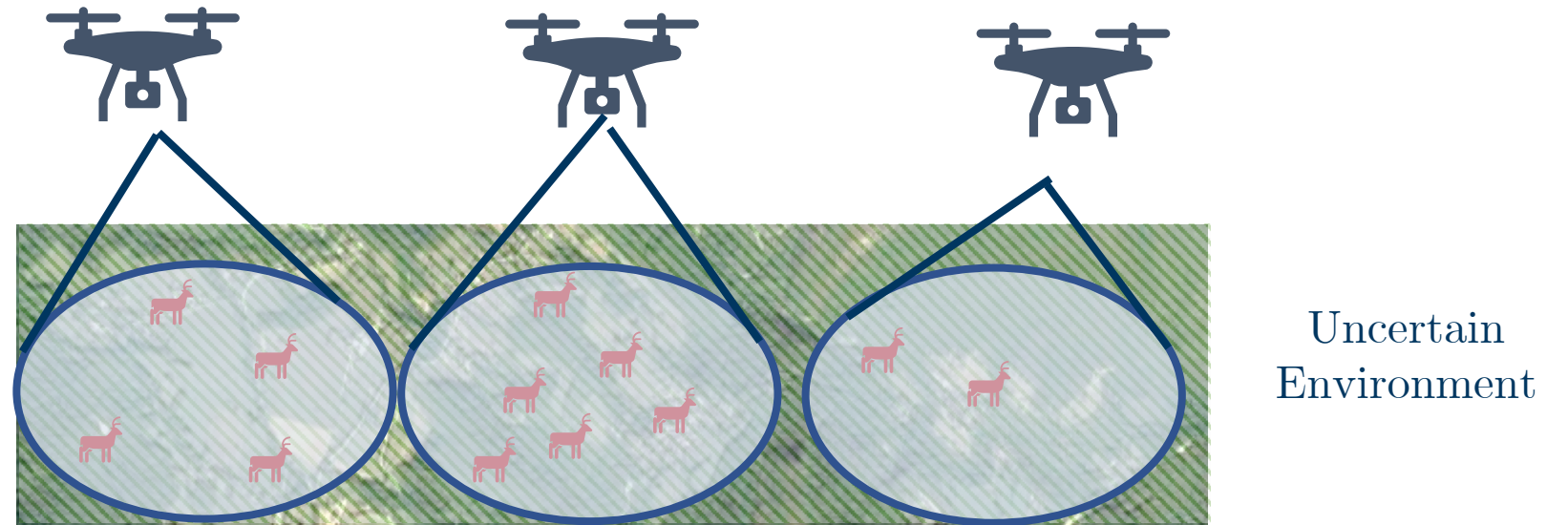


Uncertain  
Environment

Goal: surveil number of deer with a distributed drone team

# Informing Decision-Makers

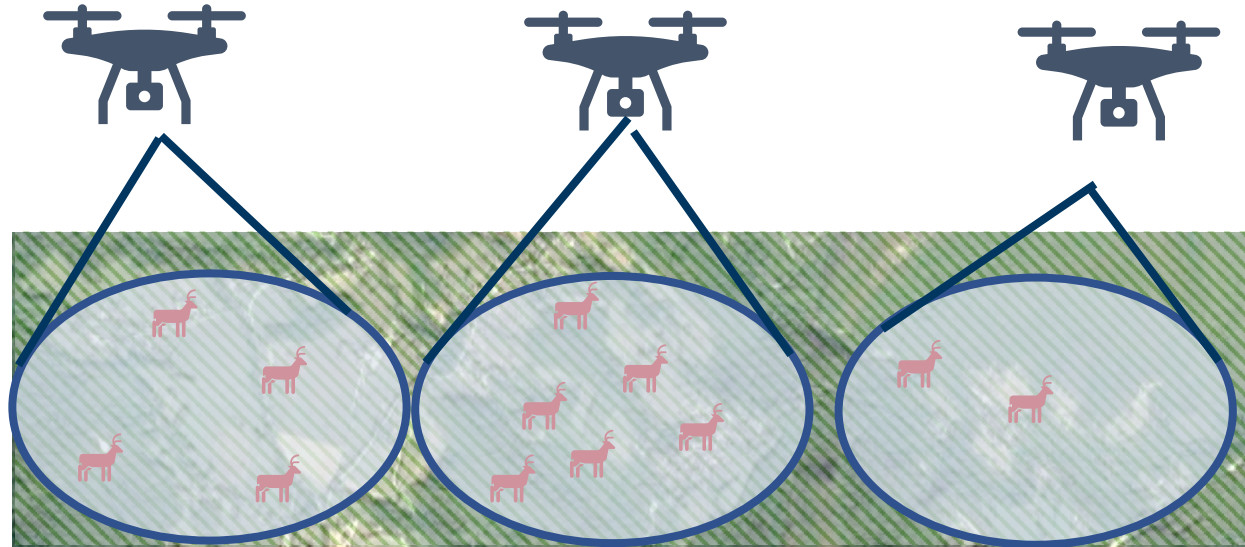
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Goal: surveil number of deer with a distributed drone team

# Informing Decision-Makers

---



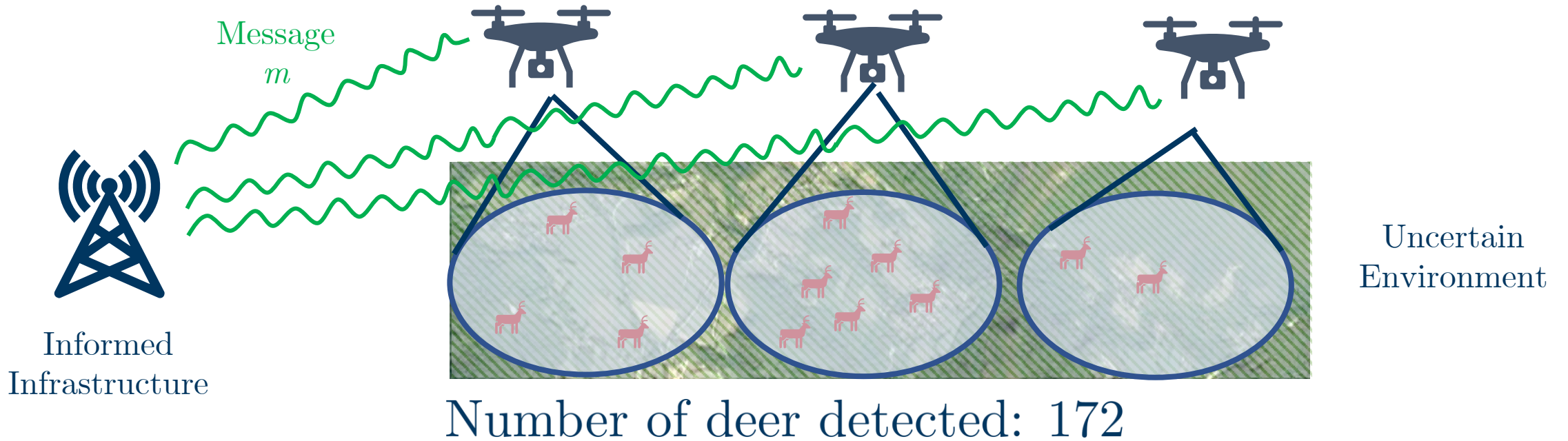
Uncertain  
Environment

Number of deer detected: 172

Goal: surveil number of deer with a distributed drone team

# Informing Decision-Makers

---



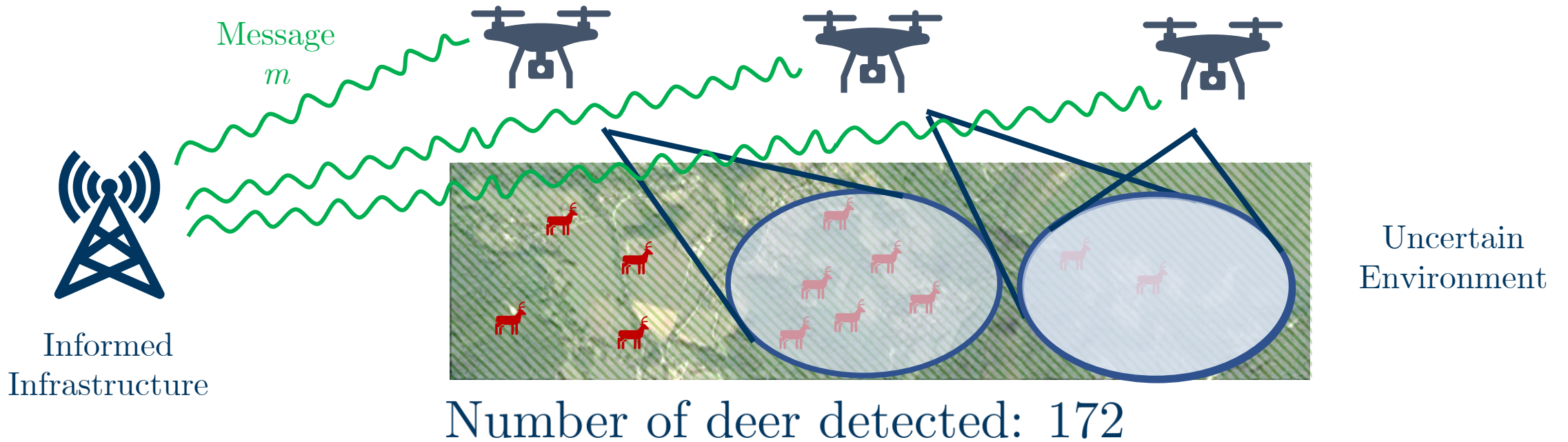
Goal: surveil number of deer with a distributed drone team

Proposal: send **messages** to inform drones of areas with more deer throughout day



# Informing Decision-Makers

---

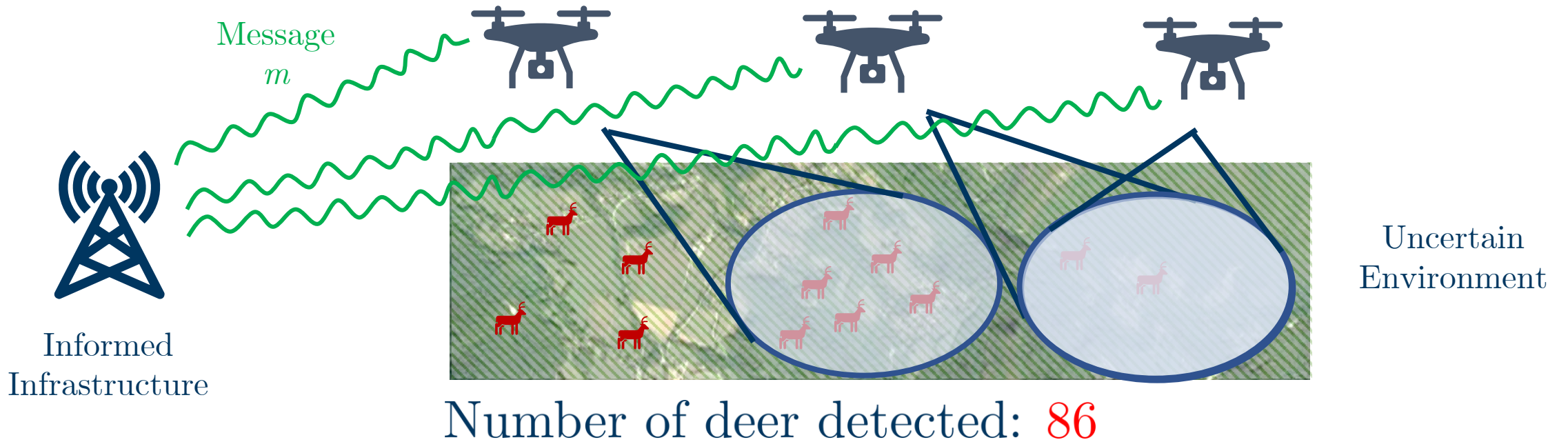


Goal: surveil number of deer with a distributed drone team

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# Informing Decision-Makers

---

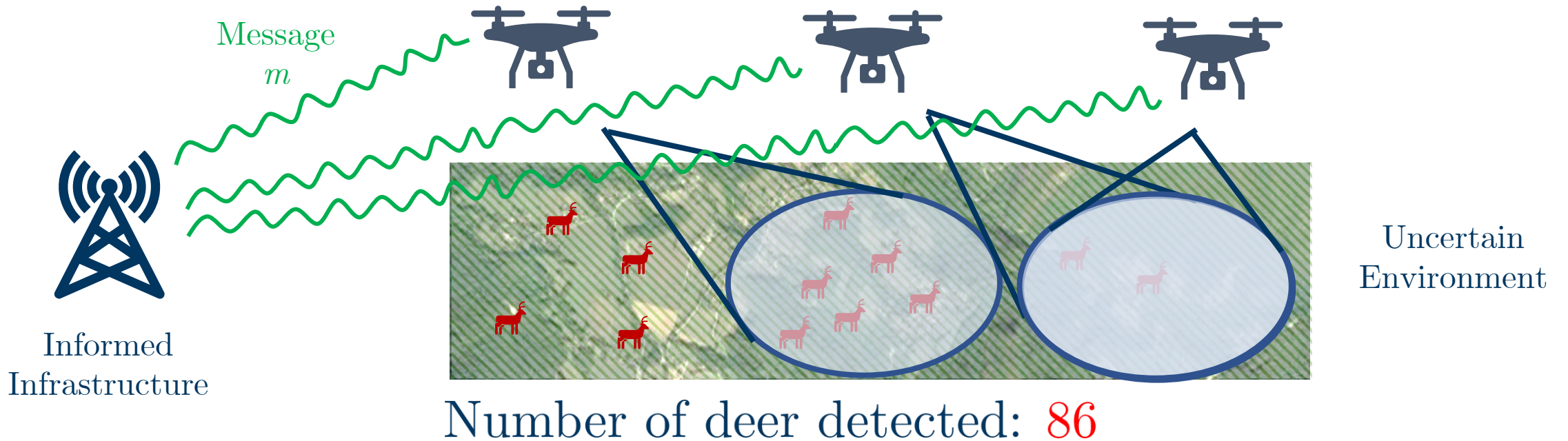


Goal: surveil number of deer with a distributed drone team

Proposal: send **messages** to inform drones of areas with more deer throughout day

Result: far *fewer* deer were detected?

# Informing Decision-Makers



Goal: surveil number of deer with a distributed drone team

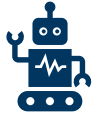
Proposal: send **messages** to inform drones of areas with more deer throughout day

Result: far *fewer* deer were detected?

Why did revealing information worsen system performance?

# Information Signaling

---



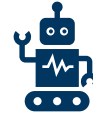
Agent 1

$$U_1(a; \alpha)$$



Agent 2

$$U_2(a; \alpha)$$



Agent 3

$$U_3(a; \alpha)$$

...



Agent  $n$

$$U_n(a; \alpha)$$

System state  $\alpha$  affects agent costs

# Information Signaling



System state  $\alpha$  affects agent costs

**Player knowledge:** full information on state

$\alpha$

**Nash Equilibrium**  $(a^{\text{NE}}, \alpha)$ :

$$U_i(a^{\text{NE}}; \alpha) \leq U_i(a'_i, a_{-i}^{\text{NE}}; \alpha)$$

$$\forall a'_i \in \mathcal{A}_i,$$

$$i \in N = \{1, \dots, n\}$$

# Information Signaling



System state  $\alpha$  affects agent costs

**Player knowledge:** a prior belief

$$\mu_0(x) = \mathbb{P}[\alpha = x]$$

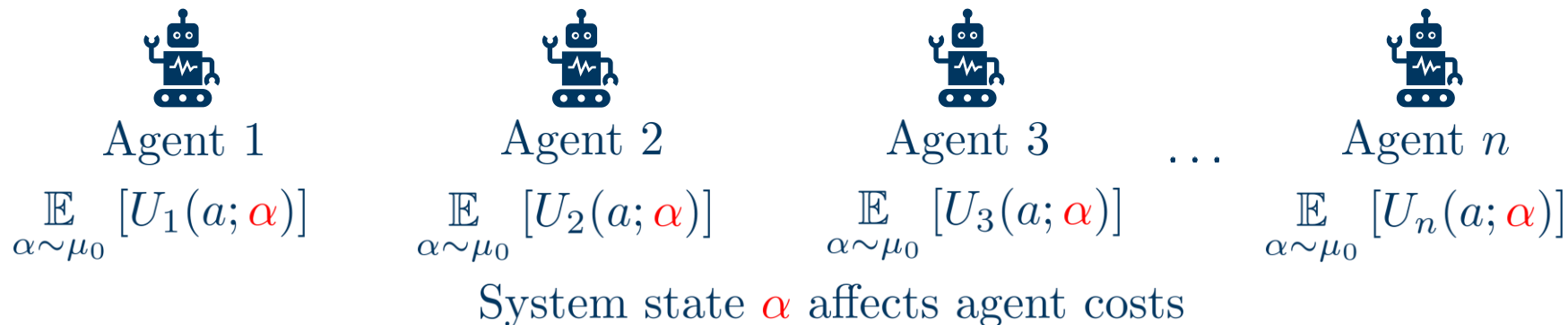
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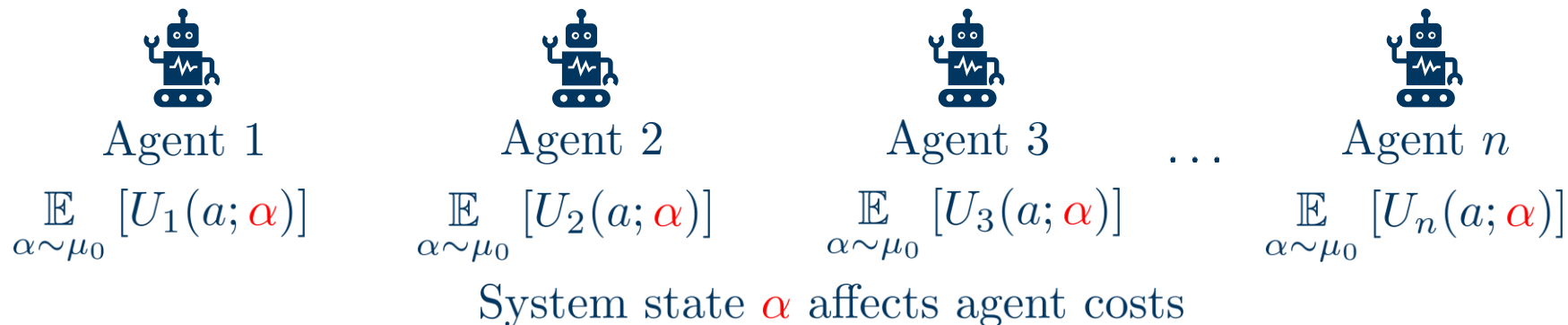
**Uninformed Nash Equilibrium**  $(a^{\text{NE}}, \mu_0)$ :

$$\mathbb{E}_{\alpha \sim \mu_0} [U_i(a^{\text{NE}}; \alpha)] \leq \mathbb{E}_{\alpha \sim \mu_0} [U_i(a'_i, a_{-i}^{\text{NE}}; \alpha)]$$

$$\forall a'_i \in \mathcal{A}_i,$$

$$i \in N = \{1, \dots, n\}$$

# Information Signaling



**Player knowledge:** partial information via signals

$$m = \pi(\alpha) \quad \pi : \text{supp}(\alpha) \rightarrow M$$

Message                      Signaling policy

**Uninformed Nash Equilibrium**  $(a^{\text{NE}}, \mu_0)$ :

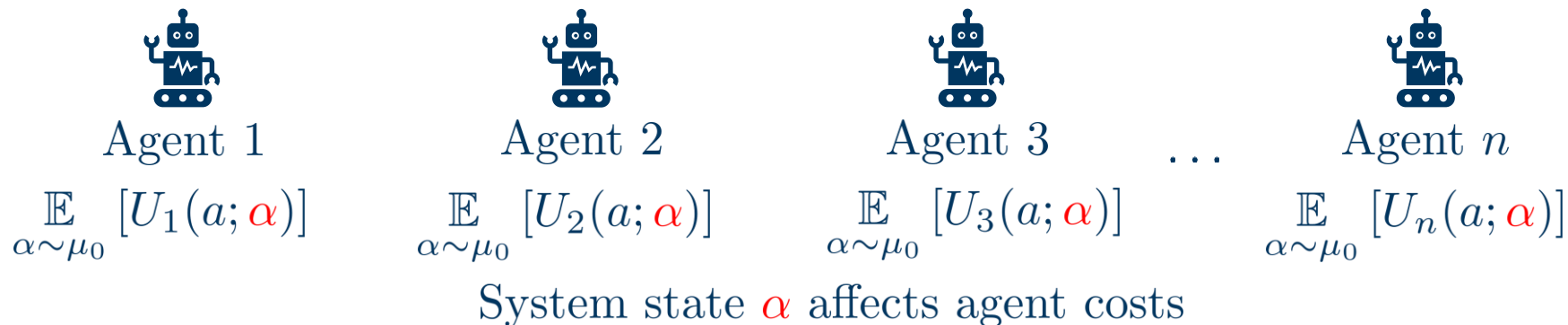
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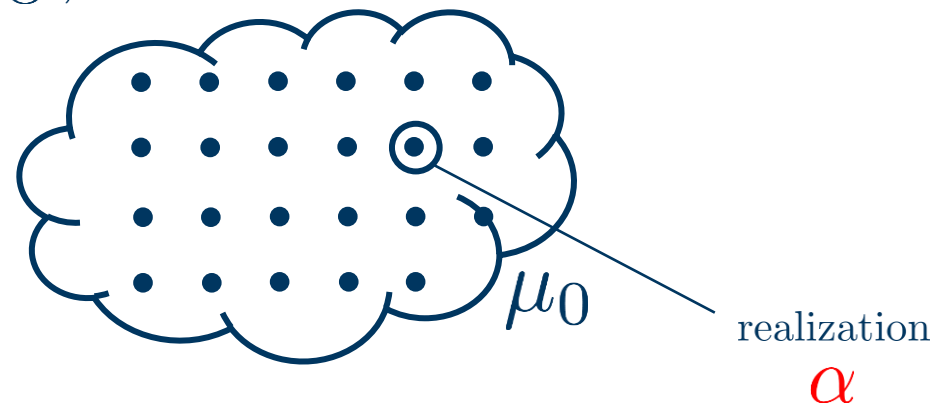
Message                      Signaling policy

**Uninformed Nash Equilibrium**  $(a^{\text{NE}}, \mu_0)$ :

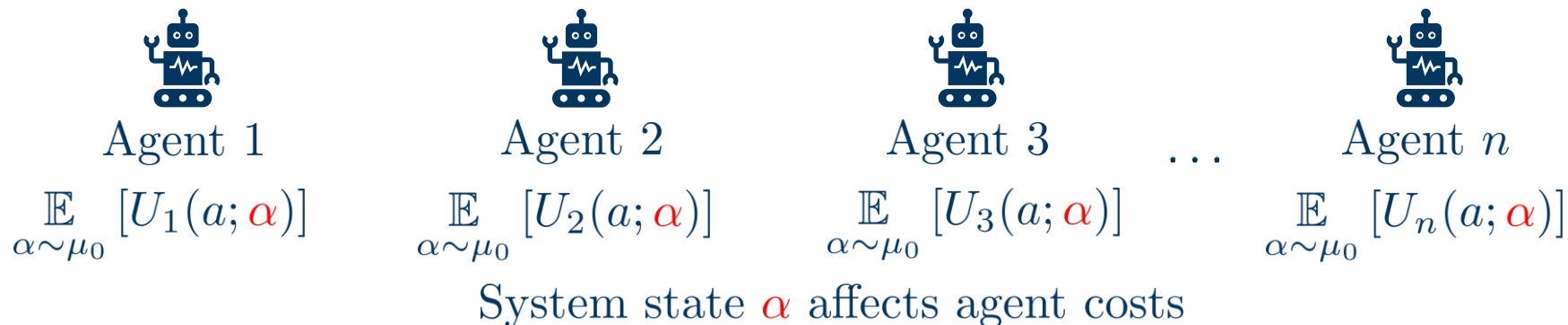
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E.g.,



# Information Signaling



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Message                      Signaling policy

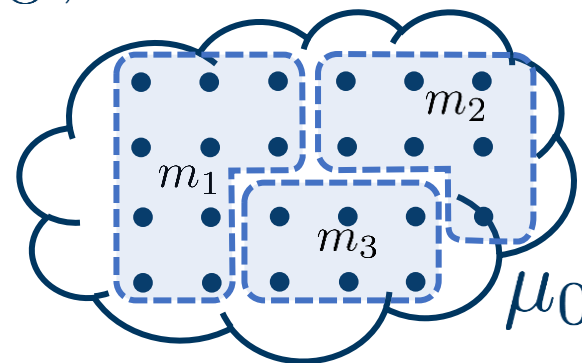
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$$\mathbb{E}_{\alpha \sim \mu_0} [U_i(a^{\text{NE}}; \alpha)] \leq \mathbb{E}_{\alpha \sim \mu_0} [U_i(a'_i, a_{-i}^{\text{NE}}; \alpha)]$$

$$\forall a'_i \in \mathcal{A}_i,$$

$$i \in N = \{1, \dots, n\}$$

E.g.,



Posterior Belief

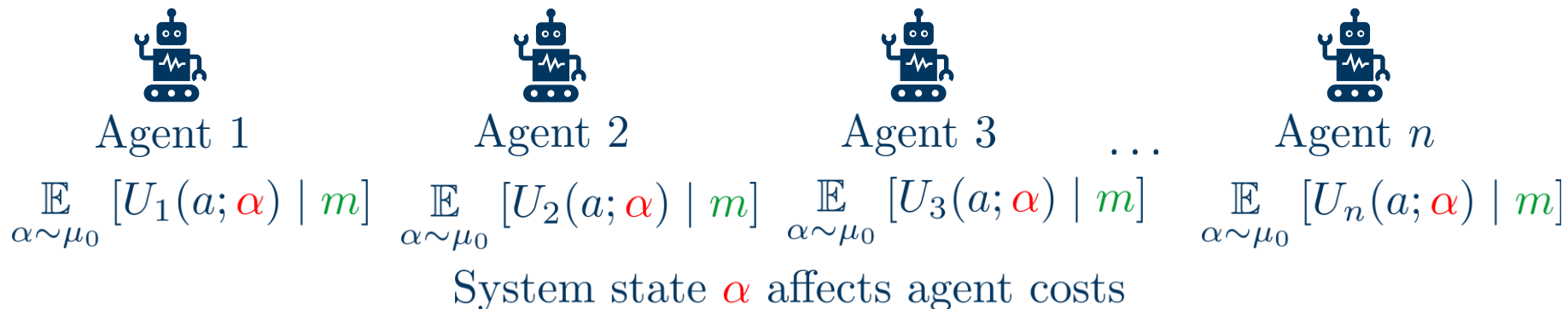
$$\mu_m$$

→ message  
 $m = \pi(\alpha)$

realization

$\alpha$

# Information Signaling



**Player knowledge:** partial information via signals

$$m = \pi(\alpha) \quad \pi : \text{supp}(\alpha) \rightarrow M$$

Message                      Signaling policy

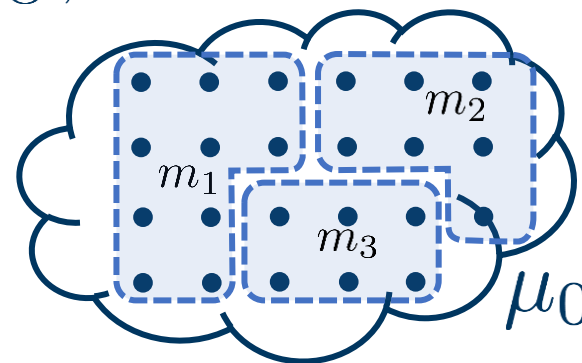
**Bayes-Nash Equilibrium** ( $\{a^{\text{NE}}(m)\}_{m \in M}, \mu_0, \pi$ ):

$$\mathbb{E}_{\alpha \sim \mu_0} [U_i(a^{\text{BNE}}; \alpha) | m] \leq \mathbb{E}_{\alpha \sim \mu_0} [U_i(a'_i, a_{-i}^{\text{BNE}}; \alpha) | m]$$

$$\forall a'_i \in \mathcal{A}_i, m \in M$$

$$i \in N = \{1, \dots, n\}$$

E.g.,



Posterior Belief

$$\mu_m$$

→ message  
 $m = \pi(\alpha)$

realization

$\alpha$

# The Value of Informing

---

Informed System Performance:  $W(a^{\text{BNE}}(\pi))$

Uninformed System Performance:  $W(a^{\text{NE}}(\mu_0))$

# The Value of Informing

---

$$\begin{array}{l} \text{Informed System Performance:} \\ \text{Uninformed System Performance:} \end{array} \quad \frac{W(a^{\text{BNE}}(\pi))}{W(a^{\text{NE}}(\mu_0))} =: \begin{array}{l} \text{VoI}(\pi) \\ \text{(value of informing)} \end{array}$$

# The Value of Informing

---

$$\begin{array}{l} \text{Informed System Performance:} \\ \text{Uninformed System Performance:} \end{array} \quad \frac{W(a^{\text{BNE}}(\pi))}{W(a^{\text{NE}}(\mu_0))} =: \text{VoI}(\pi) \quad \begin{array}{l} \\ \text{(value of informing)} \end{array}$$

Multiple Equilibria  $\rightarrow$  Multiple Perspectives

## Optimistic Perspective

$$\text{VoI}^+(\pi) = \frac{\max_{a^{\text{BNE}} \in \text{BNE}(\pi)} \mathbb{E}[W(a^{\text{BNE}})]}{\max_{a^{\text{NE}} \in \text{NE}(\mu_0)} \mathbb{E}[W(a^{\text{NE}})]}$$

Gain of *best-case* equilibrium performance

## Pessimistic Perspective

$$\text{VoI}^-(\pi) = \frac{\min_{a^{\text{BNE}} \in \text{BNE}(\pi)} \mathbb{E}[W(a^{\text{BNE}})]}{\min_{a^{\text{NE}} \in \text{NE}(\mu_0)} \mathbb{E}[W(a^{\text{NE}})]}$$

Gain of *worst-case* equilibrium performance

# Bounding the Value of Informing

---

 Covering Problems  $w_r(|a|_r) = \alpha_r \mathbb{1}[|a|_r \geq 1]$  

System Objective

$$U(a; \alpha) = W(a; \alpha) = \sum_{r \in \mathcal{R}} \alpha_r \mathbb{1}[|a|_r > 0]$$

Thm 2.1 [BLF, D.Paccagnan, J.R.Marden LCSS\*]

For a resource allocation game with the covering objective:

$$1 \leq \text{VoI}^+(\pi) \leq |M|$$

$$1/2 \leq \text{VoI}^-(\pi) \leq 2|M|$$

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Observations:

- Revealing information only helps best-case
- A richer message space leads to greater opportunities for improvement ( $|M|$ )
- Revealing information can *worsen worst-case*



# Bounding the Value of Informing

 Covering Problems  $w_r(|a|_r) = \alpha_r \mathbb{1}[|a|_r \geq 1]$  

System Objective

$$U(a; \alpha) = W(a; \alpha) = \sum_{r \in \mathcal{R}} \alpha_r \mathbb{1}[|a|_r > 0]$$

Optimal Pessimistic Utility Design

$$U(a; \alpha) = \sum_{r \in \mathcal{R}} \alpha_r (|a|_r - 1)! \frac{\frac{1}{(n-1)(n-1)!} + \sum_{i=|a|_r}^{n-1} \frac{1}{i!}}{\frac{1}{(n-1)(n-1)!} + \sum_{i=1}^{n-1} \frac{1}{i!}}$$

Thm 2.1 [BLF, D.Paccagnan, J.R.Marden LCSS\*]

For a resource allocation game with the covering objective:

$$1 \leq \text{VoI}^+(\pi) \leq |M|$$

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Observations:

- Revealing information only helps best-case
- A richer message space leads to greater opportunities for improvement ( $|M|$ )
- Revealing information can *worsen worst-case*

Thm 2.2 [BLF, D.Paccagnan, J.R.Marden LCSS\*]

For a resource allocation game with the covering objective and pessimistic design:

$$1 - \frac{1}{e} \leq \text{VoI}^+(\pi) \leq \left(1 - \frac{1}{e}\right)^{-1} |M|$$

$$1 - \frac{1}{e} \leq \text{VoI}^-(\pi) \leq \left(1 - \frac{1}{e}\right)^{-1} |M|$$

Observation:

- Improving worst-case guarantee has negative consequences on best-case guarantee

# Trade-off in Best-case/Worst-case Guarantees

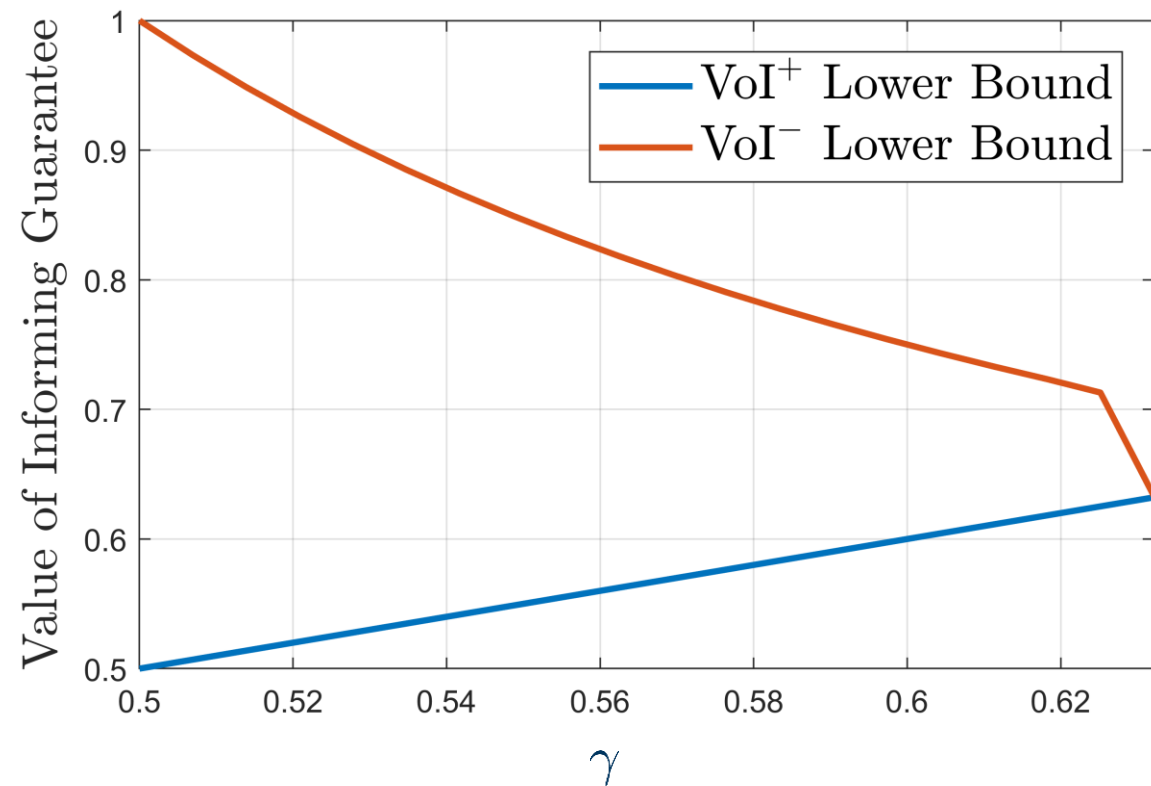
Prop. 2.3 [BLF, D.Paccagnan, J.R.Marden LCSS\*]

Given a desired  $\text{VoI}^-$  guarantee of  $\gamma \in (1/2, 1 - 1/e]$  and  $n \geq 2$ ,  $\text{VoI}^+$  satisfies

$$\text{VoI}^+ \geq Z(\gamma, n)$$

where  $Z(\gamma, n)$  equals

$$\frac{1}{1 + \max_{1 \leq j \leq n-1} j j! \left( 1 - \left( \frac{1}{\gamma} - 1 \right) \sum_{i=1}^j \frac{1}{i!} \right)}.$$



Designing utilities for better pessimistic guarantees  
worsens the optimistic guarantees

# Direction: Optimal Signaling Mechanism

---

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} W(a^{\text{BNE}}(\pi))$$

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---

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} W(a^{\text{BNE}}(\pi))$$

Optimist Optimal

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} \max_{a^{\text{BNE}} \in \text{BNE}(\pi)} W(a^{\text{BNE}})$$

Reveal everything

Pessimist Optimal

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} \min_{a^{\text{BNE}} \in \text{BNE}(\pi)} W(a^{\text{BNE}})$$

Computationally intractable

# Direction: Optimal Signaling Mechanism

---

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} \mathbb{E}_{a \in \text{BNE}(\pi)} [W(a)]$$

Average Case Performance

---

## Algorithm 1 Bandit Approach to Optimal Signals

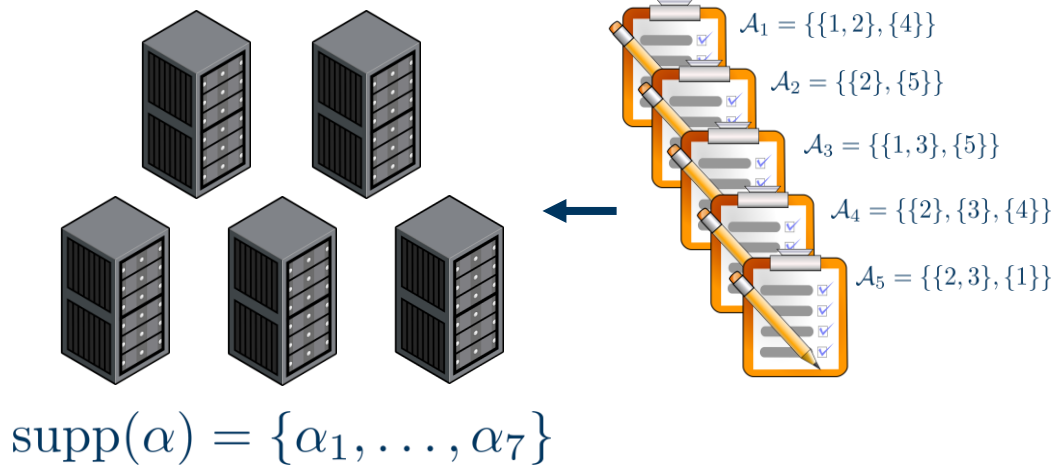
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- 1: Pick signaling policy  $\pi$  via UCB
  - 2: Sample State  $\alpha$  and initial allocation  $a_0$
  - 3: Find BNE via Best-Response  $\text{BR}(a_0, \pi(\alpha))$
  - 4: Record Equilibrium Reward
-

# Direction: Optimal Signaling Mechanism

$$\pi^* \in \arg \max_{\pi: \text{supp}(\alpha) \rightarrow M} \mathbb{E}_{a \in \text{BNE}(\pi)} [W(a)]$$

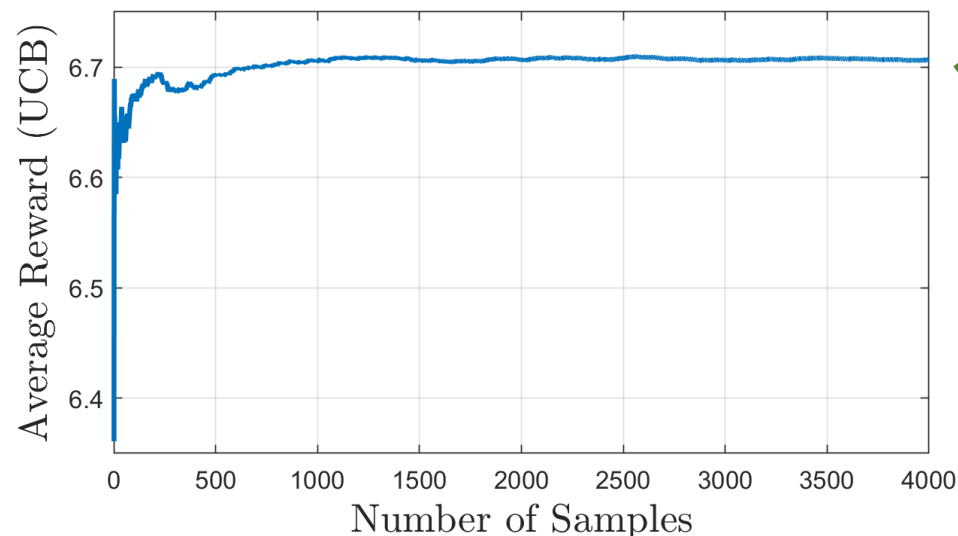
Average Case Performance



$$w_r(|a|_r) = \alpha_r |a|_r \exp(-c|a|_r)$$

## Algorithm 1 Bandit Approach to Optimal Signals

- 1: Pick signaling policy  $\pi$  via UCB
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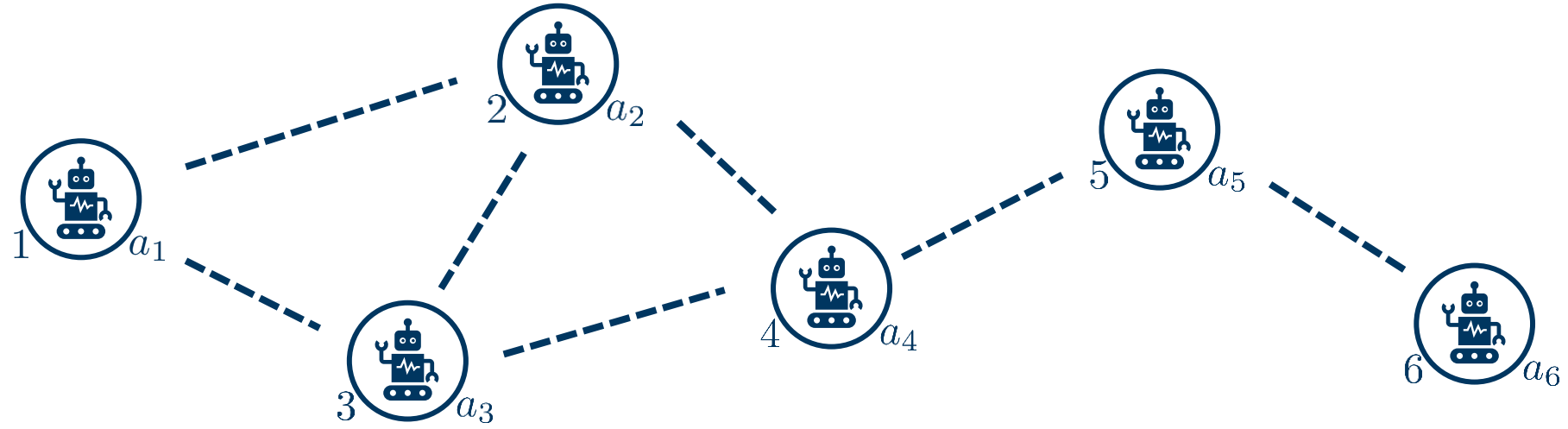
$$\pi^* = \{\{1\}, \{2, 6\}, \{3\}, \{4\}, \{5, 7\}\}$$

# III. Unreliable Communicators



# Sub-System Failures

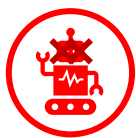
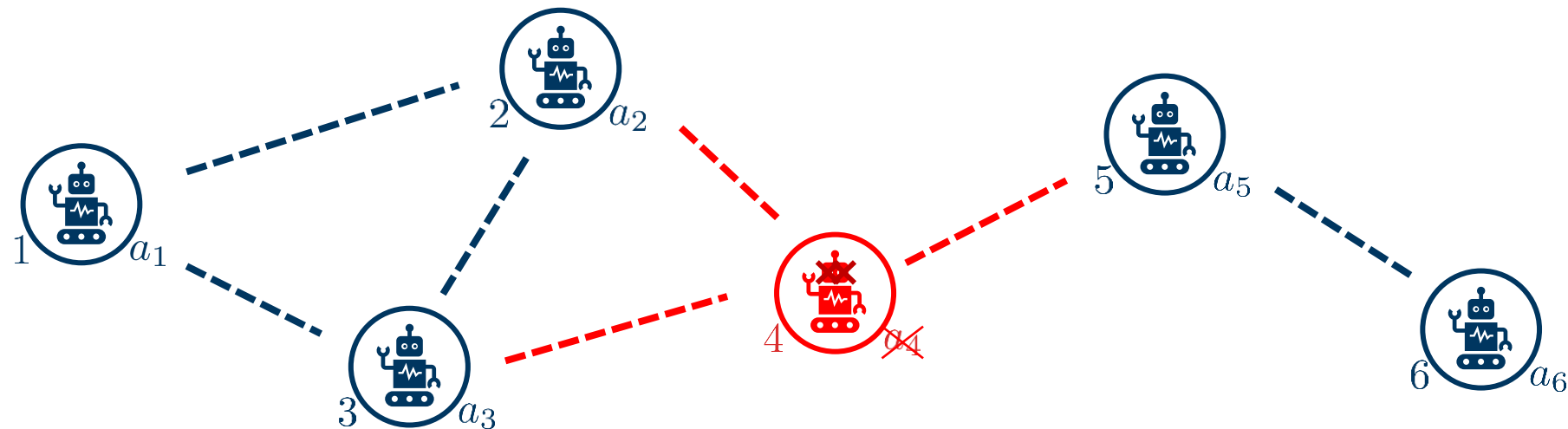
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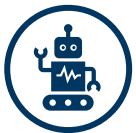


# Sub-System Failures

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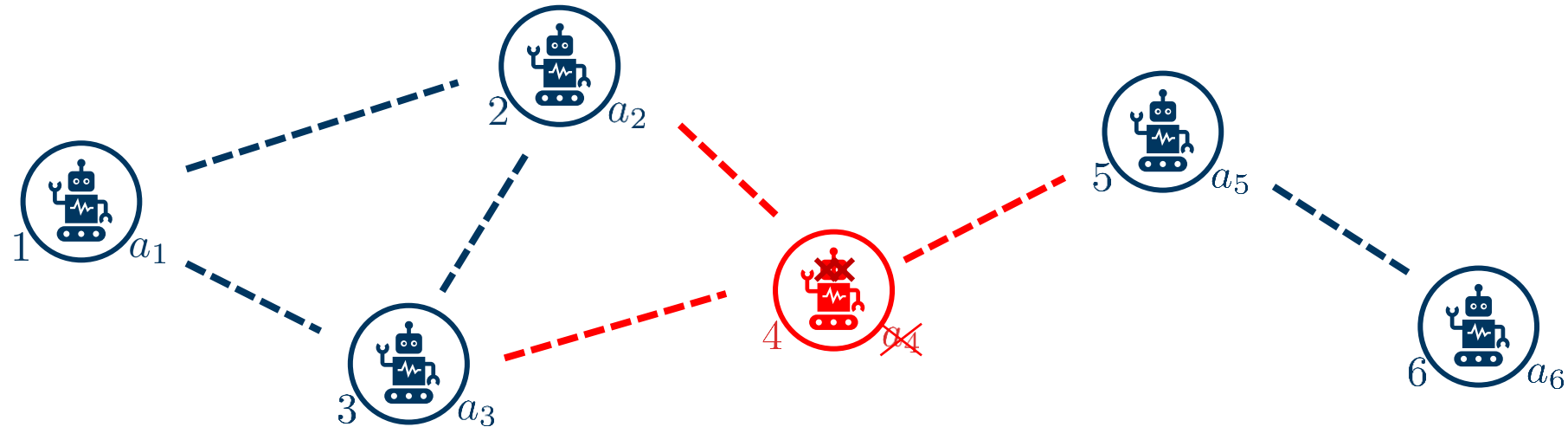


**Defective Agent:** Agent still communicates but does not contribute to the system objective

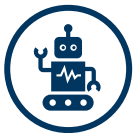


Nominal agents cannot determine which agents are defective. They must operate under the assumption some agents might be.

# Sub-System Failures



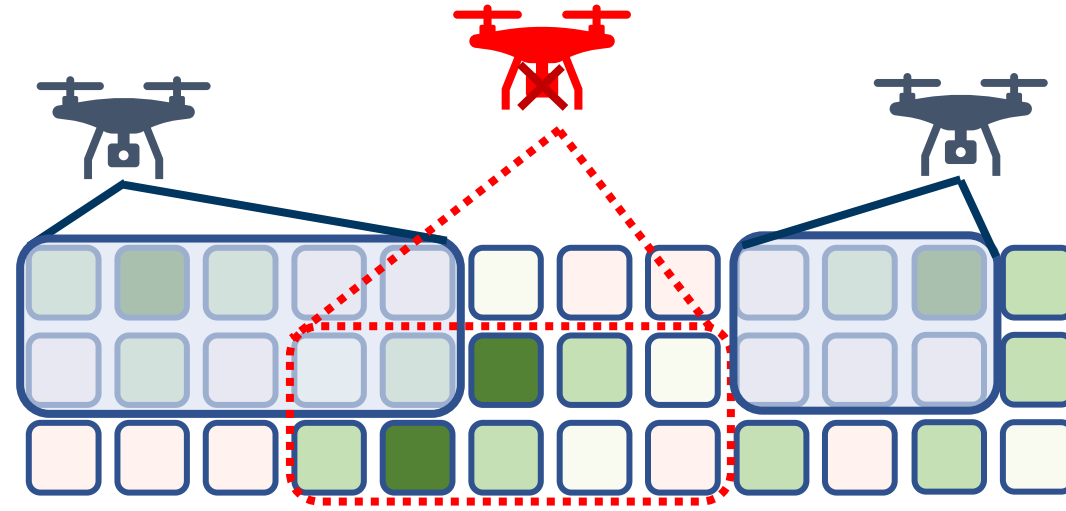
**Defective Agent:** Agent still communicates but does not contribute to the system objective



Nominal agents cannot determine which agents are defective. They must operate under the assumption some agents might be.

How should agents be designed when others may be **defective**?

# Sub-System Failures



Nominal Agents

$$i \in N$$

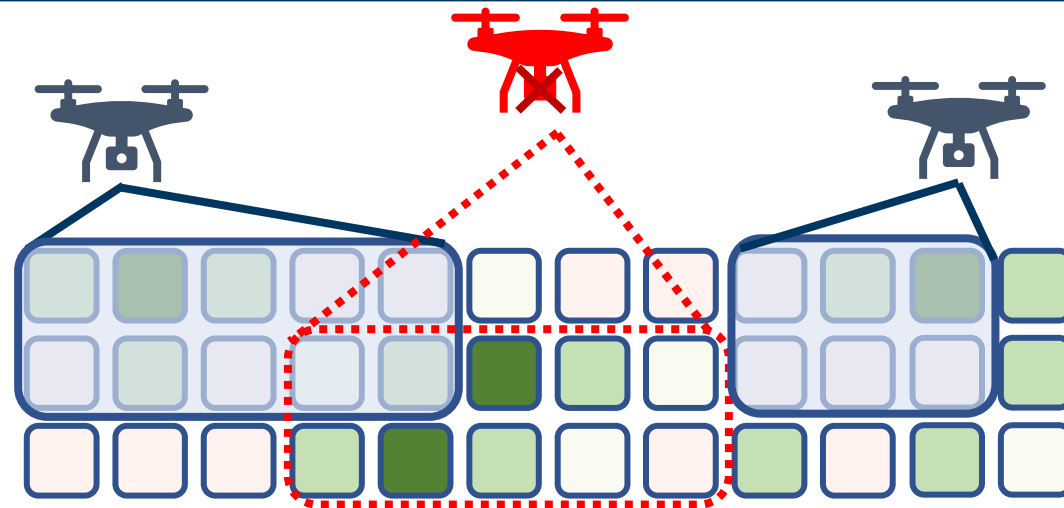


Defective Agents  
with Fixed Action

$$j \in D$$

$$z_j \in 2^{\mathcal{R}}$$

# Sub-System Failures



Nominal Agents

$$i \in N$$

System Welfare

$$W(a) = \sum_{r \in \mathcal{R}} v_r w(|a|_r)$$

Only depends on regular agents



Defective Agents  
with Fixed Action

$$j \in D$$

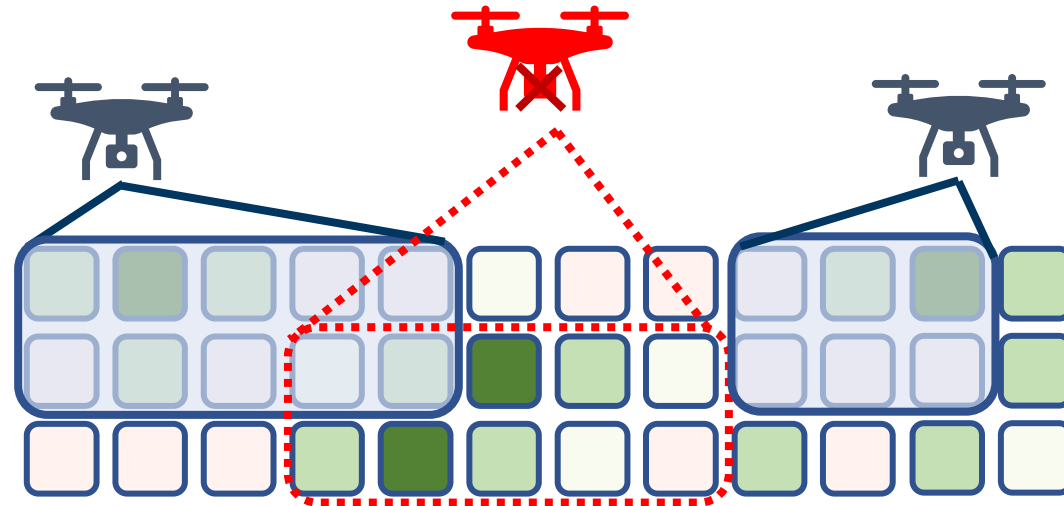
$$z_j \in 2^{\mathcal{R}}$$

Utility

$$U_i(a, z) = \sum_{r \in \mathcal{R}} v_r u(|a|_r + |z|_r)$$

Depends on regular **and** defective agents

# Sub-System Failures



Nominal Agents

$$i \in N$$

System Welfare

$$W(a) = \sum_{r \in \mathcal{R}} v_r w(|a|_r)$$

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Utility

$$U_i(a, z) = \sum_{r \in \mathcal{R}} v_r u(|a|_r + |z|_r)$$

Depends on regular *and* defective agents

Defective agents can cause nominal agents to avoid high valued resources

**Robust Design:** choose utility functions to promote more overlap

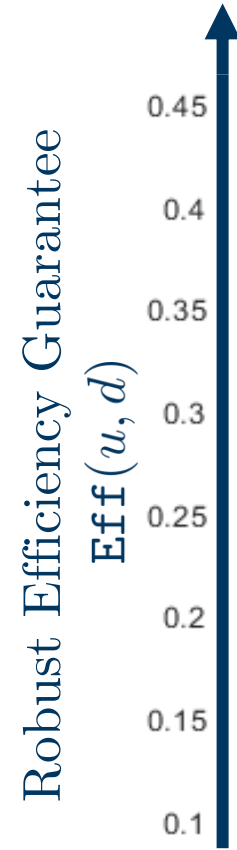
Proposition 1.3:

[CDC21,DGAA]

The optimal utility design and associated efficiency is the solution to a linear program.

# Nominal/Robust Performance Trade-off

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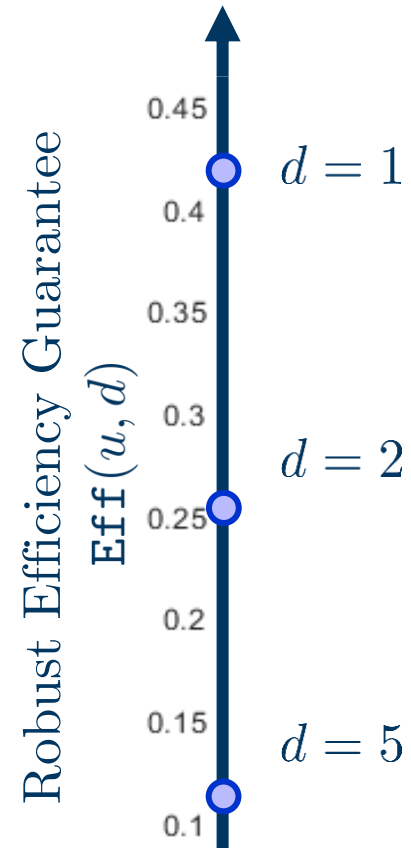


# Nominal/Robust Performance Trade-off

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● Nominal Utility Rule

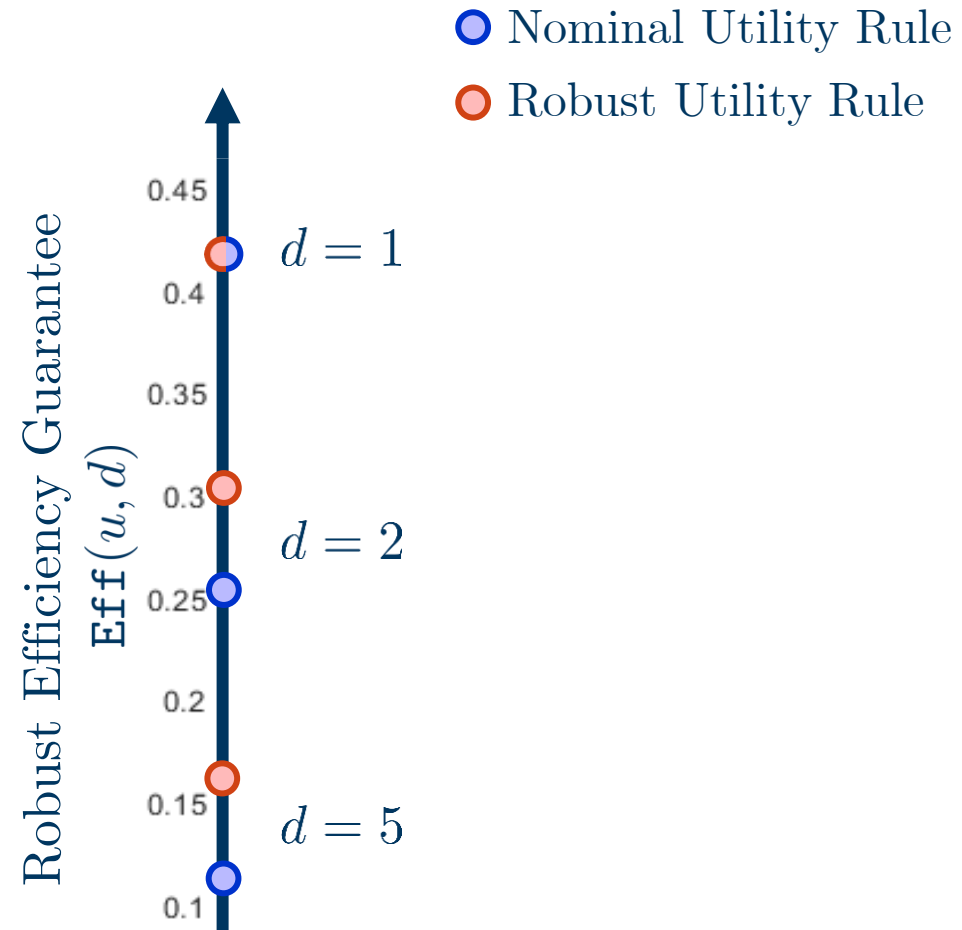
In the class of covering games,



$d := \#$  defective agents

# Nominal/Robust Performance Trade-off

In the class of covering games,



$d := \#$  defective agents



# Nominal/Robust Performance Trade-off

Theorem 1.4 [CDC21, DGAA]

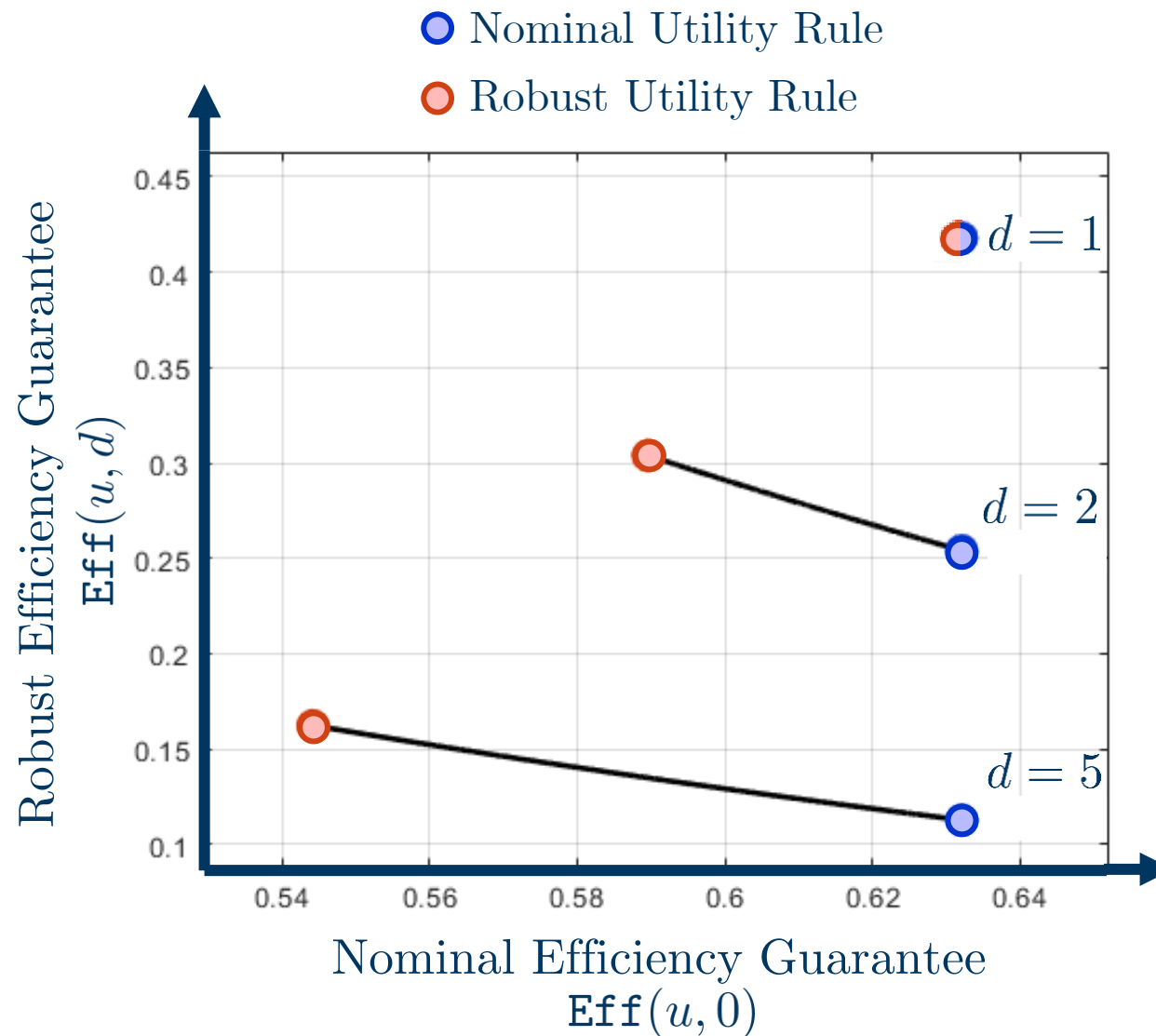
In the class of covering games, if

$$\text{Eff}(u, d) \geq \frac{Z_d + \frac{e}{e-1}}{1 + tZ_d},$$

where  $Z_d = d! \frac{e^{-\sum_{i=0}^{d-1} \frac{1}{i!}}}{e-1} - 1$  and  $t \in [0, 1]$ , then

$$\text{Eff}(u, 0) \leq \frac{(e-1)(1 + tZ_d)}{1 + (e-1)(1 + tZ_d)}.$$

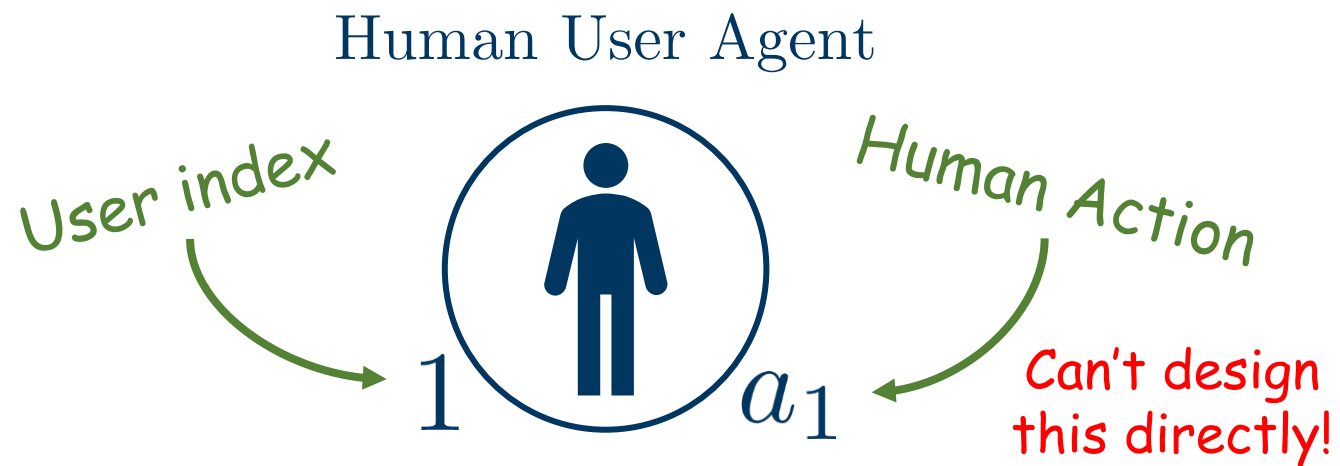
$d := \#$  defective agents



## IV. Directions & Conclusions

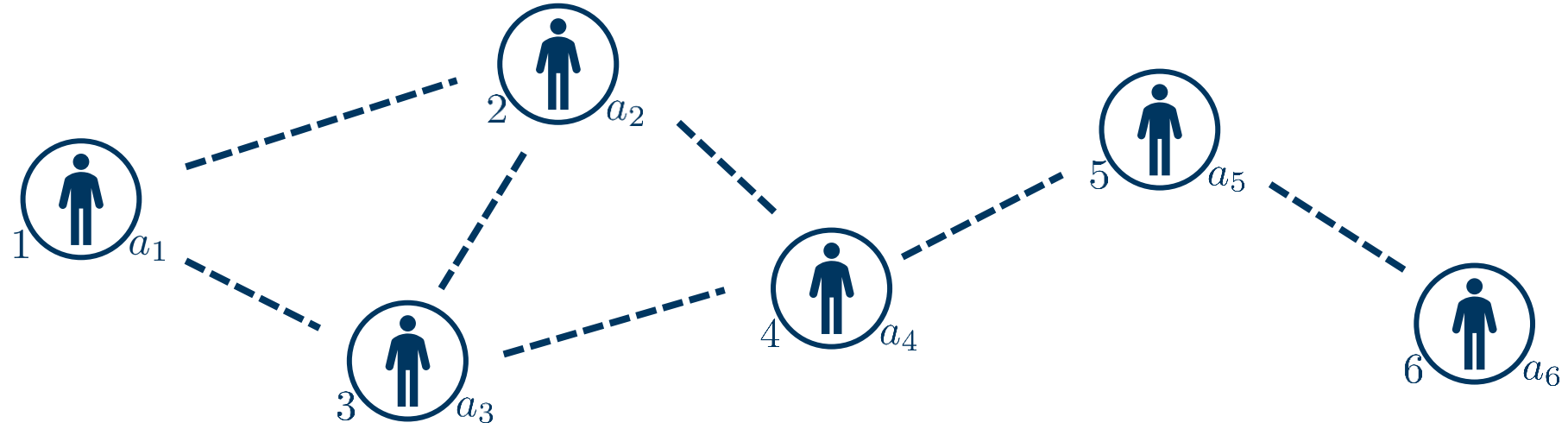
# Human Users in Multi-Agent Systems

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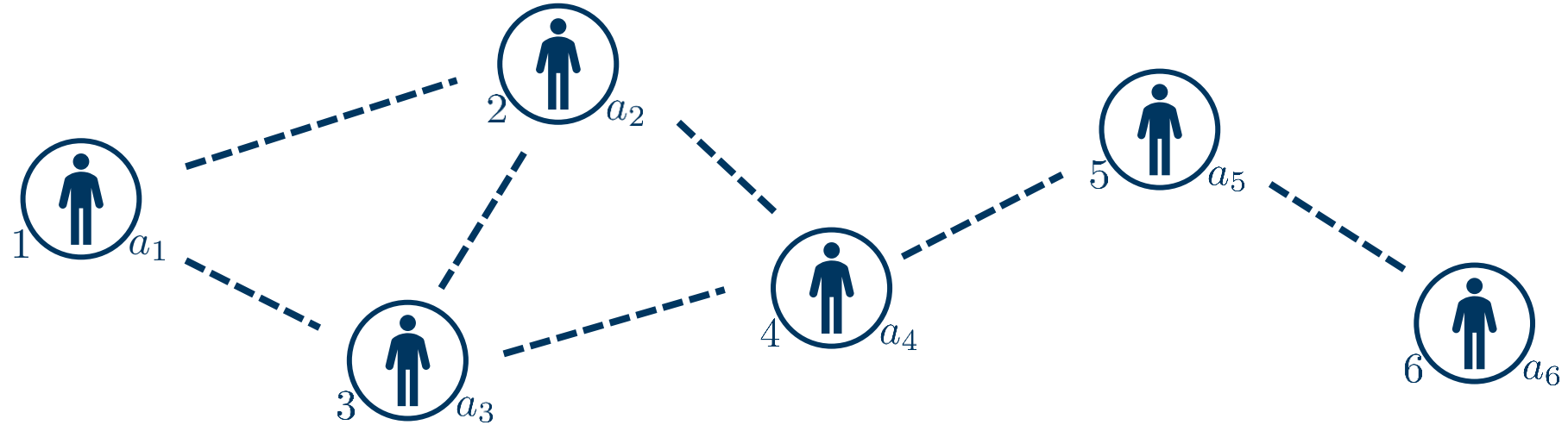
# Human Users in Multi-Agent Systems

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# Human Users in Multi-Agent Systems

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Preferences  
Beliefs



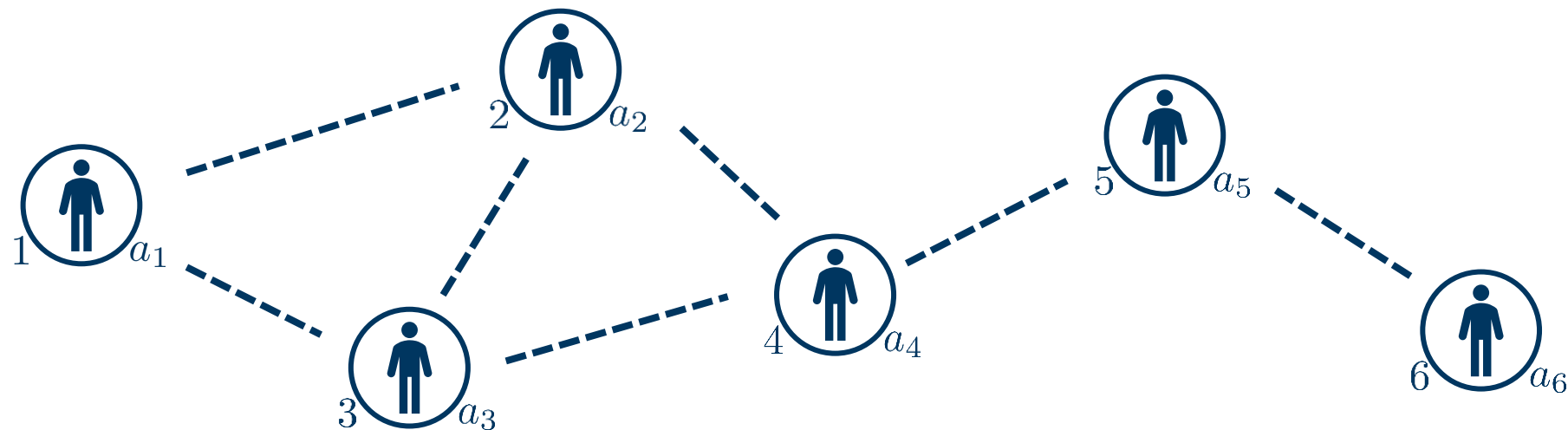
Human Decision  
 $a_i$



System Behavior  
 $a = (a_1, \dots, a_n)$

# Human Users in Multi-Agent Systems

---



Preferences  
Beliefs



Human Decision  
 $a_i$



System Behavior  
 $a = (a_1, \dots, a_n)$

How can we influence human users?  
*Preferences*                      *Beliefs*

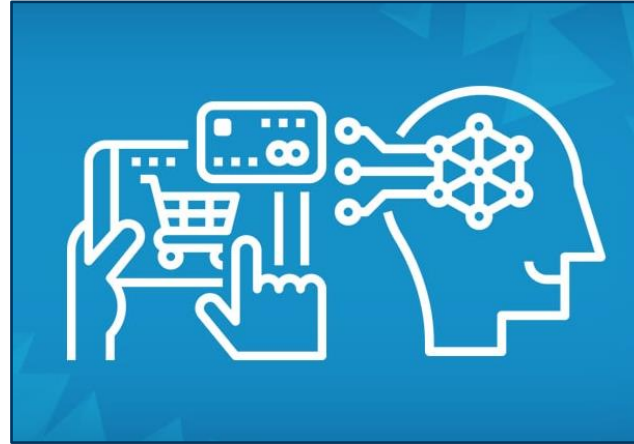
# Types of Information Communication (Social Systems)

## Agent-to-Agent



Social Networks and  
Viral Marketing

## Agent-to-Infrastructure



Recommender Systems  
and Advertising

## Unreliable Communicators



Belief Propagation and  
Fake News

Design communication mechanisms to alter *users' beliefs*  
and actions to ultimately guide system behavior

# Directions in Socio-Technical Systems



Intersection of *Engineered* and *Social* systems



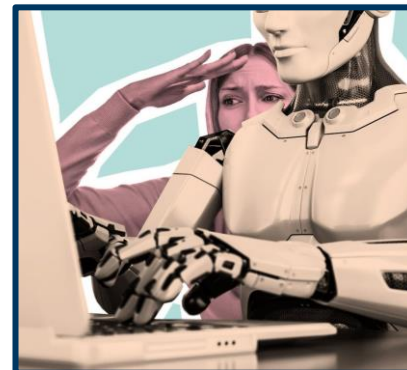
Smart services  
human demand



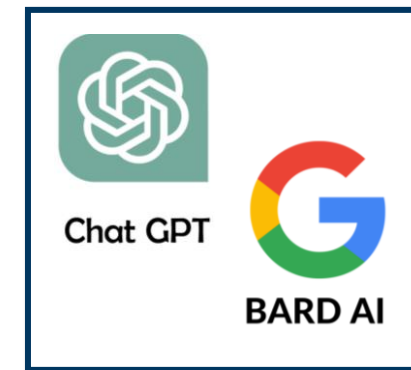
Autonomous  
mobility services



Human-robot  
interaction



Automation with  
human oversight



Human language  
AI interfaces



# Conclusion

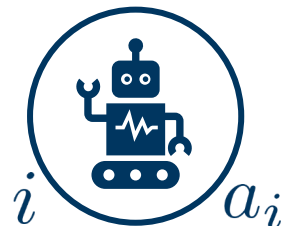
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## The Role of Communication in Distributed Systems

Share information to strategically alter system behavior

- |           |  |
|-----------|--|
| Benefits: | <ul style="list-style-type: none"><li>• Increase coordination</li><li>• Improve global objective</li><li>• Robustness to sub-system failures</li></ul> |
| Costs:    | <ul style="list-style-type: none"><li>• Increased complexity</li><li>• Unexpected behavior</li></ul>   |

Understanding benefits/costs can help us design distributed systems more intelligently



# References

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In Collaboration With:



Jason R. Marden  
UC Santa Barbara



Dario Paccagnan  
Imperial College London



Bary S. R. Pradelski  
CNRS (University of Grenoble)

- "Collaborative Coalitions in Multi-Agent Systems: Quantifying the Strong Price of Anarchy for Resource Allocation Games," **B. L. Ferguson**, D. Paccagnan, B. S. R. Pradelski, J. R. Marden. 62nd IEEE Conference on Decision and Control (submitted)
- "Robust Utility Design In Distributed Resource Allocation Problems With Defective Agents," **B. L. Ferguson** and J. R. Marden. Dynamic Games and Applications
- "Robust Utility Design In Distributed Resource Allocation Problems With Defective Agents," **B. L. Ferguson** and J. R. Marden. 60th IEEE Conference on Decision and Control
- "The Cost of Informing Decision-Makers in Multi-Agent Maximum Coverage Problems with Random Resource Values," **B. L. Ferguson**, D. Paccagnan, and Jason R. Marden. IEEE Control Systems Letters (under review)



UC SANTA BARBARA

# Information *as* Control: The Role of Communication in Distributed Systems

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May 15<sup>th</sup>, 2023

Supported by NSF, ONR, AFOSR, and ARL